



UWB Real Time Locating Systems How Secure Radio Communications May Fail in Practice

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Overview







Introduction



Introduction

Wireless communication systems are susceptible to various security threats that can compromise their reliability and impact production operations

UWB

- Ideal for short-range devices
- Can transmit information quickly over short distances
- Ability to send data through solid objects like walls and other barriers
- UWB is the preferred communication protocol for RTLS



Spectral density for UWB and narrowband - Source: FiRa Consortium



Introduction – Cont'd

Wireless communication systems are susceptible to various security threats that can compromise their reliability and impact production operations



Operation of an UWB TDoA RTLS

RTLS

- Uses UWB signals to locate stationary/mobile objects
- 3 components:
 - Tags
 - Anchors
 - Server that computes, shows and stores tag positions
- The time of arrival is analyzed to determine the position of a tag



Introduction – Cont'd

Wireless communication systems are susceptible to various security threats that can compromise their reliability and impact production operations



Cyber Threats

Networks will be **vulnerable** to attacks by cyber criminals who are seeking to exploit vulnerabilities in order to gain access to sensitive data or disrupt operations.

Source: FiRa Consortium



Motivation



Motivation

According to the FiRa consortium, in 2018 there was an increased demand for "improvements to existing modulations to increase the integrity and accuracy of ranging measurements"

- In 2020, the Institute of Electrical and Electronic Engineers (IEEE) released standard 802.15.4
- IEEE quickly followed up with the **802.15.4z amendment**, also released in 2020
- Synchronization and exchange of location data are considered "out-of-scope" by the standard
- These communications, whose design is left entirely to vendors, are critical aspects for the **overall posture of TDoA RTLS**
- Additionally, there has not been any research on UWB focusing on this specific problem

This is what motivated us to take a deeper look into how threat actors can exploit this vulnerability and disrupt environments utilizing UWB RTLS.





Industry Scope

Focus: industrial and healthcare sectors

- Highly targeted
- UWB RTLS widely used in critical applications

Examples: use cases where UWB is protecting people's lives

- Employee and patient tracking
- Geofencing
- Contact tracing

Simatic RTLS: How to create a safe working environment

During the corona pandemic, many companies are faced with question of how to keep their production running without putt the safety of employees at risk. Simatic RTLS allows for socia distancing in production and other worksites. I had a recent ta with Nicole Lauther, Siemens USA, who is coordinating the gk roll-out of the solution.



The Siemens solution can identify dangerous bottlenecks in the production environment where en

often have to spend time without keeping the necessary distance.

📮 **Ubi**sense

WHAT WE DO

Worker Safety

Ubisense tracks the real-time location and movement of both people and equipment across indoor and outdoor spaces for worker safety applications.

- Easily define 3-dimensional safe and unsafe spaces
- Alert workers in real-time to prevent accidents
- Avoid collisions by automatically disabling equipment
- Locate personnel quickly in hazardous environments
 Decord cafety incidents for angoing training
- Record safety incidents for ongoing training

Ubisense Dimension4™ sensors have been engineered to work reliably 24/7, even in harsh industrial environments. Uncompromising performance for employee safety.

LEARN MORE

Collision Avoidance

Ubisense Dimension4[™] tags and SmartSpace[®] software connect seamlessly with existing systems and devices to give complete coverage of work zones and immediate control to prevent collisions.

 Track the real-time movement of overhead cranes or heavy equipment and disable operation to prevent collisions

Examples of UWB RTLS use cases advertised by vendors



Analyzed Solutions

 Our research was performed on the following solutions, which target the industrial and healthcare sectors:





Sewio Indoor Tracking RTLS UWB Wi-Fi Kit Avalue Renity Artemis Enterprise Kit

• Both these UWB RTLS kits come equipped with a set of tags, anchors, and a server software that provides the **aforementioned safety features**



Technical Scope

- Network communications occurring in a normal RTLS infrastructure:
 - \circ UWB
 - Tags to anchors
 - Anchors to anchors
 - Ethernet/Wi-Fi/other
 - Anchors to RTLS server



As in a chain, a flaw in **any** of these communications may lead to a compromise of the **entire** infrastructure

• Up to now, security research has exclusively focused on the analysis of UWB signals. This is the first research analyzing the communications on Ethernet/Wi-Fi/etc.



Architecture of an UWB TDoA RTLS



TDoA Background and Theory



Packet Taxonomy

In a TDoA RTLS there are normally two kinds of packets exchanged among anchors and server



Synchronization packets

- Anchors' clocks are usually not in-sync, (different boot times, clock drifts, etc.)
- A reference anchor continuously sends an UWB signal that is received by all nonreference anchors
- The reference anchor sends a synchronization packet containing the transmission timestamp, the non-reference anchors a synchronization packet containing the receiving timestamp
- The server uses this information to build a common notion of time



Positioning packets

- A tag emits an UWB signal that is received by all anchors
- All anchors send the timestamps at which they received the UWB signal from the tag to the central positioning server, inside positioning packets



Algorithm Details

The routine implemented in TDoA RTLS can usually be divided in two major phases



Clock Synchronization

- There are many synchronization algorithms in literature. This work used the Linear Interpolation algorithm
- The basic idea is to convert all anchors' timestamps to a common clock domain, so that they can be compared. These converted timestamps are called Global Times (GT)



Position Estimation

- We cannot immediately derive the distances from the GTs as we are missing the transmission instants
- We can, however, correlate the difference of GTs to the difference of distances. This is why the algorithm is called Time Difference of Arrival

Delta(i,j,t) = (GT(reference,j,t) - GT(i,j,t)) * c = GT(reference,j,t) * c - GT(i,j,t) * c = d(reference,j,t) - d(i,j,t)

All details in our whitepaper! Download it from the briefing page, or from the nozominetworks.com website



Algorithm Details – Cont'd

The routine implemented in TDoA RTLS can usually be divided in two major phases

Position Estimation – Cont'd

• Eventually, a non-linear system of equations can be set up to solve for Xj,t, Yj,t, and Zj,t, i.e., the position of tag j at the instant t:

 $\left(pTs(reference, j, t) - sTS(reference, t) - (CS(1, t) * (pTs(1, j, t) - sTS(1, t)) + ToF(1)) \right) * c = \sqrt{(Xj, t - Xreference)^2} + (Yj, t - Yreference)^2 + (Zj, t - Zreference)^2 - \sqrt{(Xj, t - X1)^2} + (Yj, t - Y1)^2 + (Zj, t - Z1)^2$...

 $\left(pTs(reference, j, t) - sTS(reference, t) - (CS(N, t) * (pTs(N, j, t) - sTS(N, t)) + ToF(N)) \right) * c = \sqrt{(Xj, t - Xreference)^2 + (Yj, t - Yreference)^2 + (Zj, t - Zreference)^2} - \sqrt{(Xj, t - XN)^2 + (Yj, t - YN)^2 + (Zj, t - ZN)^2}$



Summary

To obtain the position of a tag, the following data need to be known:

- All coordinates of the anchors involved
- Synchronization timestamps
- Positioning timestamps



Reverse Engineering of Devices Network Traffic

Network Traffic

 Both solutions use custom, unknown binary network protocols for the communications among anchors and server. No standard data structures are immediately recognizable

	270 5000 - 5000 Len		
91 4.010120831 192.108.225.11 192.108.225.2 UDP	278 5000 → 5000 Len	1=230	
94 4.013749191 192.168.225.12 192.168.225.2 UDP	332 5000 → 5000 Len	1=290	
95 4.021486274 192.168.225.14 192.168.225.2 UDP	332 5000 → 5000 Len	h=290	
96 4.027134239 192.168.225.15 192.168.225.2 UDP	96 5000 → 5000 Len	1=54	
97 4.029864276 192.168.225.13 192.168.225.2 UDP	96 5000 → 5000 Len	n=54	
98 4.030089847 192.168.225.11 192.168.225.2 UDP	96 5000 → 5000 Len	1=54	
99 4 061396680 192 168 225 14 192 168 225 2 UDP	96 5000 → 5000 Len	3=54	
100 / 067518350 102 168 225 15 102 168 225 2 UDD	96 5000 → 5000 Len	-54	
101 4 060455244 102 160 225 12 102 160 225 2 UDD	96 5000 - 5000 Lon		
	90 5000 → 5000 Len	1-04	
102 4.009904122 192.108.225.11 192.108.225.2 UDP	96 5000 → 5000 Len	1=54	
103 4.0/3242/10 192.168.225.12 192.168.225.2 UDP	96 5000 → 5000 Len	1=54	
107 4.247084388 192.168.225.15 192.168.225.2 UDP	96 5000 → 5000 Len	1=54	
108 4.248879109 192.168.225.13 192.168.225.2 UDP	96 5000 → 5000 Len	1=54	
Erame 95: 332 bytes on wire (2656 bits) 332 bytes cantured	(2656 bits) on interfac 00	100 1c 69 7a a2 fa ec 68 27 19 87 7c 9f 08 00 45 00 jzb'	
Ethornot TT Src: Microchi $9f_7c_00f_6000000000000000000000000000000000$	t: ElitoGro a2:fa:ec (1c 00		
Filternet II, Stc. Microchi_of.70.91 (00.27.19.01.70.91), DS Internet Drotocol Version 4, Srev 102 160 225 14, Det. 102	160 225 2		# 16
Internet Protocol Version 4, Src. 192.108.225.14, DSt. 192.	108.225.2 00		- # 0.
Viser Datagram Protocol, Src Port: 5000, Dst Port: 5000		130 /C 8T 19 2/ 68 55 01 21 00 aa TD 6/ 8T 19 2/ 68 1 NU !	∵g∵n
✓ Data (290 bytes)		140 C4 T1 2D 03 22 D4 07 00 00 00 55 04 08 0T 2C 00 ···+·"···	••••••
Data: 23d12736009f7c8f19276855012100aafb678f192768c4f12b0)322b4070000005504080f 00	050 14 Of b1 08 00 05 f4 00 38 bb 06 01 00 42 07 04	8 · · · · B · ·
[Length: 290]	00	060 00 b2 f7 d6 95 23 23 98 36 00 9f 7c 8f 19 27 68 ## 6	6 · · · · 'h
		070 55 01 21 00 aa fb 67 8f 19 27 68 c5 f1 22 03 fb U ! g	• 'h • • " • •
	00	80 b7 07 00 00 00 6c 04 ed 11 28 00 f0 10 9b 08 91	. (
		90 05 ef 00 3f bb 06 01 00 42 07 04 00 b2 f7 d6 95 ····	в
		3a0 23 dc 69 36 00 9f 7c 8f 19 27 68 55 01 21 00 aa # 16 1	. 'hll. L
		f_{10} =	
		100 06 01 00 43 07 04 00 D2 T7 06 95 23 T6 6D 36 00 C	•••#•Кб•
		00 91 /C 81 19 27 68 55 01 21 00 aa fb 67 81 19 27 'NU	! · · · g · · '
	00	0f0 68 c7 f1 15 03 ad bf 07 00 00 00 f6 03 fb 0f 28 h	• • • • • • • • • (
		100 00 a6 0f 91 0b 72 05 ef 00 2a bb 06 01 00 43 07 ····r·	· * · · · · C ·
		110 04 00 b2 f7 d6 95 23 df a2 31 00 9f 7c 8f 19 27 ·····#·	1 ' '
	01	120 68 55 02 20 00 bb de 8c fd 5f 04 22 c6 ea cb 50 hU	"p
	01	130 d9 07 00 00 00 f7 01 47 0c 18 00 ae 17 6b 0f 78 ·····G	····k·x
	1 01	140 04 75 00 f2 b9 07 04 00 b2 f7 d6 95	

Example of Ethernet network packet in Sewio RTLS



Packet Dissection

• By reverse engineering the server software, full packet structure was reconstructed



Example of code snippets from Sewio and Avalue RTLS servers

Sewio and Avalue RTLS dissectors



We are freely releasing PCAPs and dissectors for both Sewio and Avalue RTLS! Download them from the **briefing page**, or from **github.com/NozomiNetworks**



Security Considerations



Confidentiality

- No confidentiality in the anchors-server communications
 - The synchronization and positioning timestamps are sent in cleartext



Integrity

- No secure integrity mechanisms either
 - 。 Sewio RTLS uses CRC-16
 - Avalue RTLS performs a byte per byte sum of all packets



Sewio RTLS – Extraction of timestamps directly from the network traffic





Anchors Coordinates Prerequisite



Anchor Coordinates Prerequisite

Obtaining the anchor coordinates is the most challenging requirement. They are manually input at the first installation and never transmitted through the network



Attacker with Physical Access

- If the anchors are visible, obtaining their coordinates is simple
- If not, an attacker can still produce an estimation by measuring the power levels of the anchors' wireless signals (UWB, Wi-Fi, etc.)
- In fact, tag coordinates can be estimated even with imperfect anchor coordinates



If anchor coordinates are estimated with a <10% error, the tag coordinates are computed with an average error of <20%, i.e., **~50 cm** in a 6m x 5m room

Tag Coordinates Average Error with respect to Anchor Coordinates Error





Obtaining the anchor coordinates is the most challenging requirement. They are manually input at the first installation and never transmitted through the network



Attacker with Remote Access

 Besides timestamps, anchors transmit on the wire the power level information of the received UWB signal. We can compute two different metrics:

• First Path Power Level (FPPL)

$$FPPL = 10 * \log_{10} \left(\frac{FP1^2 + FP2^2 + FP3^2}{PAC^2} \right) - A \quad dBm$$

• Receive Power Level (RPL) $RPL = 10 * \log_{10} \left(\frac{MGC * 2^{17}}{PAC^2} \right) - A \quad dBm$



We **devised** and **present** a technique that remote attackers can apply to circumvent this obstacle

	Status: 0		
	Txs Timestamp:	16.88284033	95276
	Rxs Timestamp:	10.69435802	52153
	First Path Amp	1: 5064	
	First Path Amp	2: 14560	
	First Path Amp	3: 22328	
	Maximum Growth	CIR: 14423	
	Rx Pream Count:	: 245	
	Extra Data Type	e: 0	
	Extra Data Leng	yth: 0	
Ch	ecksum: 0x0d9c	-	

Power level information in Avalue RTLS packets



Obtaining the anchor coordinates is the most challenging requirement. They are manually input at the first installation and never transmitted through the network



Attacker with Remote Access – Cont'd

- It is not possible to directly estimate the absolute distance, due to evolving temporary conditions
- However, if in a given moment to the power level information is identical, the tag j0 that triggered those packets must be positioned about exactly at the same distance from all anchors

Delta(i0, j0, t0) = (GT(reference, j0, t0) - GT(i0, j0, t0)) * c = 0

• Considering that CS(reference, t0) = 1 and ToF(reference) = 0, we can exploit this information and estimate the time of flights, thus the distances of the other anchors from the reference

ToF(i0) = CS(i0,t0) * (pTs(i0,j0,t0) - sTS(i0,t0)) - pTs(reference,j0,t0) + sTS(reference,t0)



Obtaining the anchor coordinates is the most challenging requirement. They are manually input at the first installation and never transmitted through the network



Attacker with Remote Access – Cont'd

• Finally, to obtain the coordinates, we can leverage the following installation constraint:





- Given that the anchor map is most times a rectangle, by arbitrarily setting the reference anchor in position (0;0), the coordinates of all other anchors can be easily estimated (they are given by the two shortest distances)
- An attacker can adapt the expected shape on the basis of the number of anchors detected in the communications



Obtaining the anchor coordinates is the most challenging requirement. They are manually input at the first installation and never transmitted through the network



Attacker with Remote Access – Cont'd

- This was actually tested in the Avalue RTLS, using both the First Path Power Level (FPPL) as well as the Receive Power Level (RPL)
- The best results are obtained using FPPL with a threshold of ~1% between the lowest power level and the highest



It was possible to estimate the anchors coordinates with an error of **less than 10%** with respect to the real values

• This can be accurate enough for attack scenarios where cm-level precision is not required

Anchors Coordinates Average Error wrt First Path Power Level (FPPL) Acceptance Threshold





Adversary Tactics, Techniques, and Procedures (TTPs)



Traffic Interception

To perform any meaningful attacks against RTLSs, an attacker first needs to intercept all network packets

Intercepting traffic requires two steps:

- 1. gaining a foothold inside the anchors-server backhaul network
- 2. executing a Man in the Middle (MitM) attack



Network Access

- Both Sewio and Avalue RTLS allow either Ethernet or Wi-Fi to be used for the network backhaul
- Gaining access to an Ethernet network requires that an attacker:
 - either compromises a computer in that network
 - or surreptitiously adds a rogue device
- The complexity of these attacks varies on the basis of the RTLS deployment configuration



Deployment configurations available on Sewio RTLS



Traffic Interception – Cont'd

To perform any meaningful attacks against RTLSs, an attacker first needs to intercept all network packets



Network Access – Cont'd

- As for Wi-Fi, both solutions support WPA2-PSK
- Gaining access to a Wi-Fi network requires:
 - either the knowledge of the WPA2 password
 - or the exploitation (if any) of vulnerabilities in the wireless appliances
- As for the first point, out of the box, both solutions use a static password that can be found in the public documentation
- In case an asset owner does not change it, obtaining access to the backhaul network is simple



Default WPA2-PSK password on Avalue RTLS



Traffic Interception – Cont'd

To perform any meaningful attacks against RTLSs, an attacker first needs to intercept all network packets



Man in the Middle

 In the tests executed, it was possible to MitM both solutions via standard ARP spoofing attacks



MitM attack against Sewio RTLS

arpspoof - i attacker_eth - t server_ip anchor1_ip & arpspoof - i attacker_eth - t anchor1_ip server_ip



The attacks were completely **undetected** by the RTLS. No warnings or abnormal behavior that may alert an operator were shown.



Passive Eavesdropping Attacks

After obtaining access to an RTLS network and launching the MitM attack, an attacker can reconstruct the position of tags by executing one of the TDoA algorithms known in literature



Active Traffic Manipulation Attacks

To accomplish an active attack, an attacker first needs to do a target reconnaissance and add traffic filtering routines to the attack algorithm

Target Reconnaissance

• To successfully deceive an operator, it is important that the tag movements appear natural

If the target is a human being, faking its position with harsh, sudden movements would **warn an operator** and make them think that, at the very least, a malfunctioning is occurring

 This phase can be accomplished by simply performing a passive eavesdropping attack against the target

Active Traffic Filtering

- If the packet is a synchronization packet, it must be automatically forwarded to the destination
- If the packet is a positioning packet of a target, its timestamp must be modified (and the checksum updated). If not a target one, it must be forwarded unaltered
- Many techniques are available. Notably, we leveraged iptables NFQUEUE, a flexible userspace packet handler



Active Traffic Manipulation Attacks – Cont'd

Finally, an attacker can alter the timestamps by simply inverting all the equations previously described



Packet Information Manipulation

- In a manipulation attack, the tag coordinates are known (they are the target coordinates that an attacker wants to fake for a given tag) and the positioning timestamps are unknown
- First, the attacker derives the modified positioning timestamps according to the target coordinates
- Finally, the attacker re-computes the packet checksums and then sends the modified packets



Attack Demos 🕥



Locating and Targeting People/Assets





Geofencing





Contact Tracing

Sensmap Building:office office office office fractional sewiorus office	ilse_contact.py A FORWARD -p UDPsport 5000 -j NFQUEUEqueue-num 1 x2204612F3A84 and 0x22045FFD8CDE
ENCOUNTER 0x22045FFD8CDE and 0x2204612F3A84	x2204612F3A84 and 0x22045FFD8CDE
((Q)) Q D0-13 FB-15 9F-14	
0 00:31:301:4:79:10 10:109:73:32:17a1cc 000 0:022045FFD0CCF_M012F30A94 00:31:501:4:79:10 00:37:501:4:79:10	00 42: arp reply 192.168.225.13 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.15 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.14 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.12 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.1 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.1 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.1 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.1 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 42: arp reply 192.168.225.2 is-at 80:3f:5d:4:79:16 00 4



Remediations



Segregation and Firewall Rules



Goal

• Move the entire UWB RTLS backhaul network to a segregated network, and secure the access to the network both physically and logically

This is now mandated by some RTLS vendors

Physical access

- Restrict physical access to the device to qualified personnel
- Disable unused physical interfaces of the device. Unused interfaces could be used to gain access to the operating site.

Software - Safety functions

- Only use protocols that are required to operate the device.
- Restrict access to the device with a firewall or rules in an ACL (Access Control List).
- Using VLANs gives you good protection against DoS attacks. Check whether this is
 practicable.
- Activate the access logging function (external). Use the central logging function to record changes and access.
- Configure a SysLog server to save all logs to a central location.

SIMATIC RTLS4030G Operating Instructions, 04/2021, C79000-G8976-C515-06

Siemens RTLS4030G operating instructions



Advantages

- Allows the problem to be mitigated relatively quickly
- Can be enacted by deploying traditional solutions such as VLANs, IEEE 802.1X, firewall rules



Challenges

- Some RTLS servers expose core network services on all interfaces. Firewall rules must be set to allow as few services as possible on the management interface
- Does not protect from a physical MitM (either via wire tap, or wireless sniffer if wireless password is compromised)



Intrusion Detection Systems



Goal

 Detect signs of MitM attacks. Leverages the fact that MitM attacks are unavoidable to obtain the timestamps

This option was successfully tested on both Sewio and Avalue RTLS



Detection of a MitM attack by an IDS



Advantages

- Plug-and-play solution
- Allows the problem to be mitigated very quickly



Challenges

 Does not protect from a physical MitM (either via wire tap, or wireless sniffer if wireless password is compromised)



Traffic Encryption



Goal

 Add a traffic encryption layer on top of the existing communications, to protect even against a physical MitM

This option was successfully tested on the Avalue RTLS for a PoC using standard tools (SSH tunnel and Socat)



SSH tunnel PoC on Avalue RTLS



Advantages

- The closest mitigation to completely solving the problem
- Allows basic RTLS functionalities to remain unaltered



Challenges

- In Avalue RTLS, it was necessary to reduce the number of syncs per second to counteract the higher load, at the expense of a reduced accuracy
- Entirely depends on the accessibility of the RTLS server and anchors from the vendor



Summary & Key Takeaways





Wireless technology increases efficiency + productivity while reducing unnecessary cabling infrastructure costs

- IEEE 802.15.4z has out of scope areas, creating security loopholes
- Nozomi Networks Labs discovered zero-days in two popular UWB RTLS
- UWB RTLS is used for personnel tracking, geofencing, and contact tracing
- Threat actor TTPs are MitM and eavesdropping
 or manipulation tactics
- Mitigations include segregation and firewall rules, IDS, and traffic encryption



Black Hat Sound Bytes

Key Takeaways



Weak security requirements in critical software can lead to **safety issues** that cannot be ignored



There are attack surfaces out there that no one is looking at, but **they have significant consequences if compromised**



Exploiting secondary communications in UWB RTLS can be **challenging**, **but it is doable**





Thank You!

Questions?

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