



# Breaking the Illusion of Key Zeroization: How OS, Libraries, and Hardware Keep Your AES Keys Alive

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# Everything everywhere all encrypted

At Rest

In Transit

In Use

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At Rest

Data in storage is encrypted.

(Windows BitLocker, Mac FileVault, Linux LUKS)



In Transit

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In Transit

Data on network is encrypted.

(TLS, SSH, Signal, Telegram)



In Use

# Everything everywhere all encrypted

## At Rest

Data in storage is encrypted.

(Windows BitLocker, Mac FileVault, Linux LUKS)



## In Transit

Data on network is encrypted.

(TLS, SSH, Signal, Telegram)



## In Use

Data in memory is encrypted?

(Confidential Computing, Intel SGX, AMD SEV)

# Zeroization (key term)

- Zeroization: overwriting keys in memory.
  - Keys must be zeroized after use.

# Zeroization (key term)

- Zeroization: overwriting keys in memory.
  - Keys must be zeroized after use.
- Ransomware without zeroization is NOT threat.

# Why zeroization is important?

Zeroization is required from 3 perspectives

**defense-in-depth**

Residual keys expand the attack surface. Reliable zeroization reduces the chance of successful exploitation.

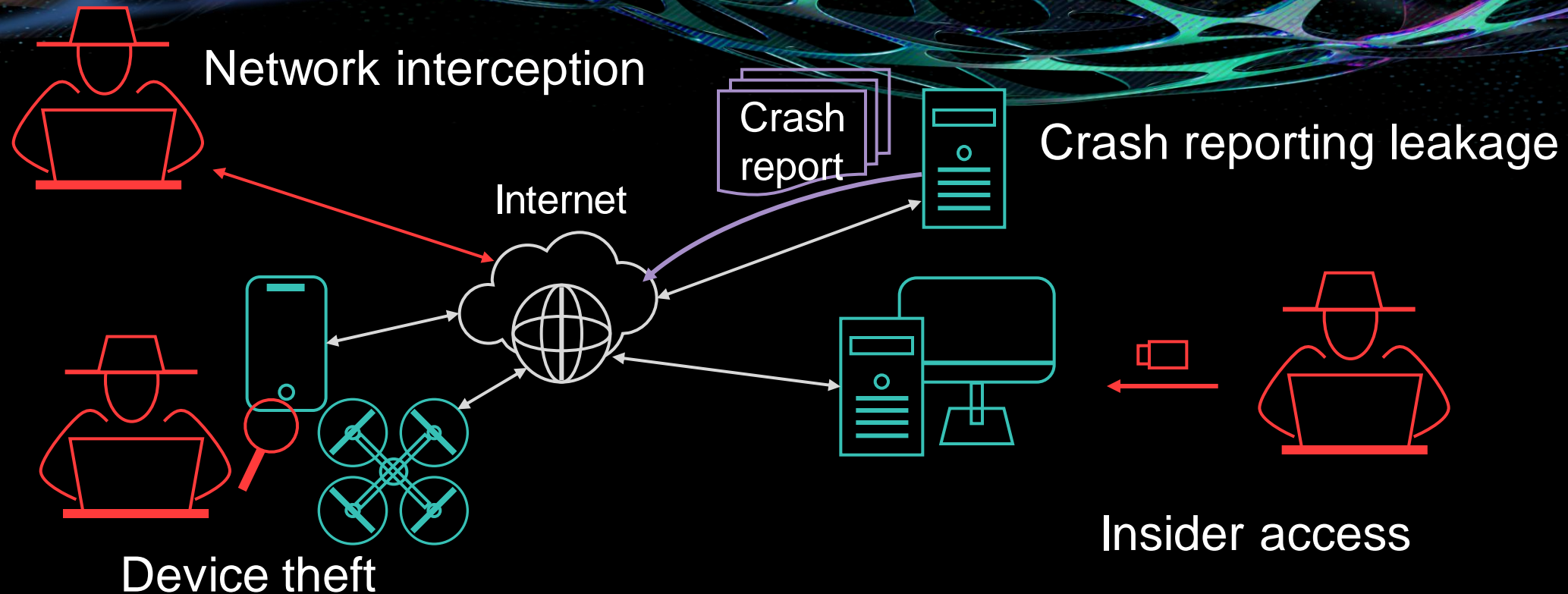
**forward secrecy**

If key lifetime is unknown, past communications may be exposed after system compromise.

**certification**

Standards require zeroization of Sensitive Security Parameters (e.g. FIPS 140-3).

# Potential key exposure scenarios



We focus on key zeroization as a post-compromise defense for maintaining forward secrecy.

# Think about a real-world compromise scenario

The system is compromised, and data is exfiltrated.

The data is safe because it is encrypted 😊

What about the keys? Did they exist in the system during the attack?

For executives:

Can we assure our customers and regulators that the data is safe?

For engineers:

Can we explain how keys are zeroized in our system?

# Do we assume that the keys simply disappear ?

OS and API seem to support zeroization

- free() clears memory
- process exit removes secrets
- reboot erases everything
- keys are not copied elsewhere
- the OS takes care of it

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## THIS IS THE ILLUSION OF ZEROIZATION

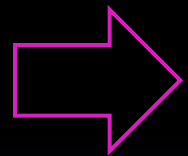
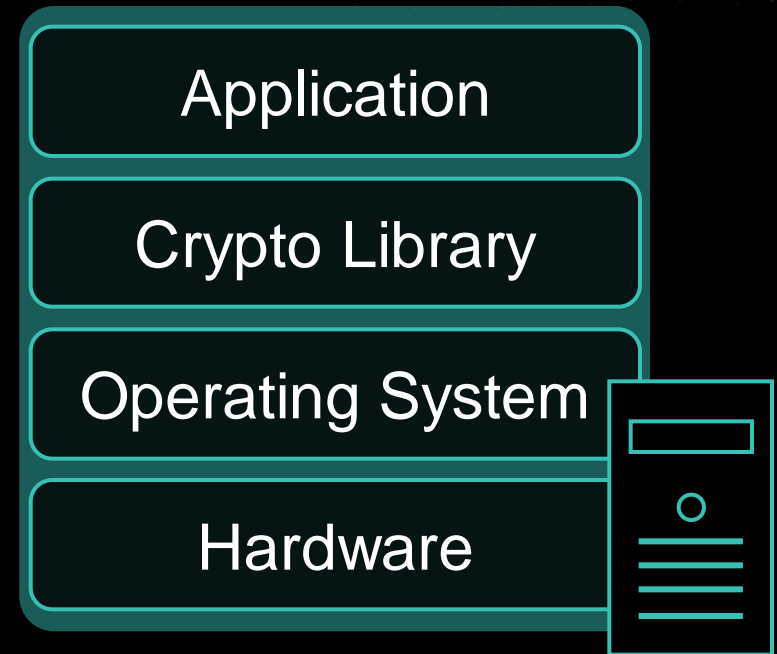
Breaking this illusion requires understanding the entire system stack.

compliance • cryptography • applications • libraries • OS • hardware

# Why a system-wide perspective is necessary?

Key lifetime is influenced by interactions across multiple system layers

- Zeroization is usually implemented in applications or libraries
- OS holds (copies) App's state



Zeroization can't be examined from user space only!

# Research questions: Understanding zeroization failures across layers

RQ1

When a process terminates **normally**, are the cryptographic keys in user space zeroized?

RQ2

When a process terminates **abnormally**, are the cryptographic keys in user space zeroized?

RQ3

Are cryptographic keys **duplicated into memory** regions that the process cannot read or write?

RQ4

Do residual keys persist after a **system reboot**?

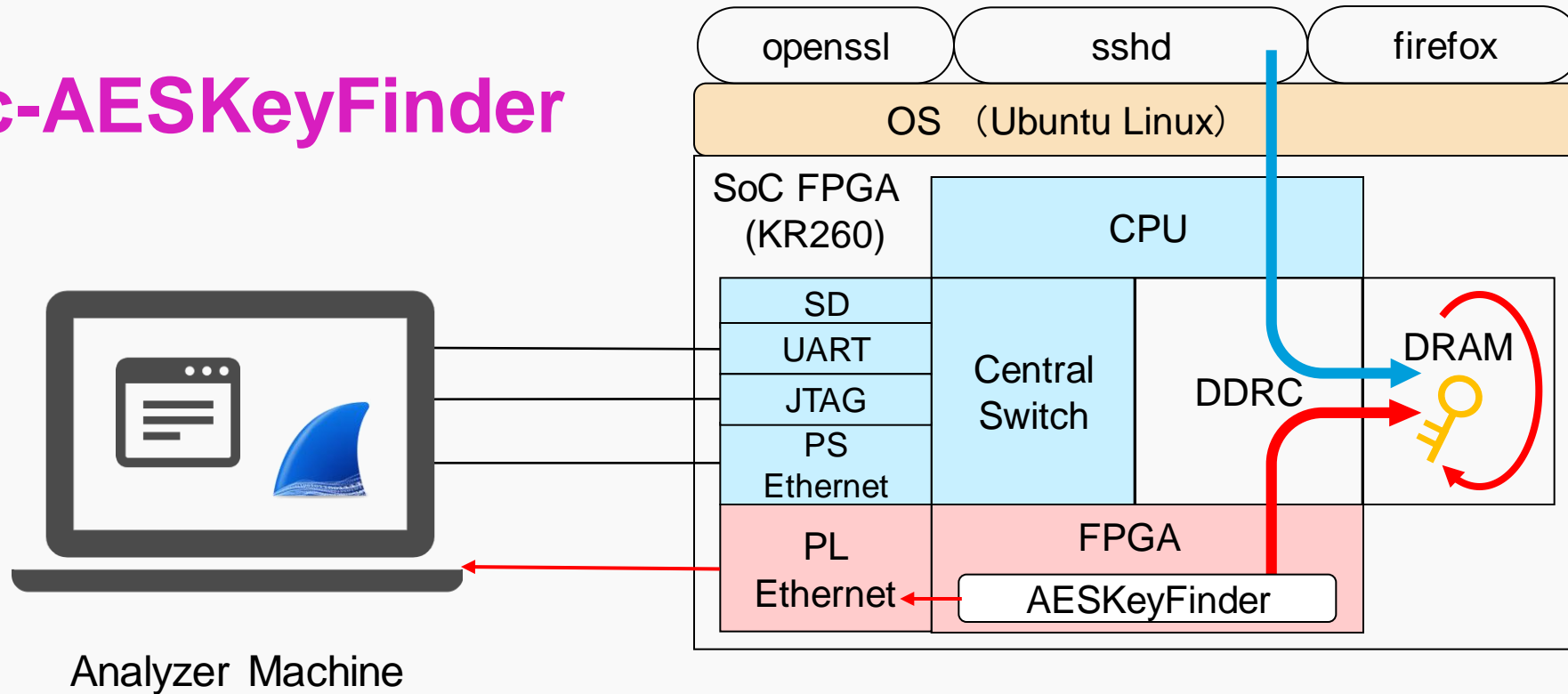
# Investigation method

We continuously scan physical memory using an FPGA.

This allows us to observe keys beyond OS control.

Our tool:

**Periodic-AESKeyFinder**



# Periodic-AESKeyFinder design

## FPGA-implemented AESKeyFinder

- AESKeyFinder is a forensic tool for AES keys, and is known for its cold boot attack paper
- Small Modification: Partial key schedule matching

Full key schedule

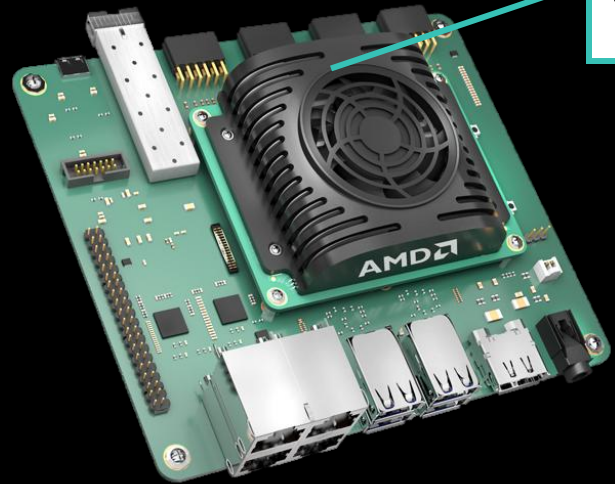
```
1 845e74b00: 0fc3b56c309d162d 55cbdd1966fd3458
2 845e74b10: eba2bc81e46109ed d89455c08d5f88d9
3 845e74b20: b5a297925e002b13 e0694a8b38fd1f4b
4 845e74b30: d643455363e1d2c1 0ed71093eebe5a18
5 845e74b40: 69099950bf4adc03 8960d3db87b7c348
6 845e74b50: 6fe4952506ed0c75 613385b6e853566d
7 845e74b60: 27e65ae74802cfc2 ae86893ccfb50c8a
8 845e74b70: 8400d1c2a3e68b25 e53354744bb5dd48
9 845e74b80: b53f9947313f4885 1bb9107bfe8a440f
10 845e74b90: a5af871310901e54 409cd3675b25c31c
11 845e74ba0: 303647179599c004 2b8f576c6b13840b
12 845e74bb0: 0000000000000000 0000000000000000
13 845e74bc0: 0000000000000000 0000000000000000
```

Partial key schedule

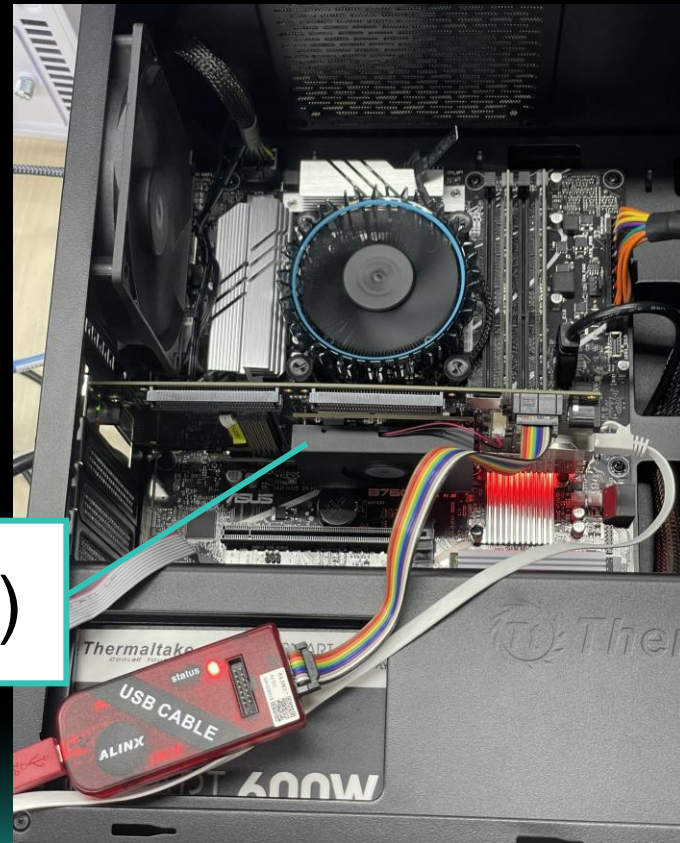
```
1 8033e9d90: 0fc3b56c309d162d 55cbdd1966fd3458
2 8033e9da0: eba2bc81e46109ed d89455c08d5f88d9
3 8033e9db0: 05f006f2f39d44f3 2bf60ccac89f43a6
4 8033e9dc0: 17f6ea6224e7872e cad13992c8c8f466
5 8033e9dd0: 27e65ae74802cfc2 ae86893ccfb50c8a
6 8033e9de0: 8400d1c2a3e68b25 e53354744bb5dd48
7 8033e9df0: b53f9947313f4885 1bb9107bfe8a440f
8 8033e9e00: a5af871310901e54 409cd3675b25c31c
9 8033e9e10: 05f006f2f39d44f3 2cf60ccac89f43a6
10 8033e9e20: 05f006f2f39d44f3 2df60ccac89f43a6
11 8033e9e30: d823a4ef35a44bdc 99101d42a81c5753
12 8033e9e40: d3e21b5769b7e96a 89ab4b04e6efd5ae
13 8033e9e50: 0000000000000000 0000000000000000
```

# Experimental setups

SoC FPGA



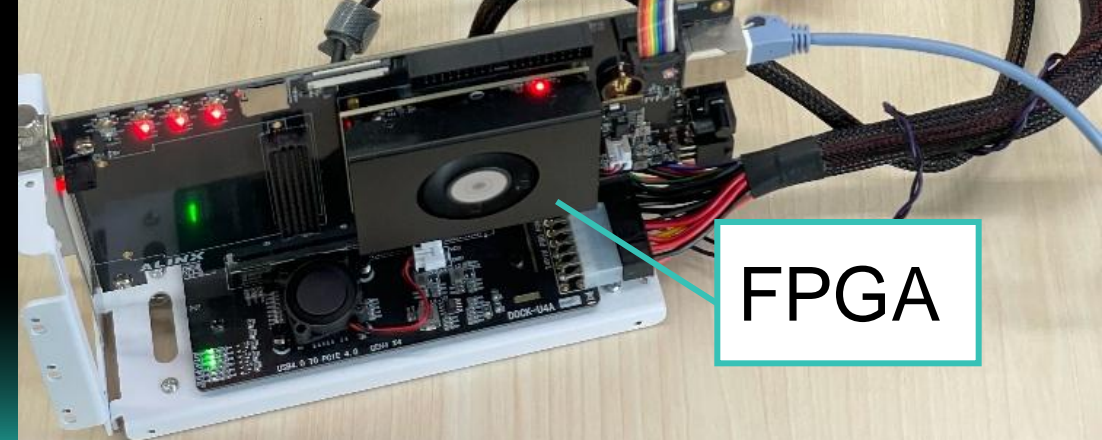
Intel Desktop (PCIe)



Laptop (Thunderbolt)



FPGA



# Demonstration

We test whether AES keys remain in memory after three actions.

**RQ1: Normal** termination: Exit the web browser with close button

**RQ2: Abnormal** termination: Fore terminate the browser (pkill)

**RQ4: Reboot** experiment: Reboot the PC

No.	Time	Source	Destination	Protocol	Length	Info
1568	3.842486579	Xilinx_32:36...	Broadcast	KeyFin...	64	AES128 found! Key count=1
1571	3.842487168	Xilinx_32:36...	Broadcast	KeyFin...	64	AES128 found! Key count=1
1572	3.842487255	Xilinx_32:36...	Broadcast	KeyFin...	64	AES128 found! Key count=1
1574	3.842487335	Xilinx_32:36...	Broadcast	KeyFin...	64	AES128 found! Key count=1
1575	3.842487416	Xilinx_32:36...	Broadcast	KeyFin...	64	AES128 found! Key count=1

Use the **Ocean-colored keys** as a makers for scan cycles. They appear every ~2 seconds, indicating one full scan of 8 GB memory.

```
Frame 15: 64 bytes on wire (512 bits), 64 bytes captured (512 bits) on interface en...
Ethernet II, Src: Xilinx_32:36:30 (00:5d:03:32:36:30), Dst: Broadcast (ff:ff:ff:ff:ff:ff)
...
Destination: Broadcast (ff:ff:ff:ff:ff:ff)
...
Source: Xilinx_32:36:30 (00:5d:03:32:36:30)
...
Type: Unknown (0x7020)
...
KeyFinder
Packet Number: 10
Number of keys found: 0
Elapsed Cycles: 00000003cb80e0c
Elapsed Laps: 2
Packet Size: 64
- Key Information
  - AES Key
    AES Key Length: 128
    KeySchedule Bitmap: 07ff
    Physical Address: 0106938f0
    Raw Key: 6e0f6785b882b247b7b0be8bed206e677
```

# Demonstration

GNU GRUB version 2.06

```
Ubuntu
Advanced options for Ubuntu
Windows Boot Manager (on /dev/nvme0n1p1)
UEFI Firmware Settings
*Periodic-KeyFinder
```

No.	Time	Source	Destination	Protocol	Length	Info
-----	------	--------	-------------	----------	--------	------

Use the ↑ and ↓ keys to select which entry is highlighted.  
Press enter to boot the selected OS, 'e' to edit the commands before booting or 'c' for a command-line. ESC to return previous menu.

# Demonstration of zeroization failures

Experiment with multiple combinations of apps, libraries, operating systems, and architectures

Project	Library	OS	Exit	SIGKILL	SIG-term	SIG-core	Context Switch	Core dump
OpenSSH	OpenSSL	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✗	✗
OpenSSH	OpenSSL	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	✗
OpenSSH	LibreSSL	Windows 11 (x86-64)	✓	✗	-	-	✓	-
OpenSSL	OpenSSL	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✗	✗
OpenSSL	OpenSSL	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	✗
Firefox	LibNSS	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✗	✗
Firefox	LibNSS	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	✗
Chromium	BoringSSL	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✓	-
Chromium	BoringSSL	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	-
Edge	Schannel	Windows 11 (x86-64)	✓	✗	✗	-	✓	-
7-Zip	Original	Windows 11 (x86-64)	✗	✗	✗	-	✓	-

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OpenSSH	OpenSSL	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	✗
OpenSSH	LibreSSL	Windows 11 (x86-64)	✓	✗	-	-	✓	-
OpenSSL	OpenSSL	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✗	✗
OpenSSL	OpenSSL	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	✗
Firefox	LibNSS	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✗	✗
Firefox	LibNSS	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	✗
Chromium	BoringSSL	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✓	-
Chromium	BoringSSL	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	-
Edge	Schannel	Windows 11 (x86-64)	✓	✗	✗	-	✓	-
7-Zip	Original	Windows 11 (x86-64)	✗	✗	✗	-	✓	-

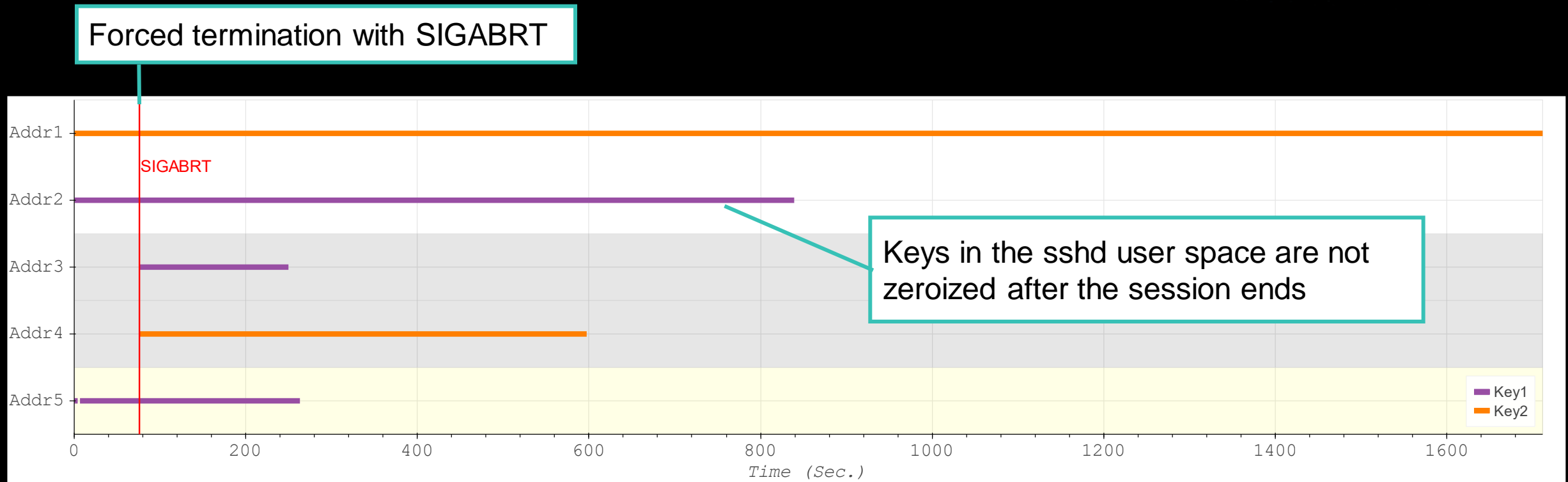
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OpenSSH	LibreSSL	Windows 11 (x86-64)	✓	✗	-	-	✓	-
OpenSSL	OpenSSL	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✗	✗
OpenSSL	OpenSSL	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	✗
Firefox	LibNSS	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✗	✗
Firefox	LibNSS	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	✗
Chromium	BoringSSL	Ubuntu 22.04 (AArch64)	✓	✗	✗	✗	✓	-
Chromium	BoringSSL	Ubuntu 22.04 (x86-64)	✓	✗	✗	✗	✓	-
Edge	Schannel	Windows 11 (x86-64)	✓	✗	✗	-	✓	-
7-Zip	Original	Windows 11 (x86-64)	✗	✗	✗	-	✓	-

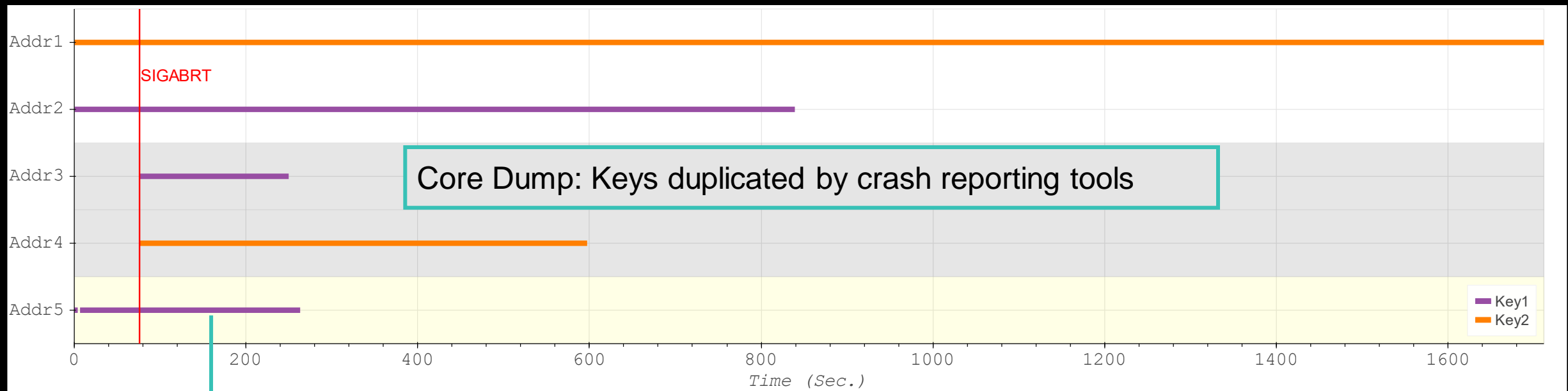
# Zeroization upon abnormal app. termination (RQ2)

- Key appearance timeline when sending SIGABRT to sshd during an SSH session



# Duplication of detected cryptographic keys (RQ3)

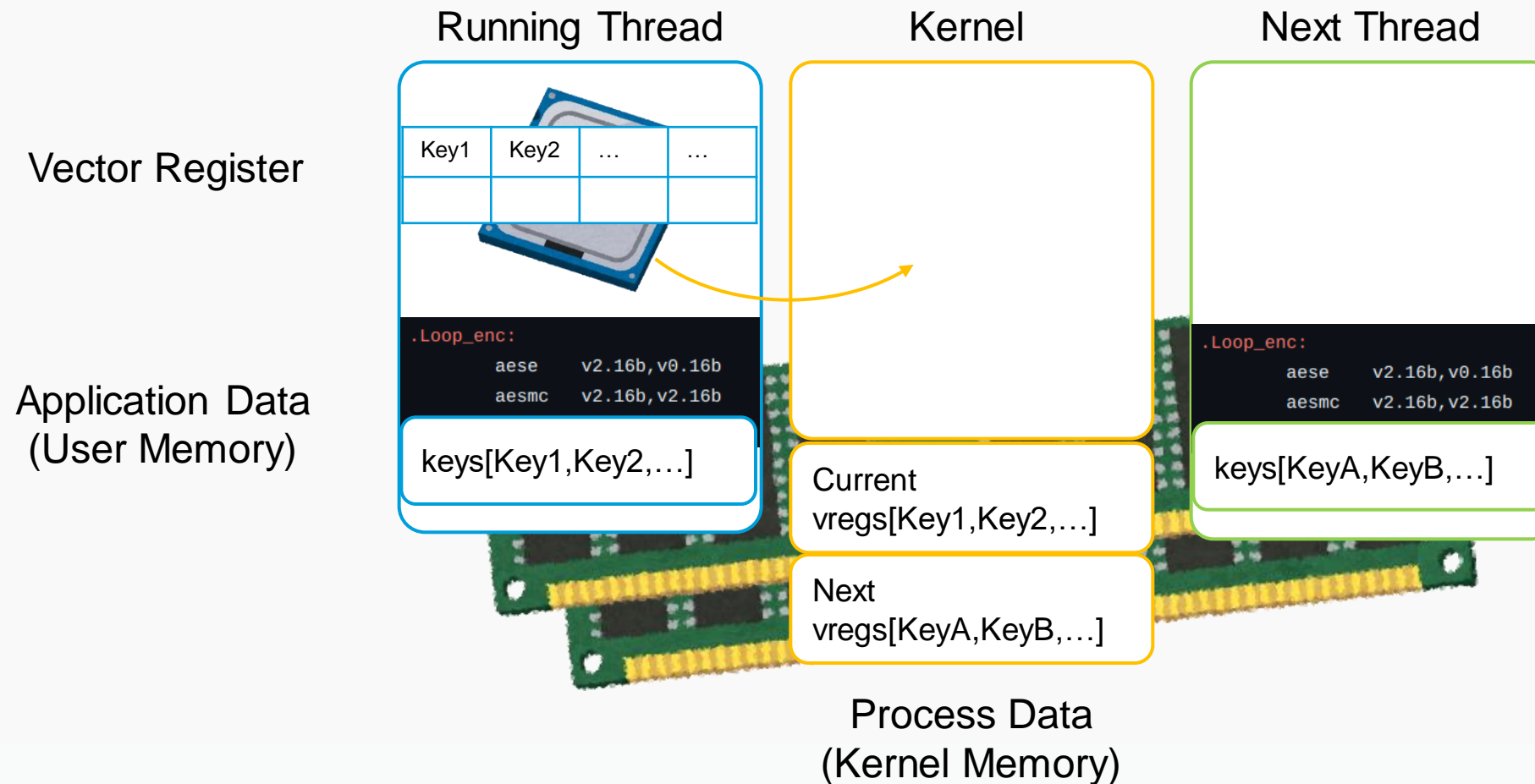
- Core Dump
- Context Switch



Keys duplicated due to context switches

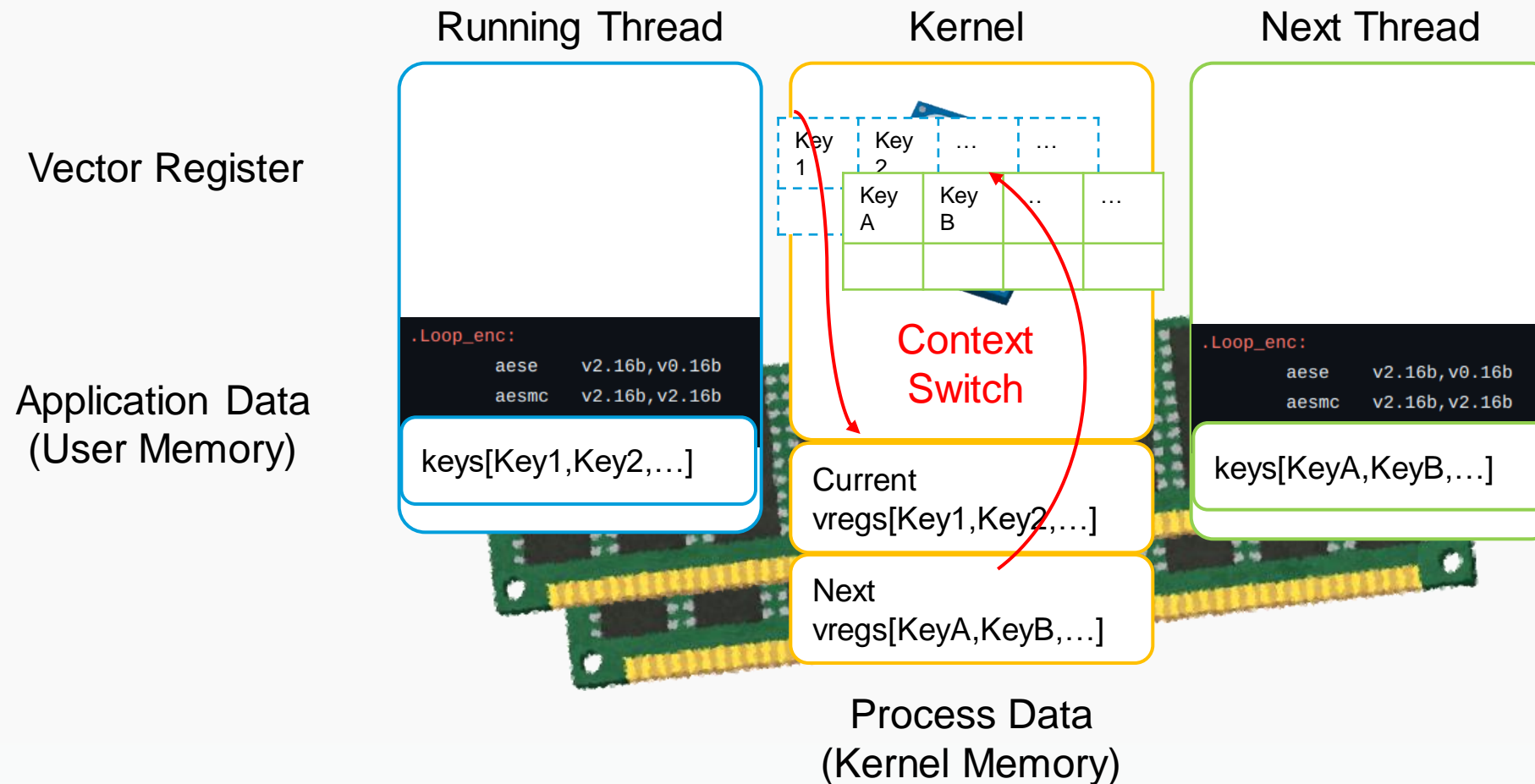
# Key duplication caused by vector-register context switching (1/3)

- Context switch: switching threads running on the CPU
- When an interrupt occurs, the running thread yields the CPU to the kernel.



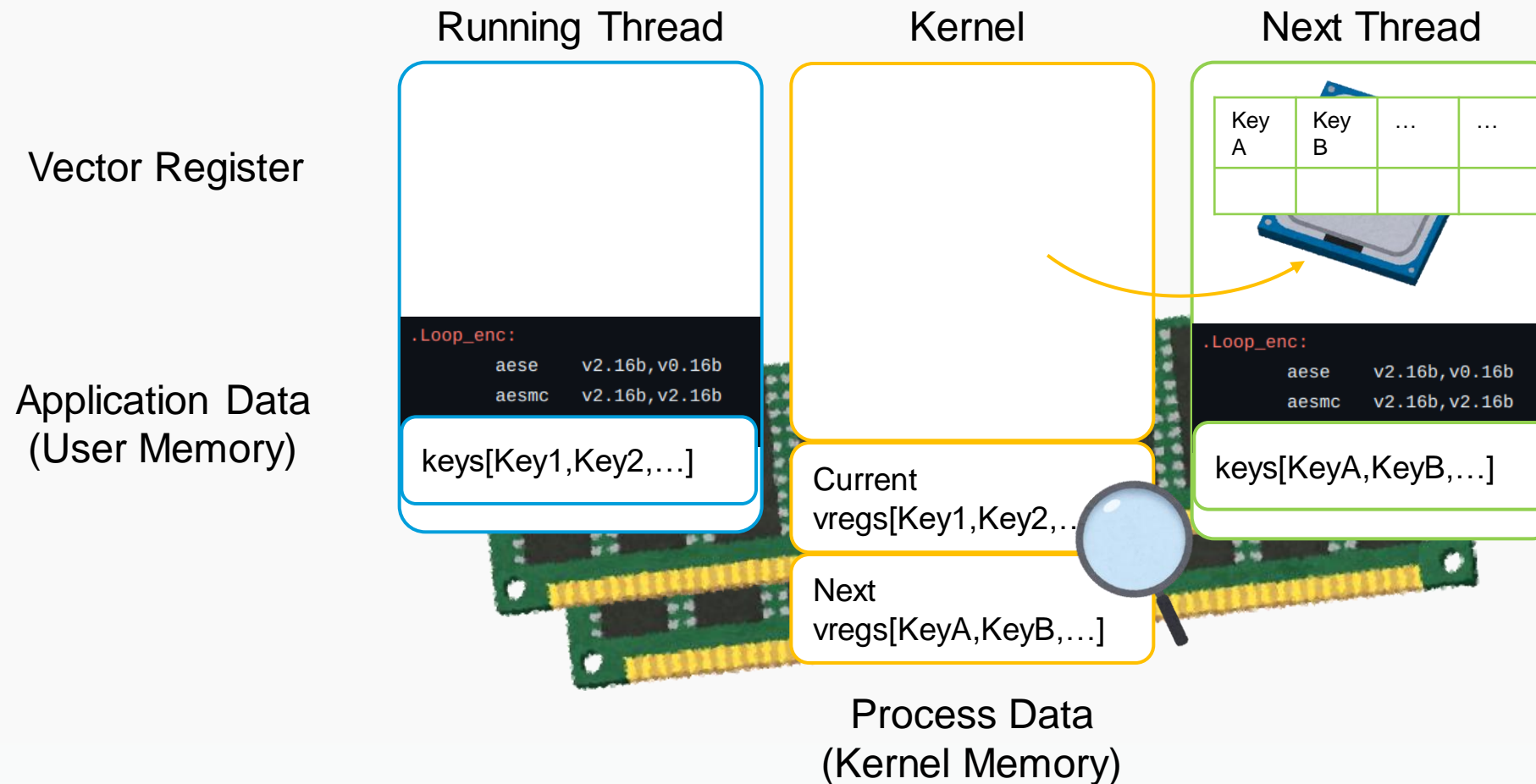
# Key duplication caused by vector-register context switching (2/3)

The register contents of the previous process are saved into kernel memory, and the next process's register contents are loaded into the CPU to resume execution.



# Key duplication caused by vector-register context switching (3/3)

Keys duplicated into kernel space during context switching were observed.



# Key persistence after OS reboot (RQ4)

- FIPS-compliant Linux zeroizes on reboot.
- Other environments sometimes retain keys.

OS (Environment)	Warm boot	After kernel crash	Reset switch
Linux 5.15.0-1038-Xilinx-zynqmp(SoC)	X	X	X
Linux-5.15.0-130-fips (Intel Desktop)	✓	✓	✓
Linux-6.8.0-50-generic (Intel Desktop)	X	✓	✓
Linux-6.8.0-50-generic (AMD Desktop)	✓	✓	✓
Linux-6.8.0-50-generic (Laptop)	X	✓	-
Windows 11 Build 22621 FIPS (Intel Desktop)	X	✓	✓
Windows 11 Build 22621 (Intel Desktop)	X	✓	✓

# Key persistence after OS reboot (RQ4)

- FIPS-compliant Linux zeroizes on reboot.
- Other environments sometimes retain keys.
- PCs zeroize memory after crash or hardware reset.

OS (Environment)	Warm boot	After kernel crash	Reset switch
Linux 5.15.0-1038-Xilinx-zynqmp(SoC)	X	X	X
Linux-5.15.0-130-fips (Intel Desktop)	✓	✓	✓
Linux-6.8.0-50-generic (Intel Desktop)	X	✓	✓
Linux-6.8.0-50-generic (AMD Desktop)	✓	✓	✓
Linux-6.8.0-50-generic (Laptop)	X	✓	-
Windows 11 Build 22621 FIPS (Intel Desktop)	X	✓	✓
Windows 11 Build 22621 (Intel Desktop)	X	✓	✓

# Key persistence after OS reboot (RQ4)

- FIPS-compliant Linux zeroizes on reboot.
- Other environments sometimes retain keys.
- PCs zeroize memory after crash or hardware reset.
- Keys remain on the SoC even after resetting.

OS (Environment)	Warm boot	After kernel crash	Reset switch
Linux 5.15.0-1038-Xilinx-zynqmp(SoC)	X	X	X
Linux-5.15.0-130-fips (Intel Desktop)	✓	✓	✓
Linux-6.8.0-50-generic (Intel Desktop)	X	✓	✓
Linux-6.8.0-50-generic (AMD Desktop)	✓	✓	✓
Linux-6.8.0-50-generic (Laptop)	X	✓	-
Windows 11 Build 22621 FIPS (Intel Desktop)	X	✓	✓
Windows 11 Build 22621 (Intel Desktop)	X	✓	✓

# Who is responsible for zeroization?

Key lifetime is influenced by interactions across multiple system layers.

- Missing signal handlers
- Core dumps
- Lazy page zeroization
- Duplication via register
- No zeroization on reboot



# Recommendations and mitigations

## App. Developers

Implement signal-handler zeroization for abnormal termination.  
Exclude sensitive pages from core dumps.

## Library Developers

Zeroize temporary key material stored in registers and temporary buffers after cryptographic operations.

## OS Developers

Zeroize the kernel memory that stores saved register contexts when a process exits.  
Implement zeroization for SIGKILL.

# Recommendations and mitigations

## System Managers

Configure the system to not generate core dumps.  
Configure crash-reporting tools to avoid external transmissions.

## Hardware / Platform Vendors

Provide clear documentation on memory zeroization behavior during reboot.

## Certification Testers

Understand these system-level zeroization issues and require vendors to document how their platforms behave.

# Takeaways

- **Zeroization must be verified by observing physical memory.**
- **Keys persist, duplicate, and survive reboot in real systems.**
- **Reliable zeroization requires a system-wide approach.**

# Conclusion

- Zeroization spans compliance, applications, OS, hardware, and cryptography.
- FPGA-based live forensics revealed duplication and persistence across layers and reboots.
- We responsibly disclosed these issues to real-world systems.
- Code & tools available on GitHub

[GitHub - Tyojan/Periodic-AESKeyFinder](#)



Thank you for listening!