blackhat® USA 2022

Devils Are in the File Descriptors: It is Time To Catch Them All

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#BHUSA @BlackHatEvents

About me

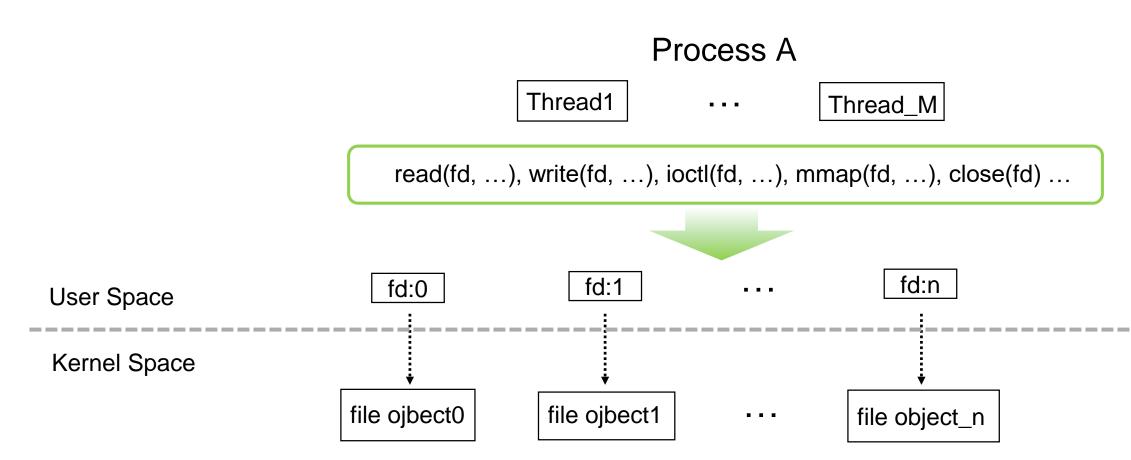
Le Wu(<u>@NVamous</u>)

- Focus on Android/Linux bug hunting and exploit
- Found 200+ vulnerabilities in the last two years
- Blackhat Asia 2022 speaker

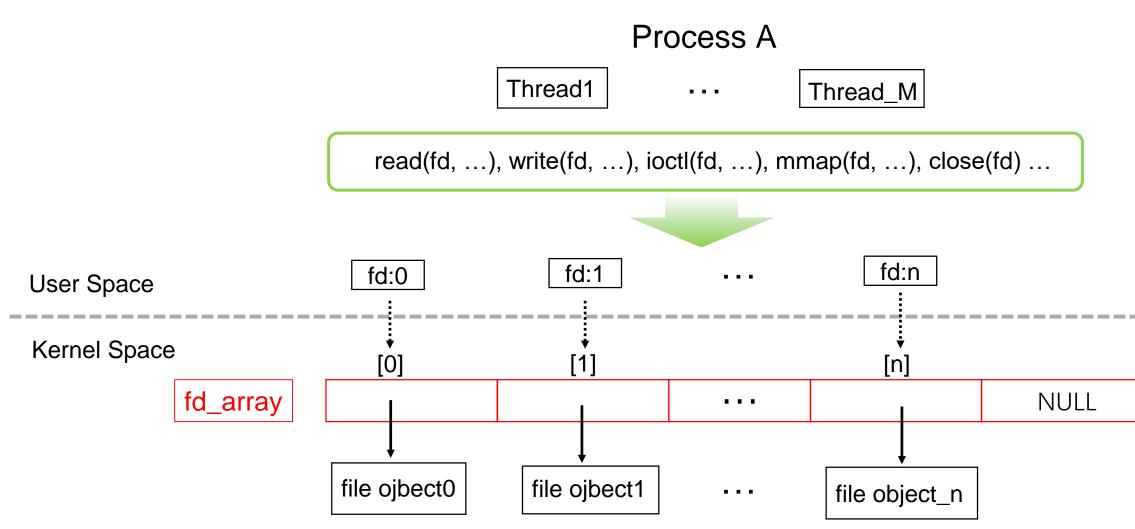


- Background
- Diving into issues in the fd **export** operations
- Diving into issues in the fd **import** operations
- Conclusion & Future work

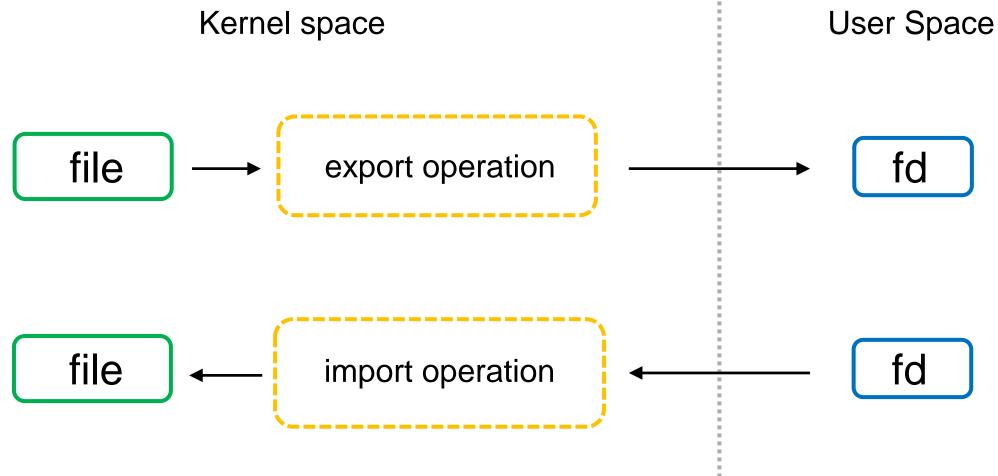
Introduction to file descriptor—— An integer in a process



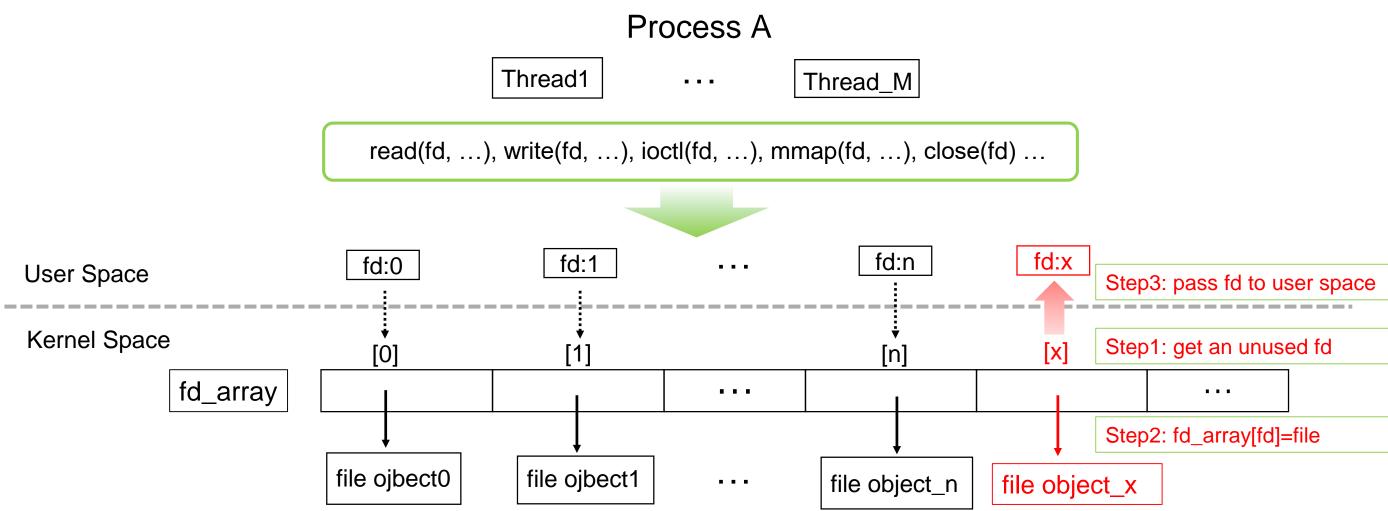
Introduction to file descriptor—— An integer in a process



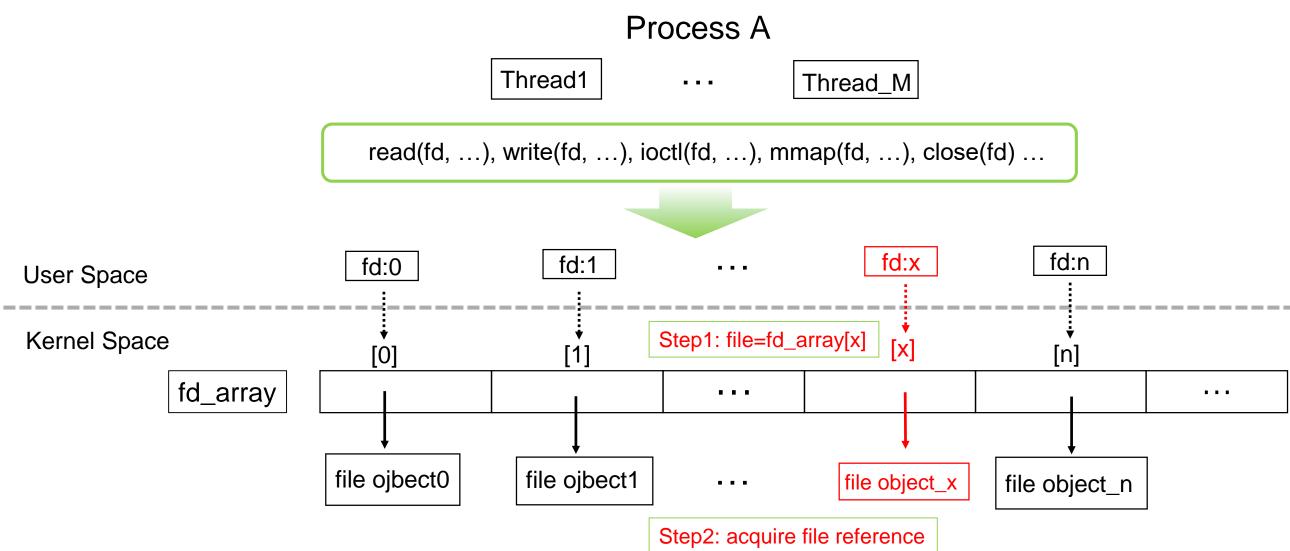
Introduction to file descriptor—fd export operation and import operation in kernel



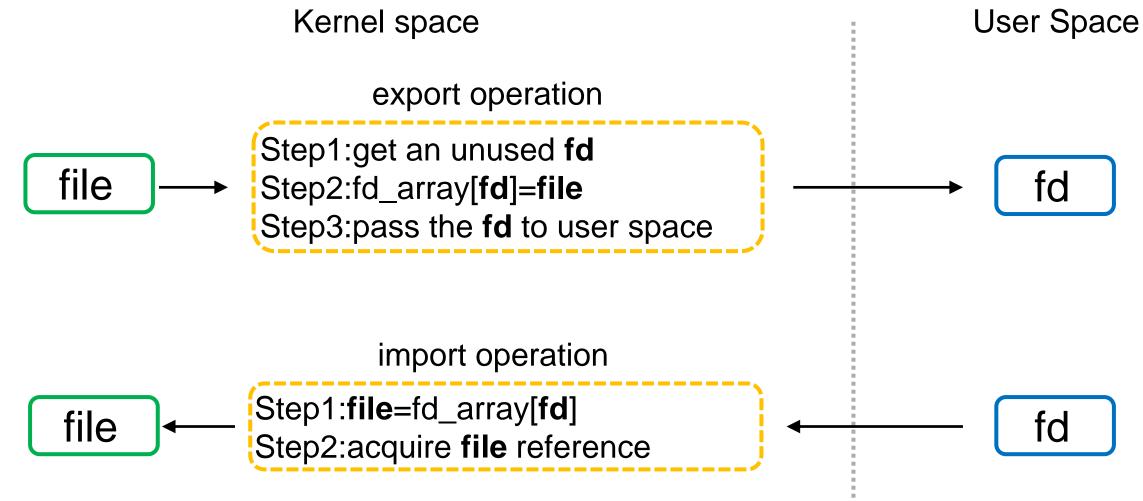
Introduction to file descriptor—— fd export operation in kernel



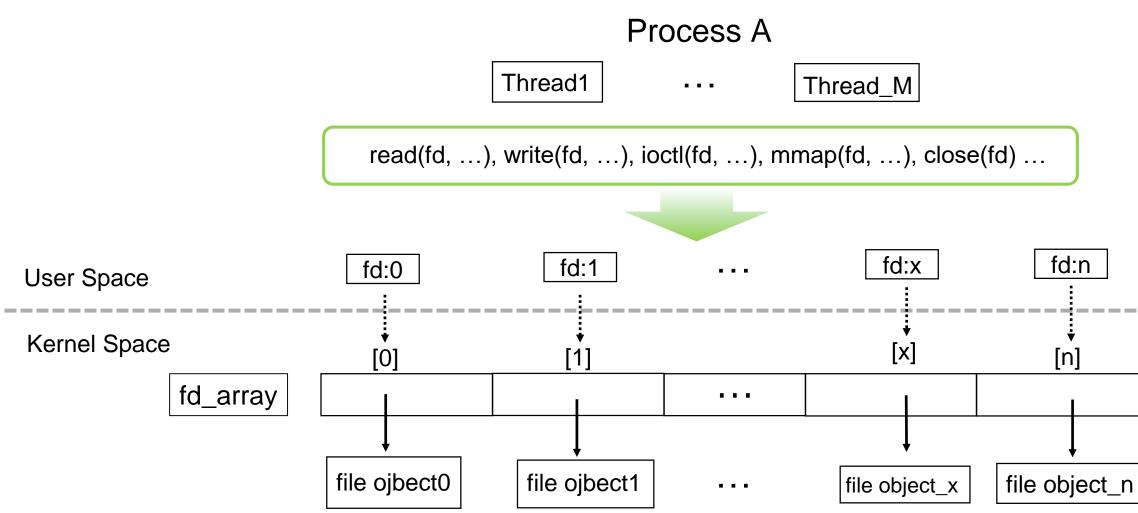
Introduction to file descriptor—— fd import operation in kernel



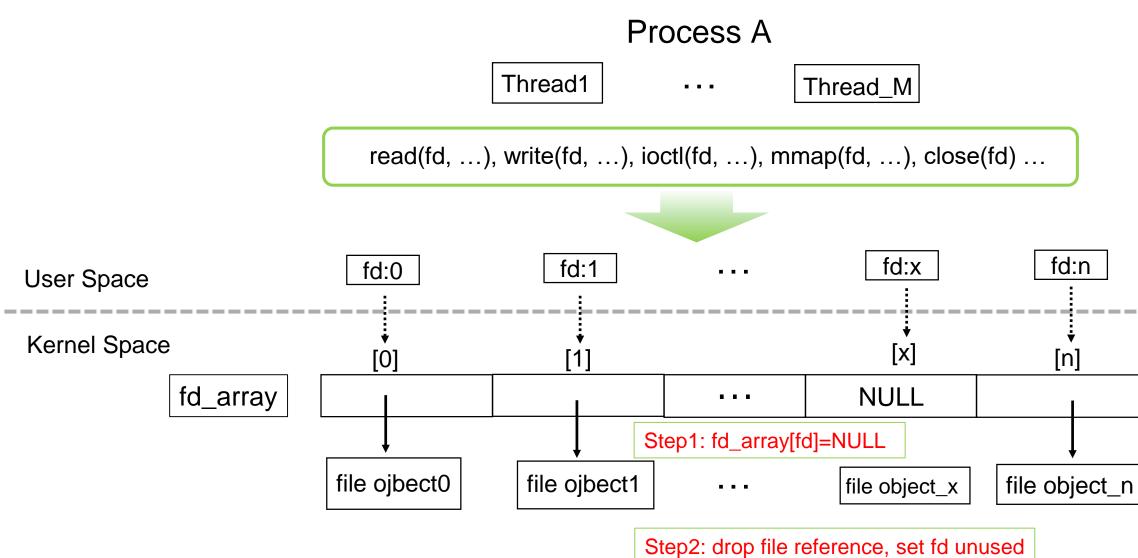
Introduction to file descriptor—fd export operation and fd import operation



Introduction to file descriptor—— User process close(fd)

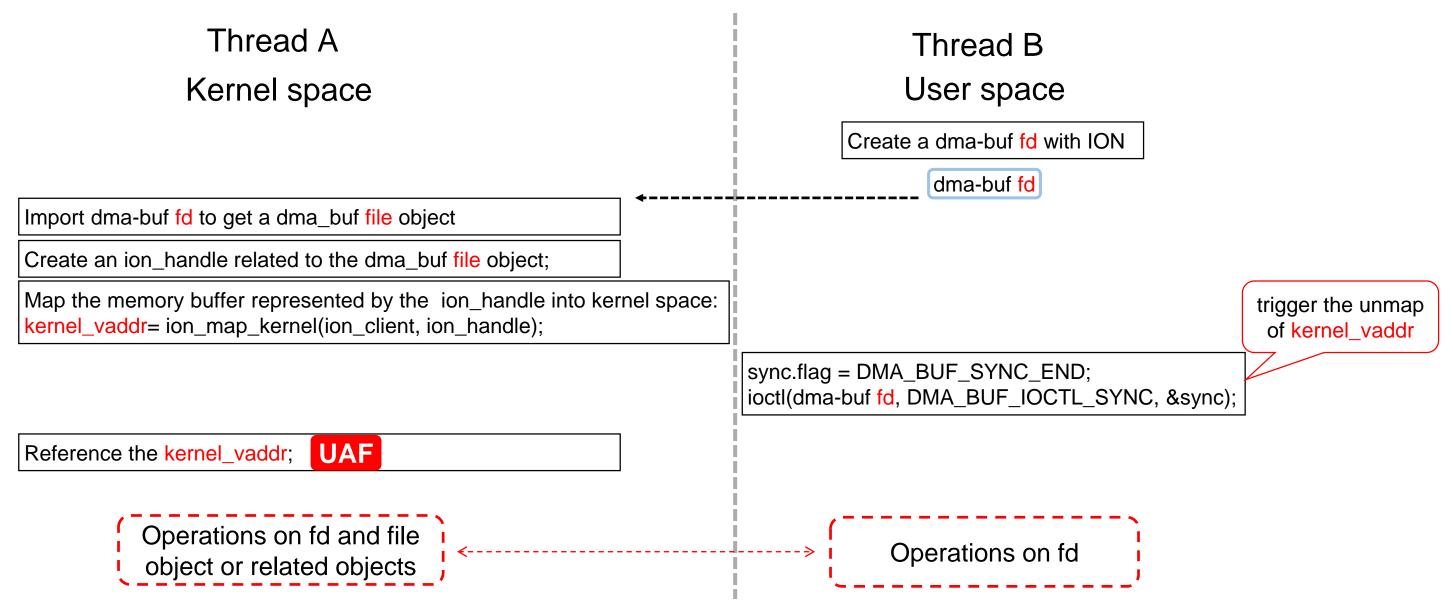


Introduction to file descriptor—— User process close(fd)



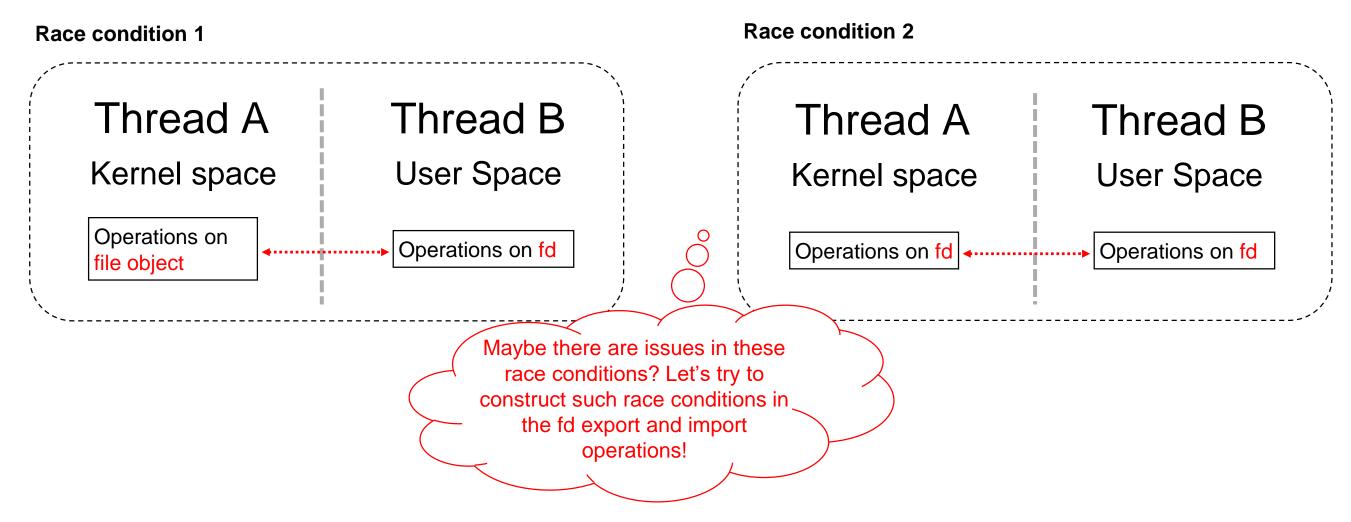
•••

Why file descriptor—Inspired by CVE-2021-0929

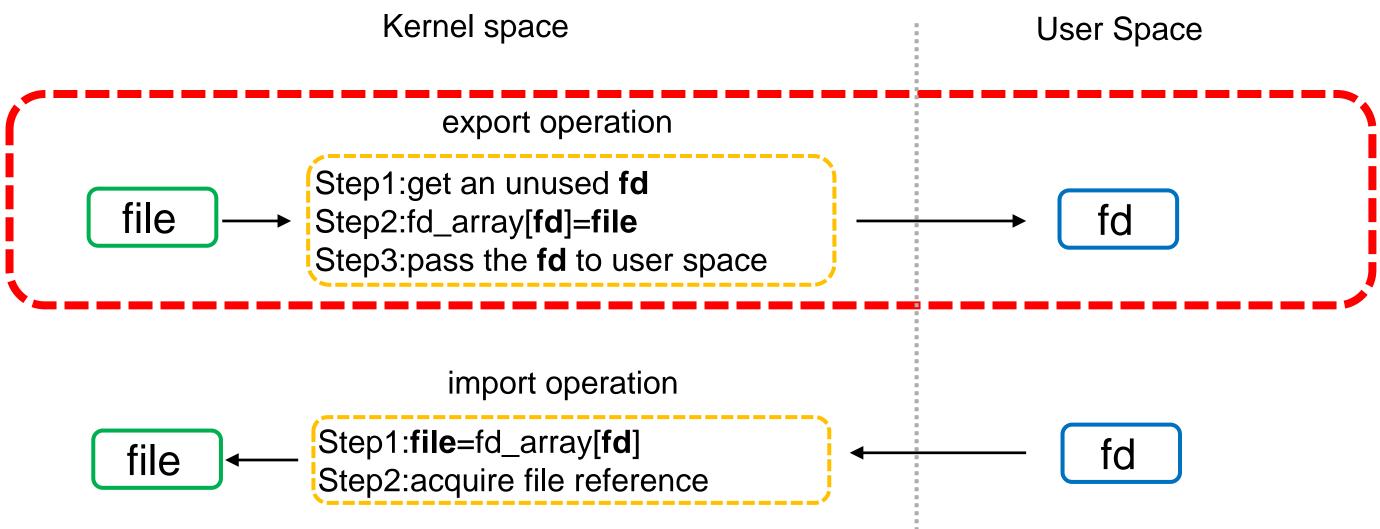


Why file descriptor—Inspired by CVE-2021-0929

A file descriptor can be shared between kernel space and user space, race condition can happen between kernel and user operations:



Diving into issues in the fd export operation



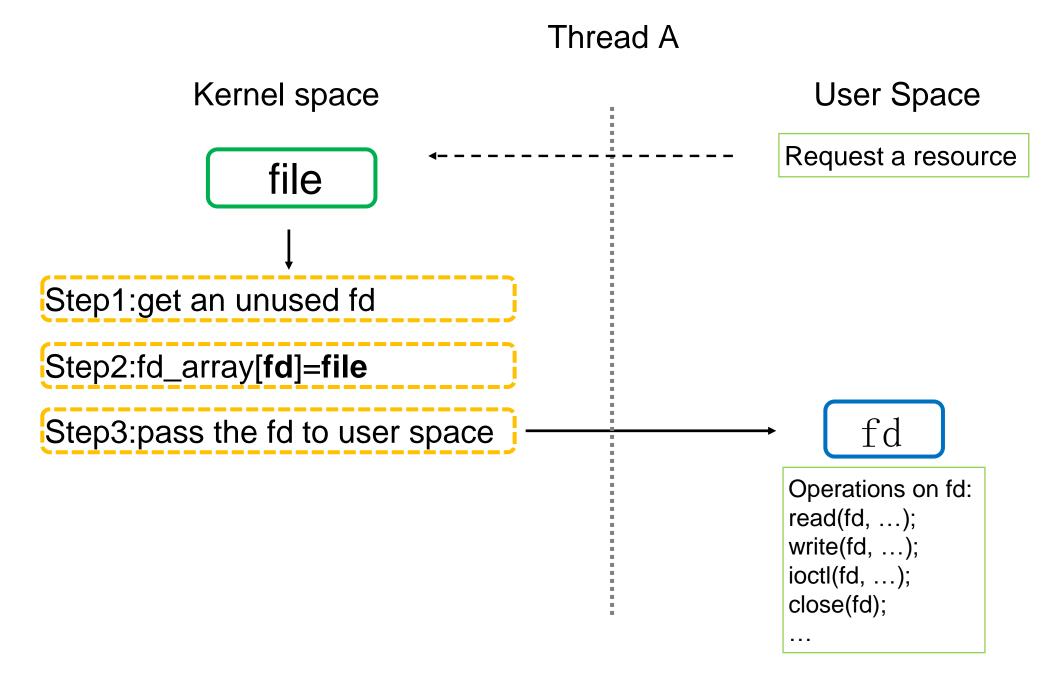


Diving into issues in the fd export operation

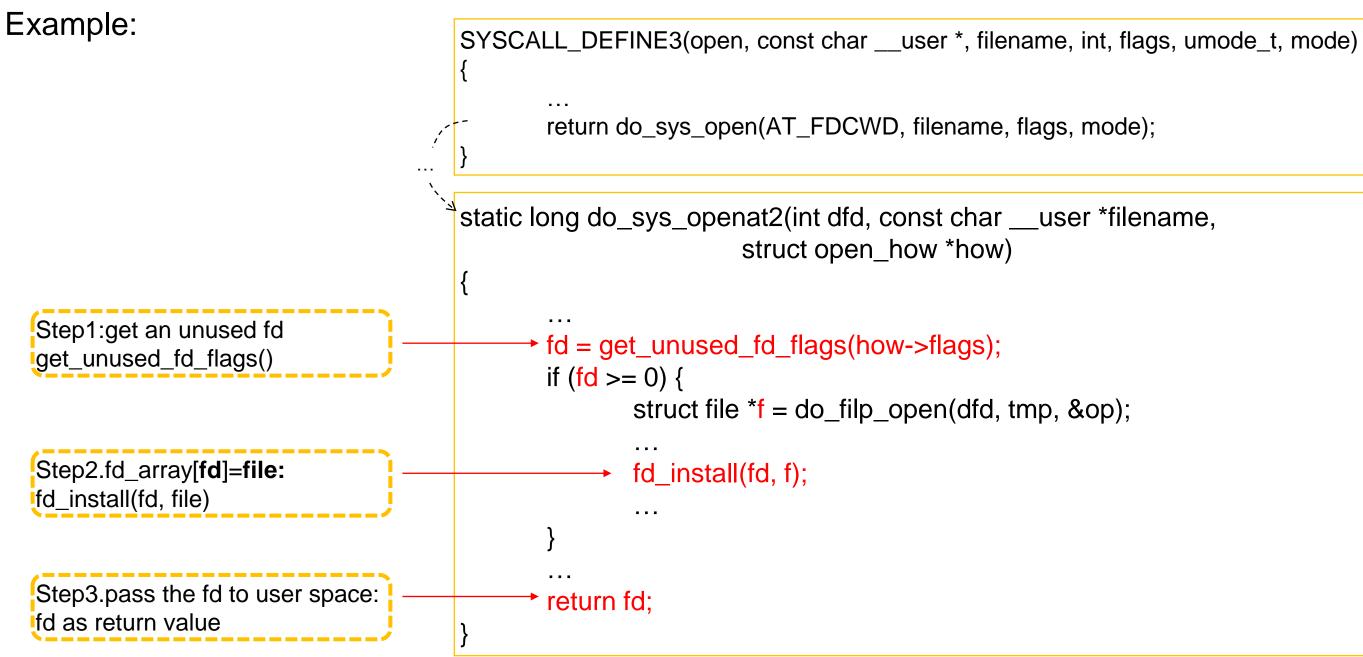
- Scenario of fd export operation
- UAF caused by race condition
- Find the issues
- Fixes



Scenario of fd export operation

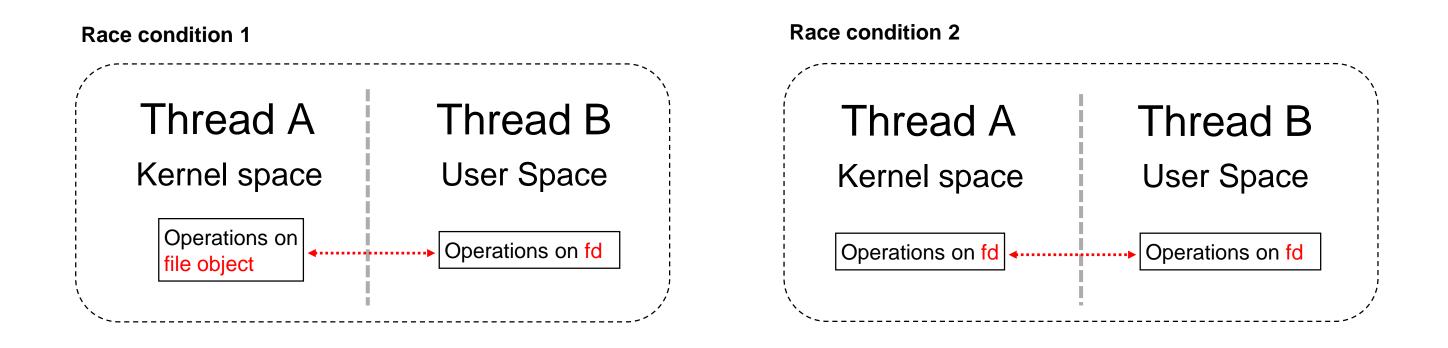


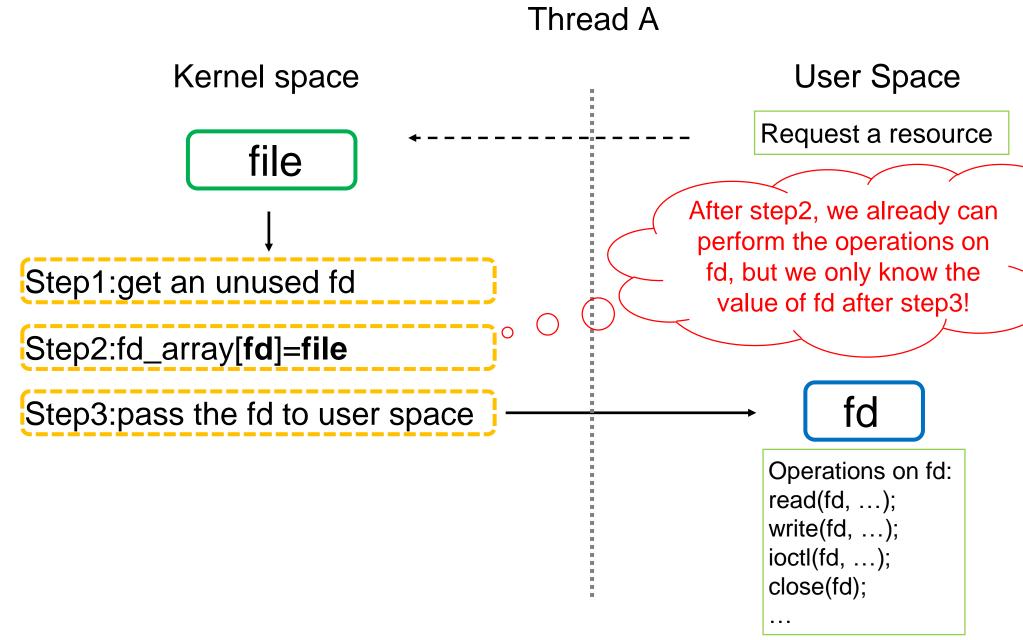
Scenario of fd export operation



Scenario of fd export operation

But this regular fd export operation is executed sequentially, which is still far from the race conditions we want to see:







Hold on! Do we have to wait for fd to be passed from kernel to know the value of it ?

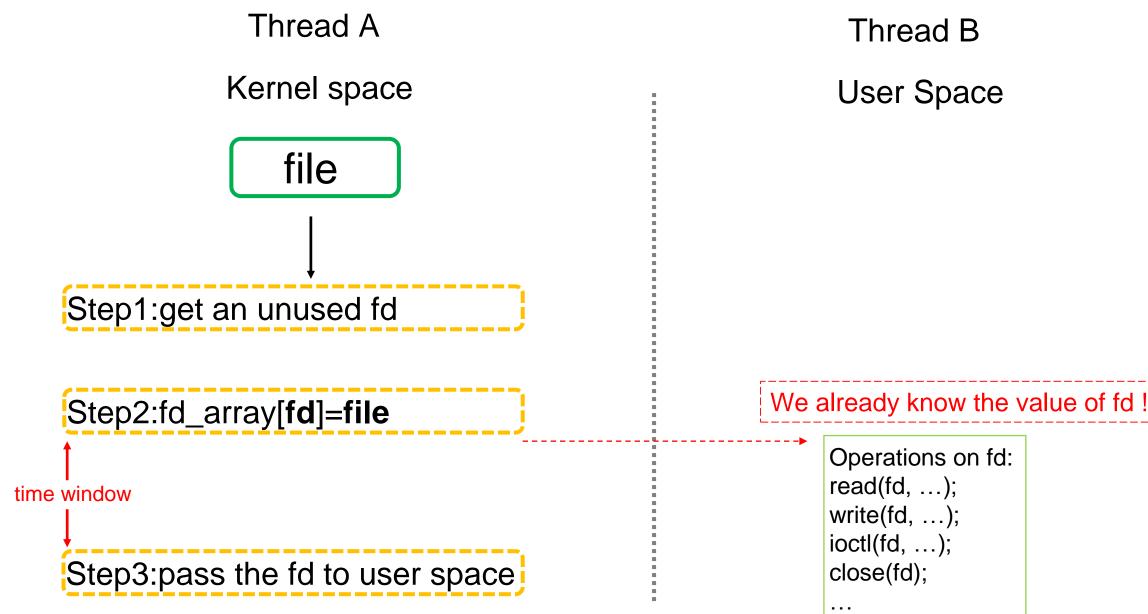
Fd is predictable:

• Assigned in ascending order

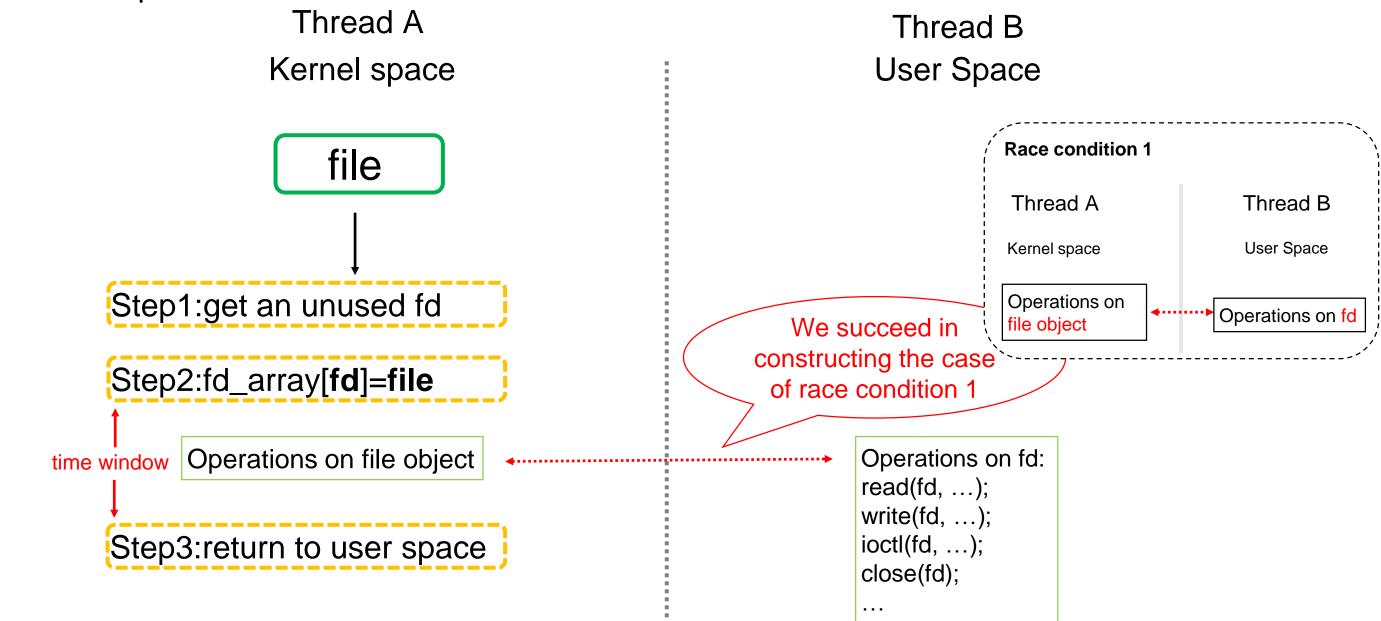
For a new process, fd 0, 1, 2 are usually occupied, 3 will be the next fd exported from kernel, and then 4, 5, 6.....

Reused after close(fd)

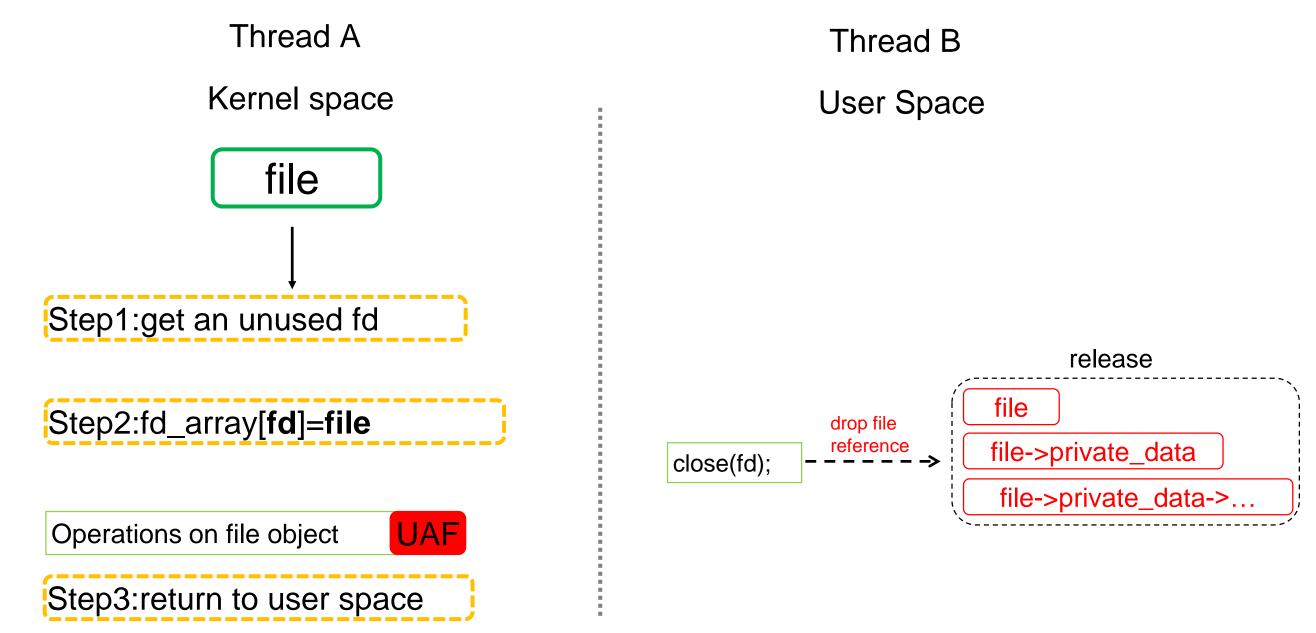
int fd = open(file_path, ...); close(fd); fd2=fd int fd2 = open(file_path2,...);



More assumption:



A potential UAF scenario:



Looking for all kinds of kernel APIs which perform the "step2":



- fd_install(fd, file)
- anon_inode_getfd()
- dma_buf_fd()

. . .

- sync_fence_install()
- ion_share_dma_buf_fd()

They all wrap fd_install(fd, file)

Try to search for the bug pattern: "reference file or related objects after the step2"

From Vendor Q:

```
static int get_fd(uint32_t handle, int64_t *fd)
          int unused fd = -1, ret = -1;
          struct file f = NULL;
          struct context *cxt = NULL;
           . . .
          cxt = kzalloc(sizeof(*cxt), GFP_KERNEL);
           . . .
          unused_fd = get_unused_fd_flags(O_RDWR);
          f = anon_inode_getfile(INVOKE_FILE, &invoke_fops, cxt, O_RDWR);
           fd = unused fd:
          fd_install(*fd, f);
           ((struct context *)(f->private_data))->handle = handle;
          return 0;
. . .
```

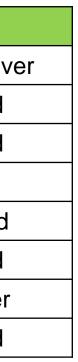
From Vendor M:

int ged_ge_alloc(int region_num, uint32_t *region_sizes) unsigned long flags; int i: struct GEEntry *entry = (struct GEEntry *)kmem_cache_zalloc(gPoolCache, ...); entry->alloc_fd = get_unused_fd_flags(O_CLOEXEC); entry->file = anon inode getfile("gralloc extra", &GEEntry_fops, entry, 0); fd install(entry->alloc fd, entry->file); return entry->alloc_fd; . . .

My assumption is correct! let's try to search for more!

I found since the end of 2021:

	From	CVE-id/issue	fd exported by function	Feature
	Vendor M	CVE-2022-21771	fd_install()	GPU related drive
		CVE-2022-21773	dma_buf_fd()	dma-buf related
		Duplicated issue#1	dma_buf_fd()	dma-buf related
	Vendor Q	CVE-2022-33225	fd_install()	
	Vendor S	Issue#1	fd_install()	sync_file related
		Issue#2	dma_buf_fd()	dma-buf related
	Linux Mainstream	Issue#1	anon_inode_getfd()	Amd GPU driver
Maybe I shou	ld	Issue#2	dma_buf_fd()	dma-buf related
pay more attention to th GPU drivers	ne			
	ARM Mali GPU driver	CVE-2022-28349	anon_inode_getfd()	can be triggered fr untrusted apps
		CVE-2022-28350	fd_install()	sync_file related, of be triggered from untrusted apps
		•	· · · · · · · · · · · · · · · · · · ·	



from , can n

CVE-2022-28349—— A Nday in ARM Mali GPU driver

```
int kbase_vinstr_hwcnt_reader_setup(
         struct kbase_vinstr_context *vctx,
         struct kbase ioctl hwcnt reader setup *setup)
         int errcode;
         int fd;
         struct kbase vinstr client *vcli = NULL;
         errcode = kbasep_vinstr_client_create(vctx, setup, &vcli);
         . . .
         errcode = anon inode getfd(
                  "[mali_vinstr_desc]",
                  &vinstr client fops,
                  vcli,
                  O RDONLY | O CLOEXEC);
         . . .
         fd = errcode;
         . . .
         list_add(&vcli->node, &vctx->clients);
         . . .
```

Affect:

Midgard GPU Kernel Driver: All versions from r28p0 – r29p0
Bifrost GPU Kernel Driver: All versions from r17p0 – r23p0
Valhall GPU Kernel Driver: All versions from r19p0 – r23p0

Android 10 devices of some vendors are affected !

CVE-2022-28350—— A Oday in ARM Mali GPU driver

```
static int kbase_kcpu_fence_signal_prepare(...)
        struct sync_file *sync_file;
        int ret = 0:
         int fd;
                                                                             Affect:
         sync_file = sync_file_create(fence_out);
                                                                             Valhall GPU Kernel Driver: All versions from r29p0 – r36p0
         fd = get_unused_fd_flags(O_CLOEXEC);
         fd_install(fd, sync_file->file);
                                                                                                Android 12 devices of some
         if (copy_to_user(u64_to_user_ptr(fence_info->fence), &fence,
                                                                                                    vendors are affected !
                           sizeof(fence))) {
                 ret = -EFAULT;
                 goto fd_flags_fail;
        return 0;
fd_flags_fail:
         fput(sync_file->file);
         . . .
         return ret;
```

Exploit of CVE-2022-28350

A known exploit <u>method</u>

Given by Mathias Krause from grsecurity for a similar vulnerability <u>CVE-2022-22942</u>:

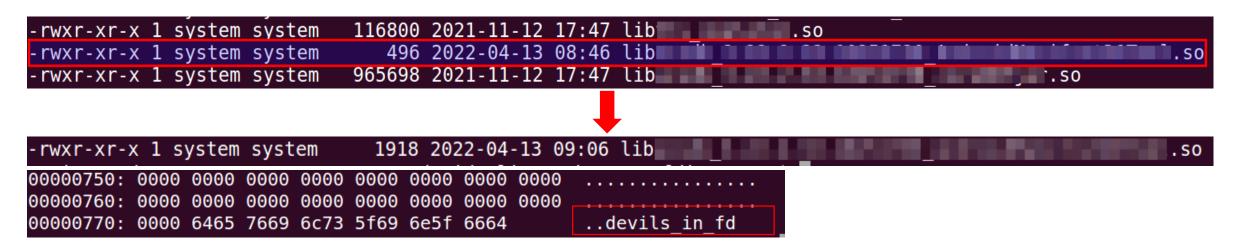
- No need for KASLR、SMEP/SMAP、 KCFI bypass ٠
- Read/write privileged files from unprivileged processes

The method won't work on Android because of SELinux 🛞

My new exploit method

- Bypass SELinux and work on the affected Android 12 devices •
- Write privileged files from untrusted apps ٠

(Details are put in the supplement part of the slides)



Find the issues

Check if the file or related objects are referenced after these functions:

- fd_install(fd, file)
- anon_inode_getfd()
- dma_buf_fd()

•

. . .

- sync_fence_install()
- ion_share_dma_buf_fd()

They all wrap fd_install(fd, file)

Fixes

Don't reference the file or related objects after step2 of fd export operation in kernel until return to user space •

```
static long do_sys_openat2(int dfd, const char __user *filename,
                   struct open_how *how)
     struct open_flags op;
     int fd = build_open_flags(how, &op);
      . . .
     fd = get_unused_fd_flags(how->flags);
     if (fd \ge 0) {
           struct file *f = do_filp_open(dfd, tmp, &op);
           if (IS_ERR(f)) {
           } else {
                 fsnotify_open(f);
                 fd_install(fd, f);
     putname(tmp);
                        return to user space directly
     return fd;
```

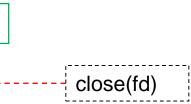
Fixes

Reference the file object or related objects with lock protection, and share the lock in file_release of fd:

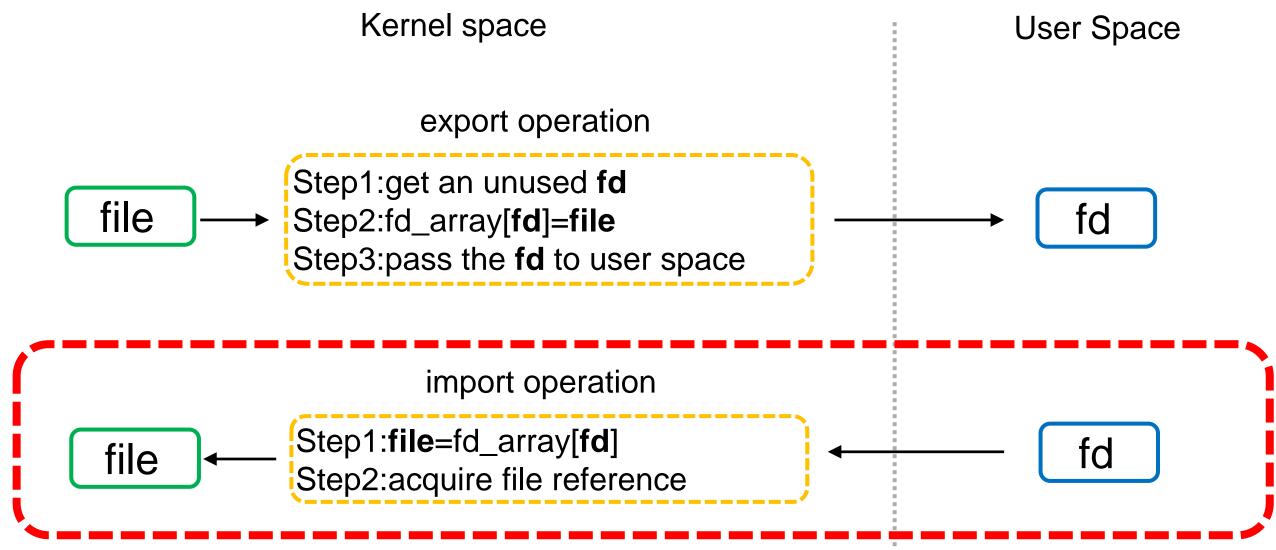
```
int file release(...) {
       int fd_export_func(...) {
                 mutex_lock(g_lock);
                                                                                               . . .
                 fd_install(file, fd);
                                                                                               mutex_lock(g_lock);
                 Reference file or related objects;
                                                                                               . . .
                 mutex unlock(g lock);
                                                                                               mutex unlock(g lock);
                 return fd;
                                                                                               . . .
     (From vendor S)
                                                                                            static int hpa_trace_task_release(struct inode *inode, struct file *file)
void hpa_trace_add_task(void)
         struct hpa_trace_task *task;
                                                                                                      struct hpa_trace_task *task = file->private_data;
                                                                                                       . . .
         mutex lock(&hpa trace lock);
                                                                                                      mutex lock(&hpa trace lock);
                                                                                                      list_del(&task->node);
          . . .
         task = kzalloc(sizeof(*task), GFP_KERNEL);
                                                                                                      mutex_unlock(&hpa_trace_lock);
                                                                                                      kfree(task);
         fd = get_unused_fd_flags(O_RDONLY | O_CLOEXEC);
                                                                                                      return 0;
         task->file = anon_inode_getfile(name, &hpa_trace_task_fops, task, O_RDWR);
```

```
fd_install(fd, task->file);
list_add_tail(&task->node, &hpa_task_list);
mutex unlock(&hpa trace lock);
```

. . .



Diving into issues in the fd import operation





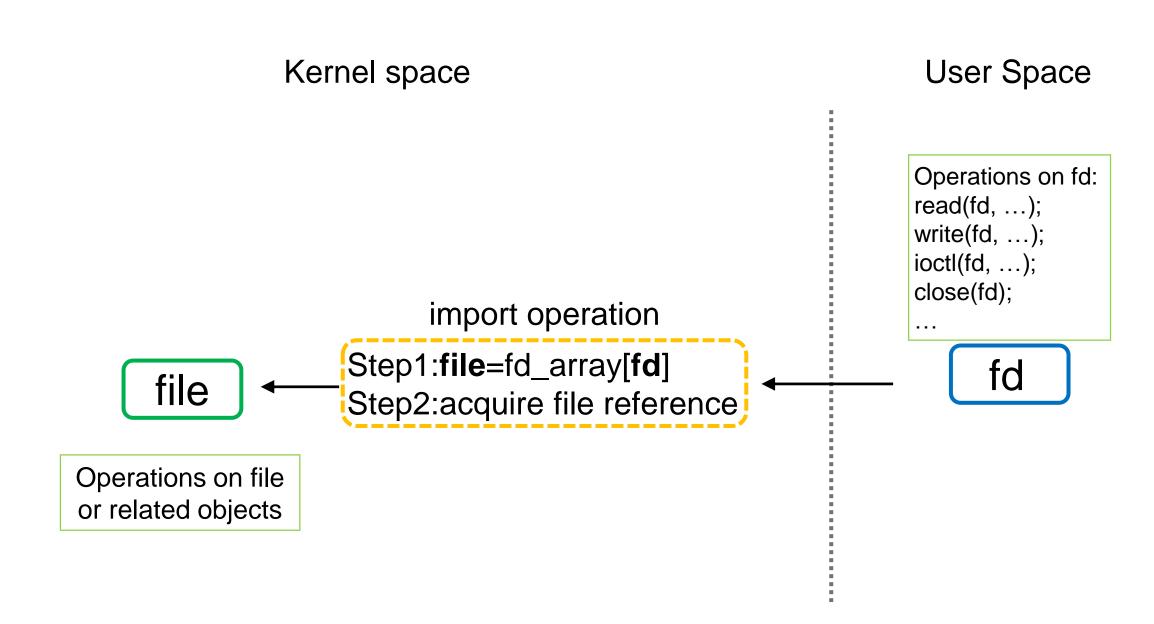
Diving into issues in the fd import operation

- Scenario of fd import operation
- Fd type confusion caused by race condition
- Find the issues
- Fixes

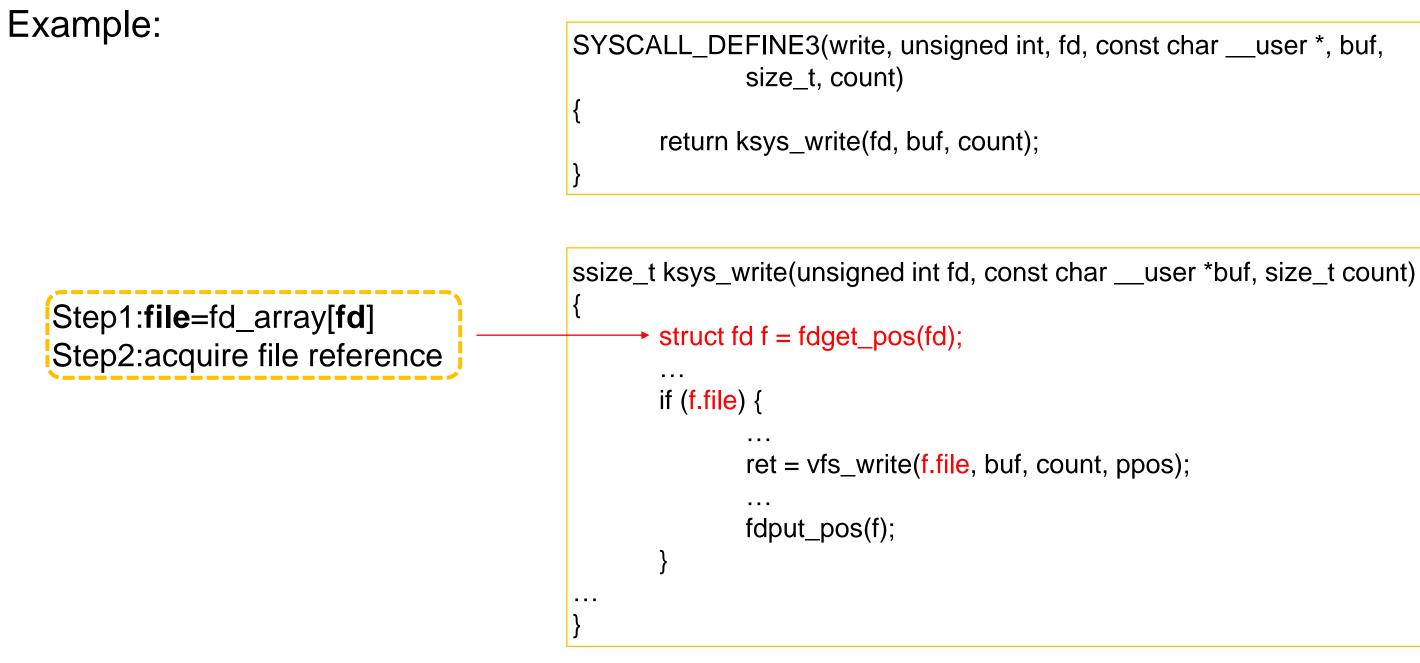


Scenario of fd import operation



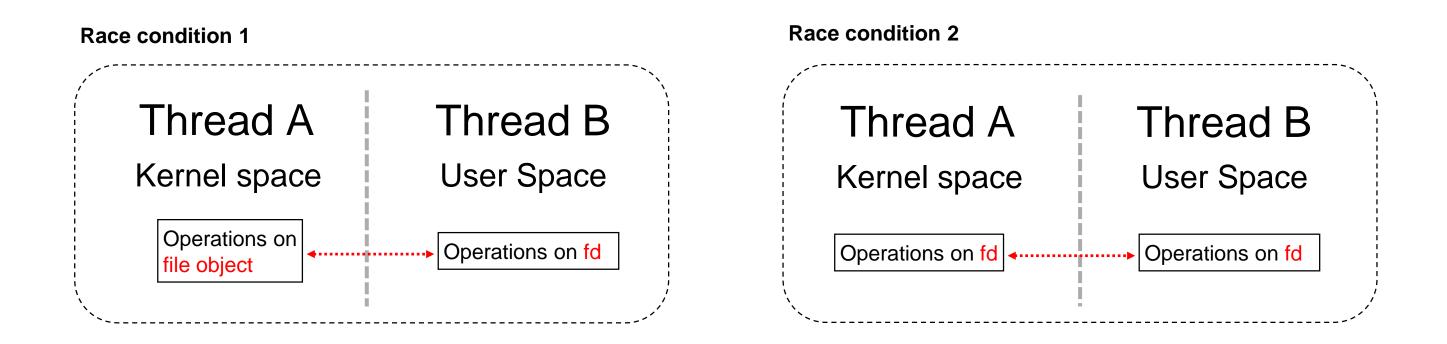


Scenario of fd import operation



Scenario of fd import operation

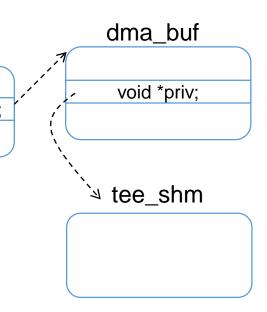
But this regular fd import operation is executed sequentially, which is still far from the race conditions we want to see:

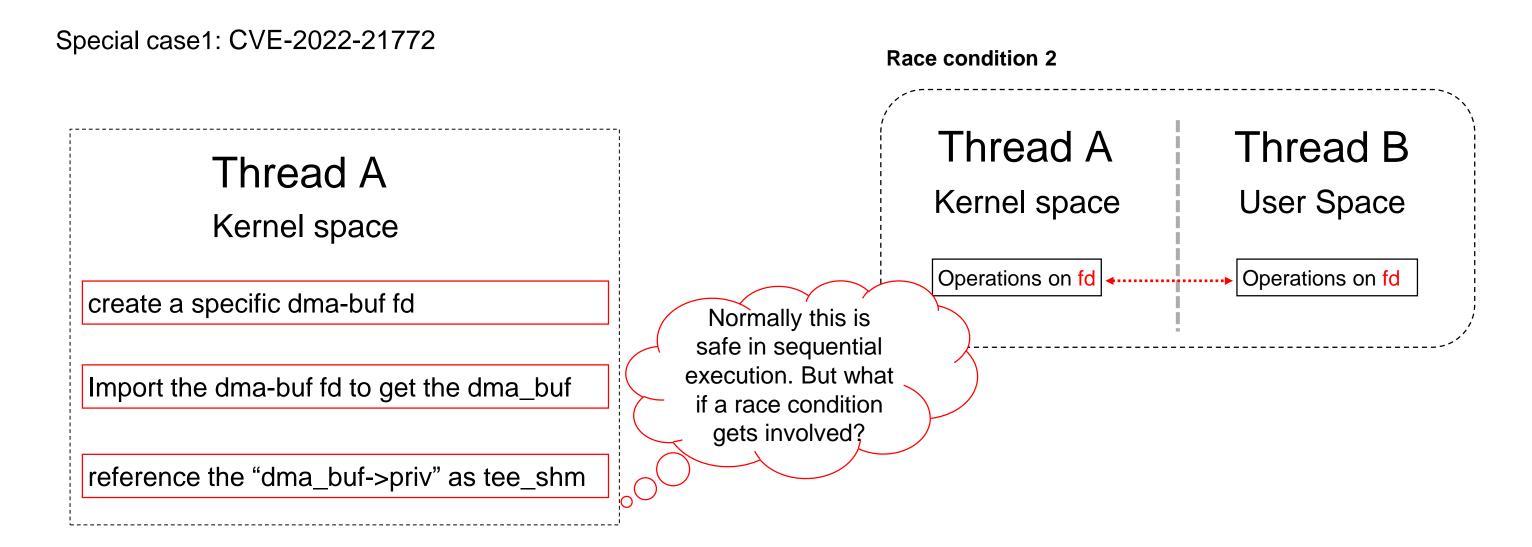


Searching for all kinds of scenarios of fd import operation in kernel...

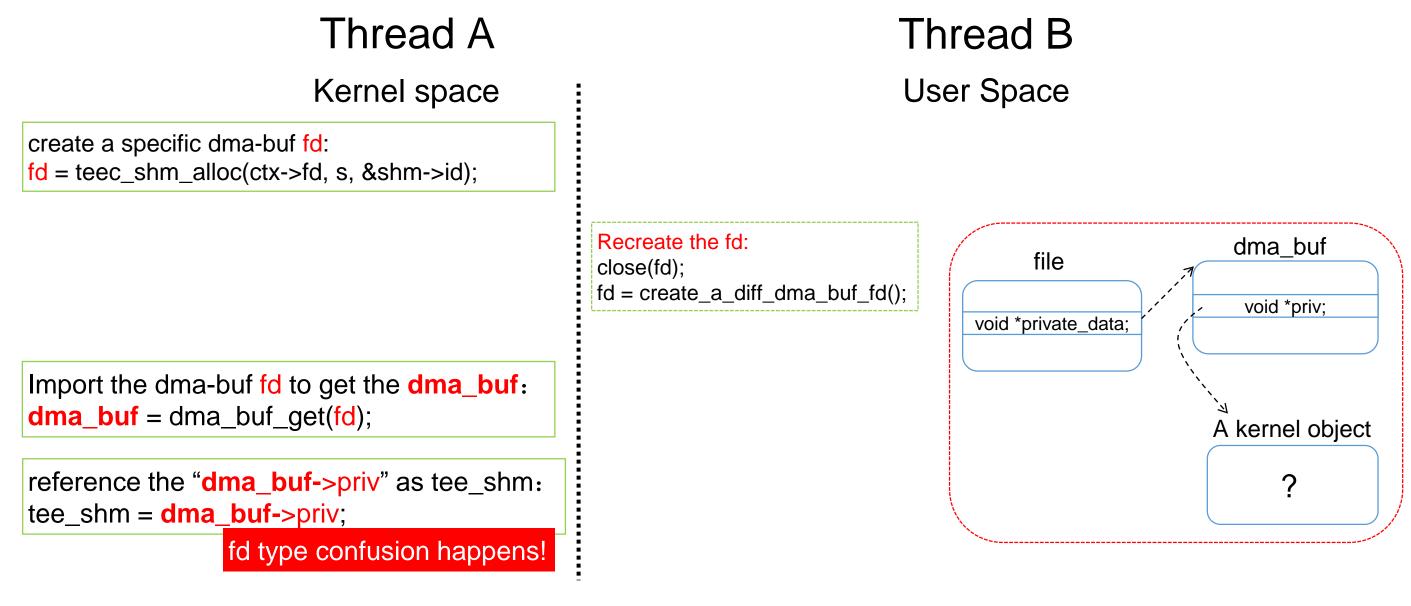
Special case1: CVE-2022-21772

```
TEEC_Result TEEC_RegisterSharedMemory(struct TEEC_Context *ctx,
                                            struct TEEC SharedMemory *shm)
{
        int fd;
        size ts;
        struct dma buf *dma buf;
        struct tee_shm *tee_shm;
                                                                                                           file
         . . .
        fd = teec_shm_alloc(ctx->fd, s, &shm->id);
                                                    create a specific dma-buf fd
         . . .
                                                                                                     void *private data;
                                                    import the dma-buf fd to get the dma_buf
        dma buf = dma buf get(fd);
        close(fd);
         . . .
                                                    reference the "dma_buf->priv" as tee_shm
        tee_shm = dma_buf->priv;
         . . .
        shm->shadow_buffer = tee_shm->kaddr;
         . . .
        return TEEC_SUCCESS;
```



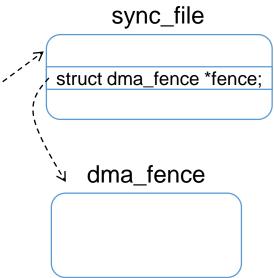


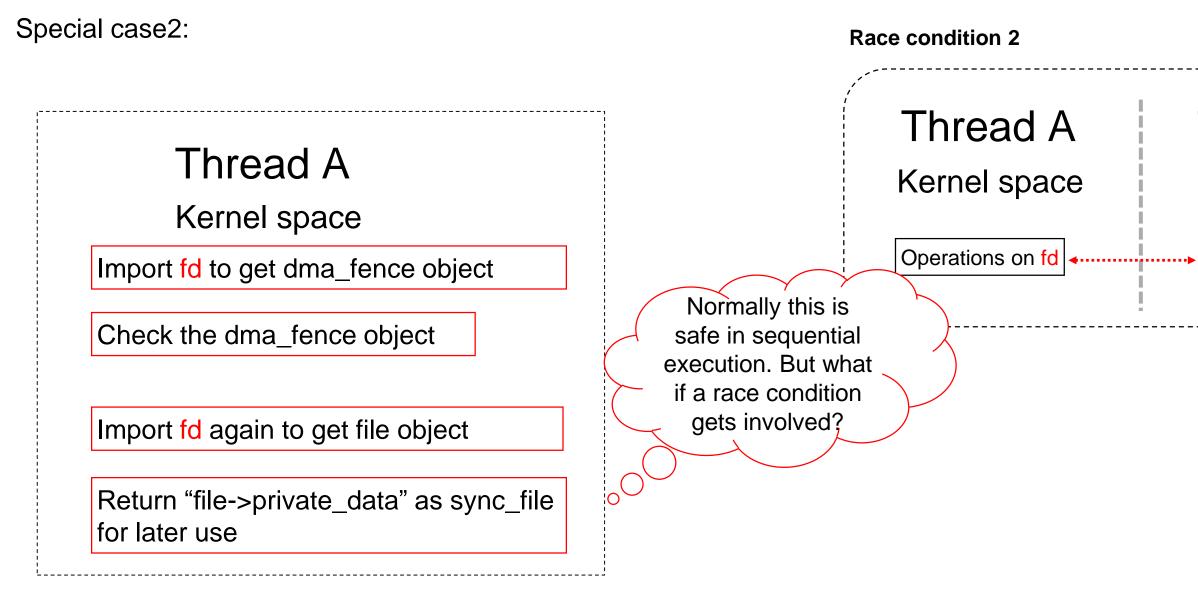
Special case1: CVE-2022-21772

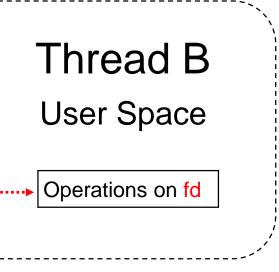


Special case2:

```
file
struct sync_file*internal_sync_fence_fdget(int fd)
                                                                                            void *private_data;
    struct file *file;
     struct dma_fence *fence = sync_file_get_fence(fd); Import fd to get dma_fence object
    /* Verify whether the fd is a valid sync file. */
    if (unlikely(!fence))
                                                                Check the dma_fence object
         return NULL;
     dma_fence_put(fence);
                                                                Import fd again to get file object
    file = fget(fd);
     return file->private_data;
                                                                Return "file->private_data" as sync_file for later use
```







Special case2:

Thread A Kernel space

Import fd to get dma_fence object:
struct dma_fence *fence = sync_file_get_fence(fd);

Check the dma_fence object

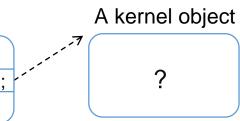
Import fd to get file object: file = fget(fd);

Return "file->private_data" as sync_file for later use

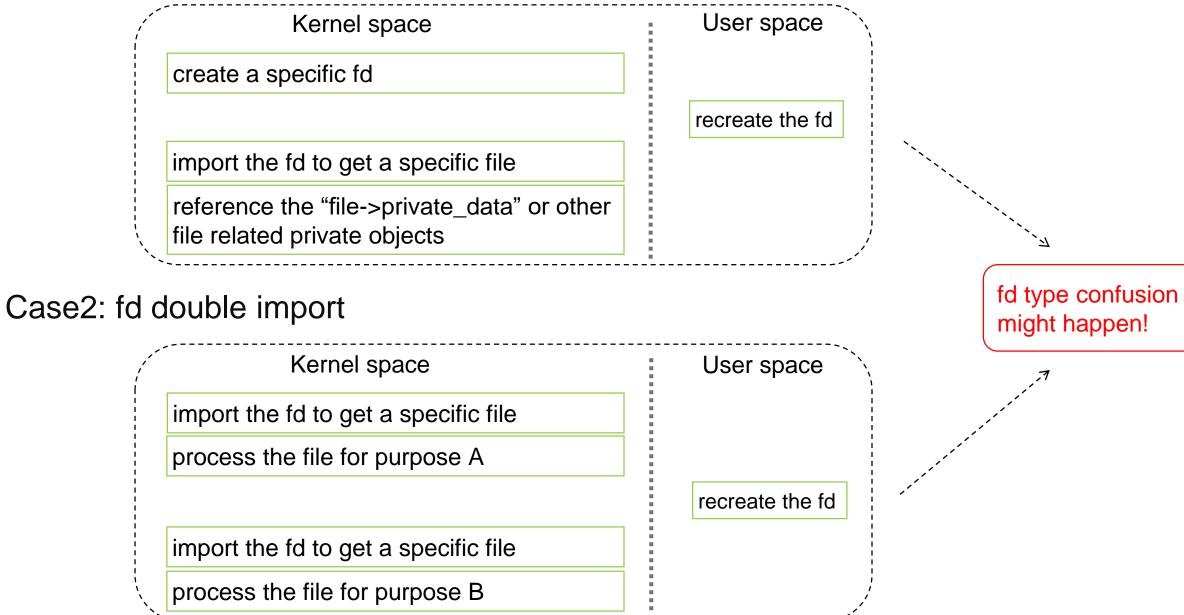
fd type confusion happens!

U	User Space		
	file		
Recreate the fd: close(fd); fd = open();	void *private_data;		

Thread B



Case1: fd time-of-create time-of-import •



Find the issues

There are still two questions that need to be answered:

- \succ Are there more issues like these?
- How to find these issues more effectively?

The difficulty of fuzzing the fd type confusion caused by race condition:

- The buggy code is lurking in kernel, the user process can barely notice it! ----->
- The race window can be tiny!

Maybe we can detect such issues at runtime by some detecting code?

CVE-2022-21772

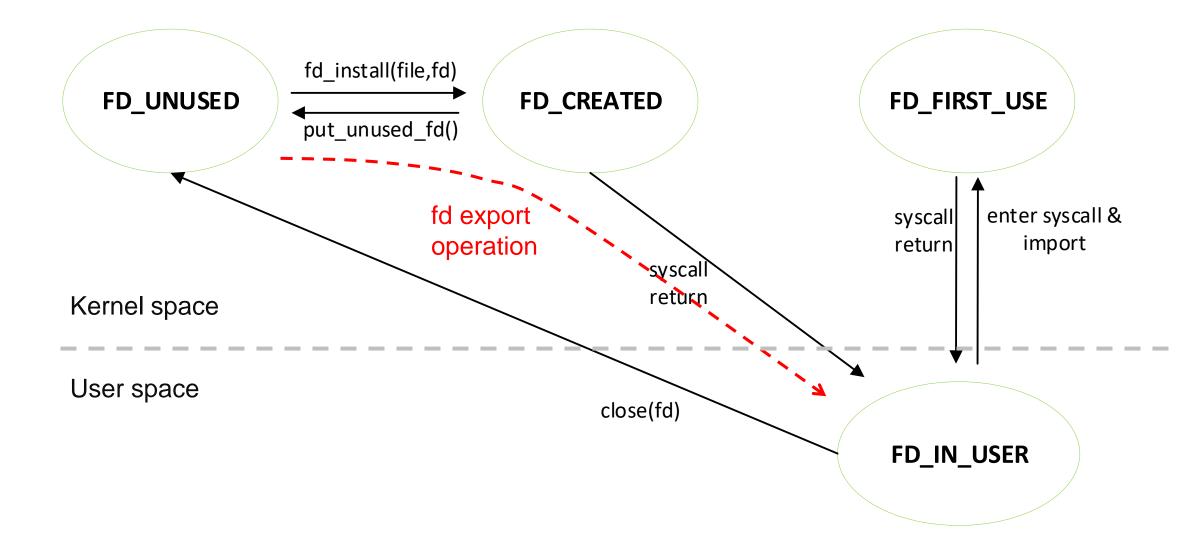
```
close(fd);
. . .
```

fd = teec_shm_alloc(ctx->fd, s, &shm->id);

dma_buf = dma_buf_get(fd);

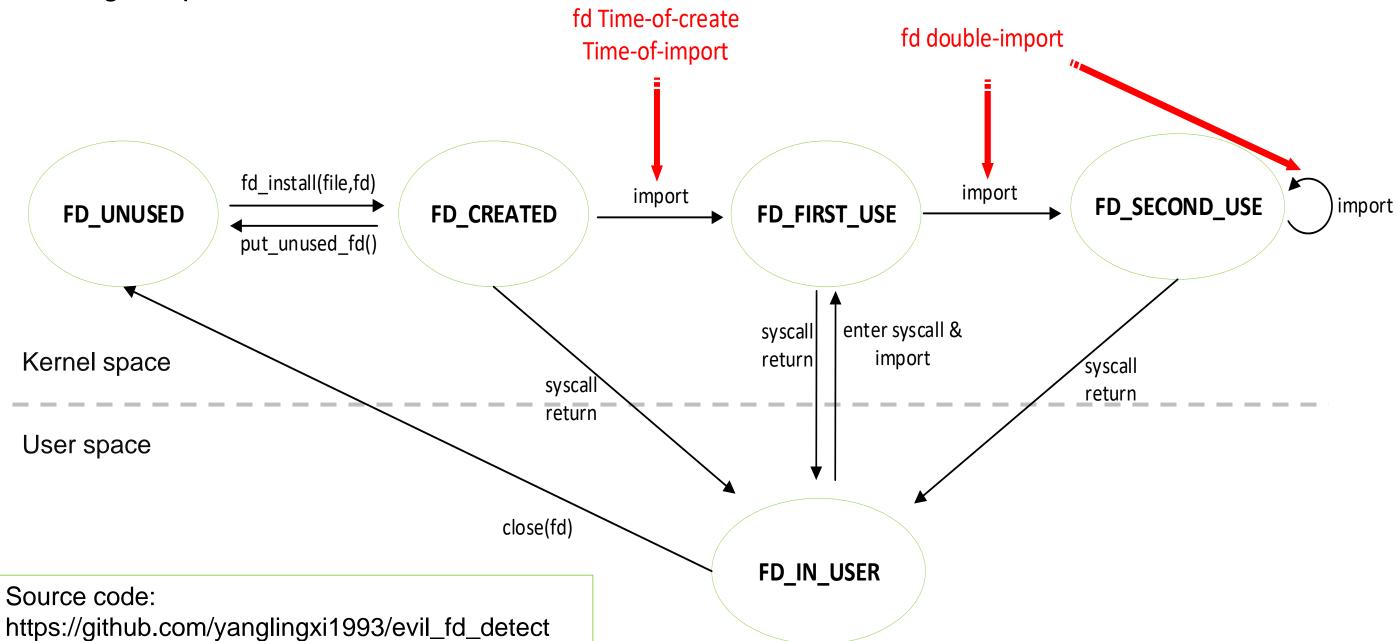
Find the issues

Regular lifecycle of an fd:



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Find the issues



Detecting the potential issues:

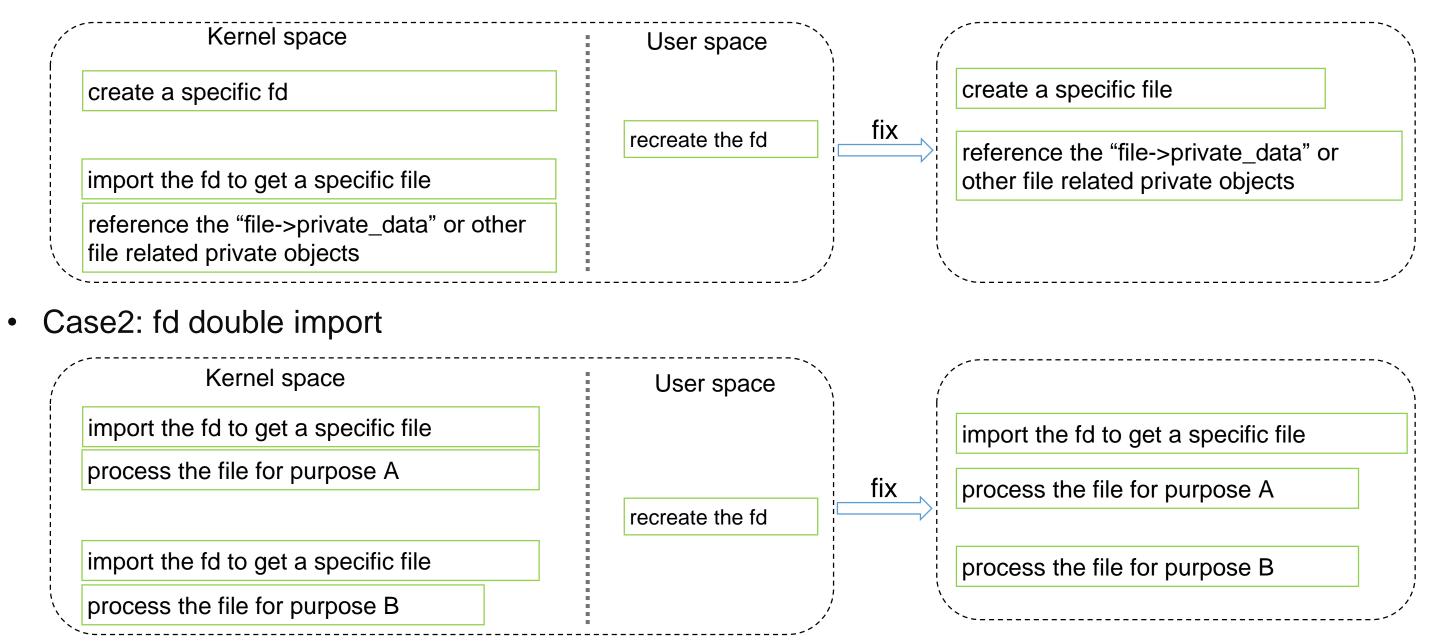
Bug hunting result

type	From	CVE-id/issue	Found by
fd time-of-create time-of-import	Vendor M	CVE-2022-21772	code auditing
		Issue#1	detect tool
		lssue#2	detect tool
	Vendor S	Issue#1	code auditing
	Vendor Q	Issue#1	detect tool
fd double import	Vendor M	CVE-2022-20082	code auditing
		Issue#1	detect tool
		lssue#2	detect tool
	Vendor Q	Issue#1	code auditing
		lssue#2	code auditing
		lssue#3	code auditing



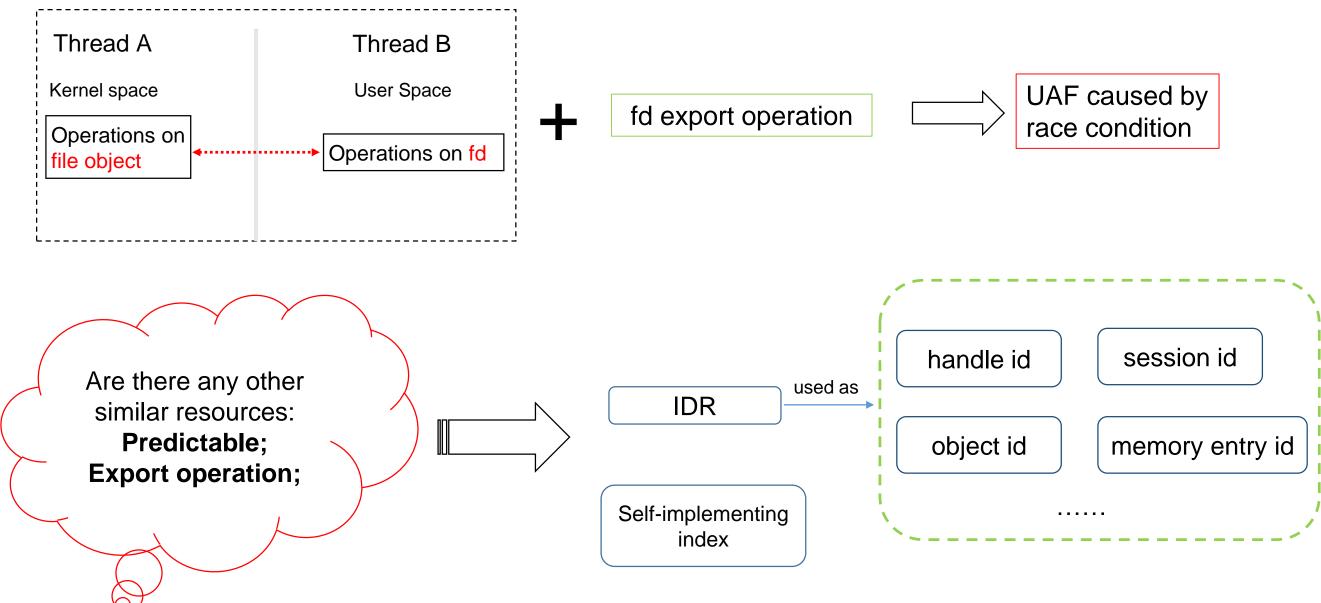
Fixes

• Case1: fd time-of-create time-of-import



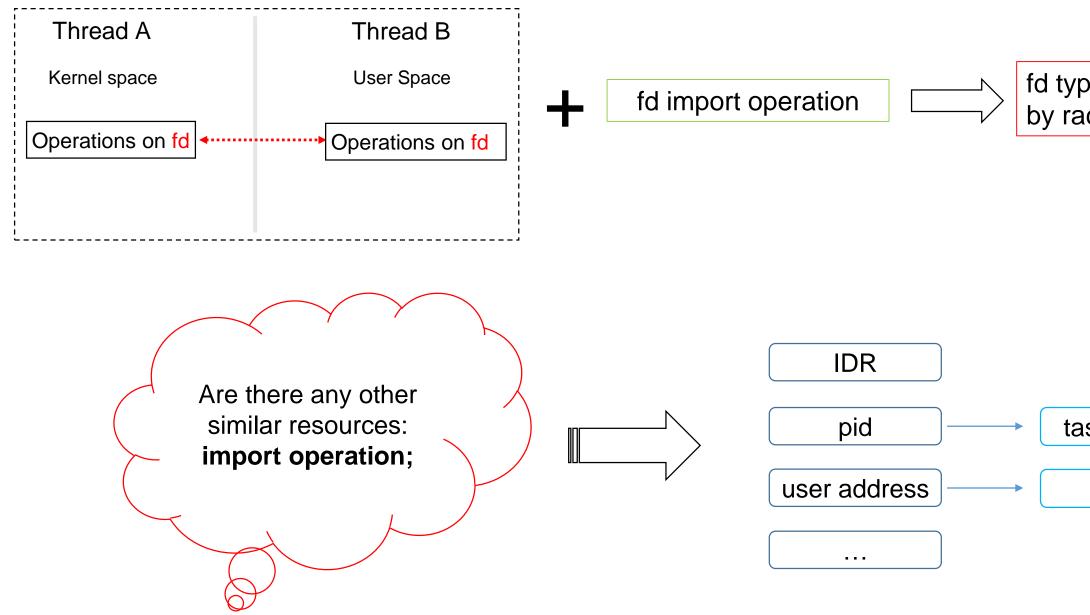
Conclusion & Future work

Race condition 1



Conclusion & Future work

Race condition 2



fd type confusion caused by race condition

task_struct

vma

Acknowledge

Thanks to 其因幡, Ye Zhang, Chenfu Bao, Shufan Yang, Lin Wu, Yakun Zhang, Zheng Huang, Tim Xia

Supplement

• Exploit of CVE-2022-28350

Small race windows can be exploitable!

- UAF caused by race condition in fd export operation
- Fd type confusion caused by race condition in fd import operation •

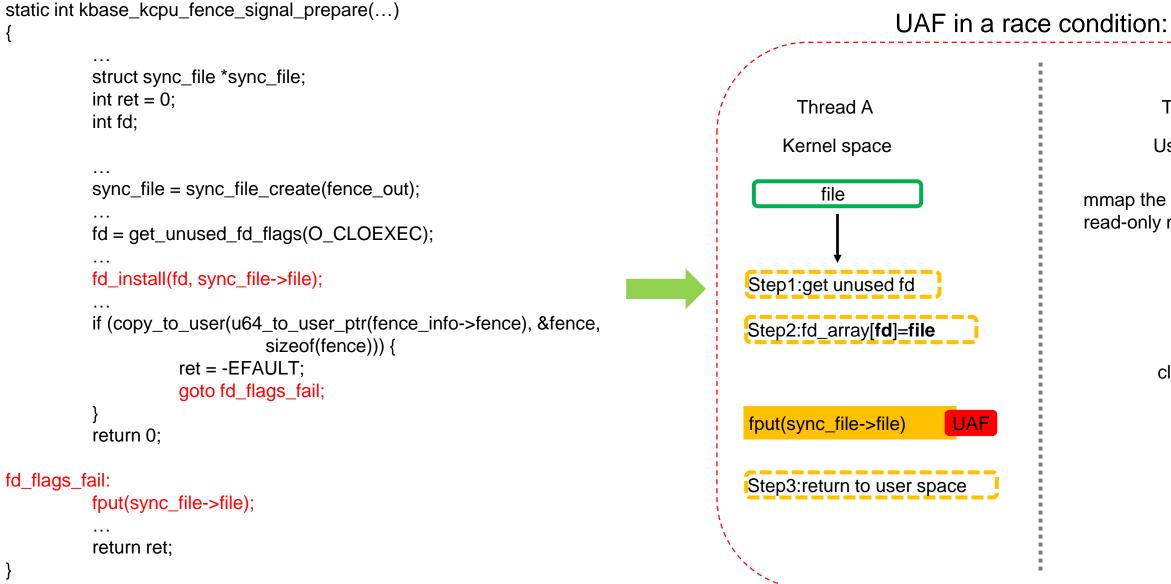
Supplement

• Exploit of CVE-2022-28350

• Small race windows can be exploitable!

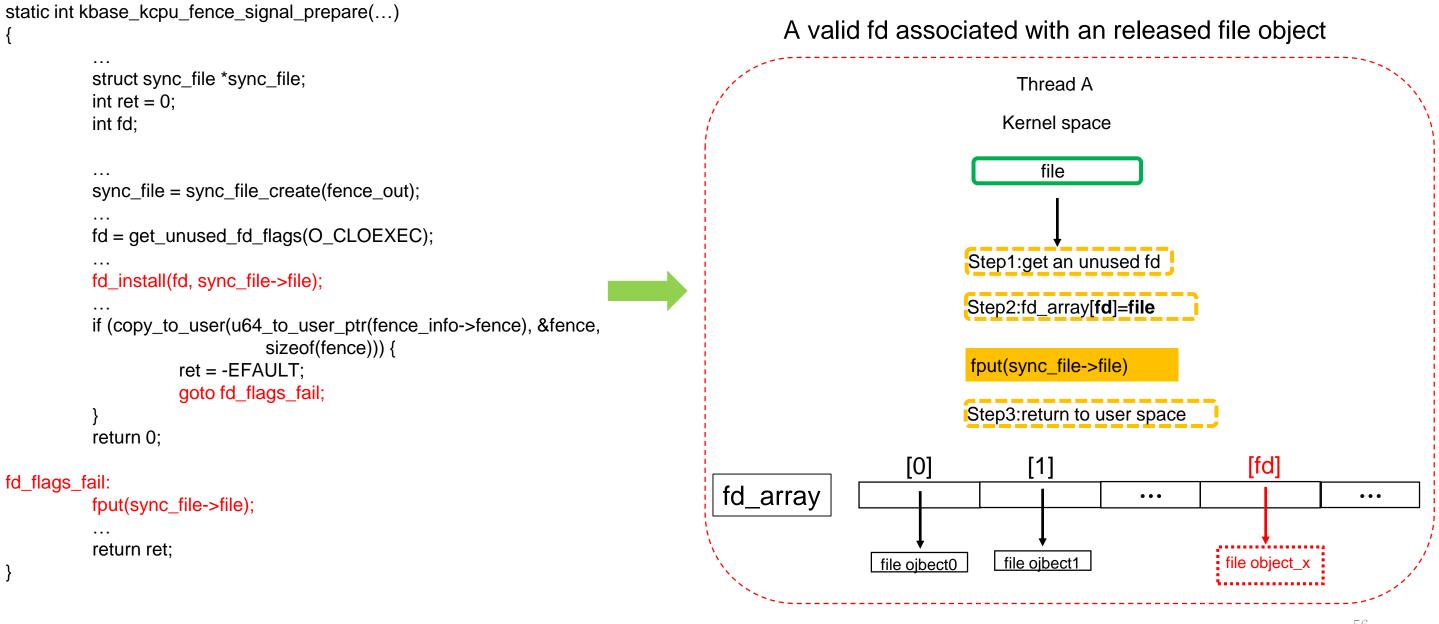
- UAF caused by race condition in fd export operation
- Fd type confusion caused by race condition in fd import operation •

What will CVE-2022-28350 lead to?



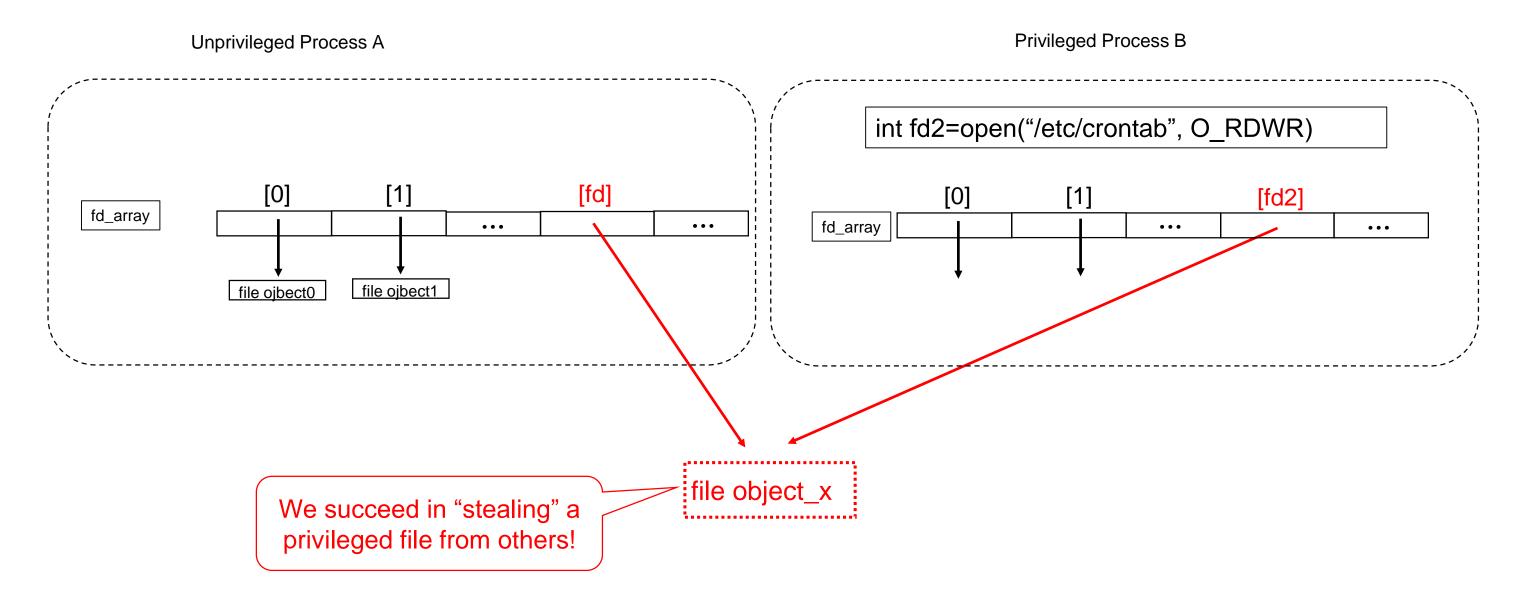
Thread B User Space mmap the "fence_info->fence" to read-only memory close(fd);

But the CVE-2022-28350 can do more:

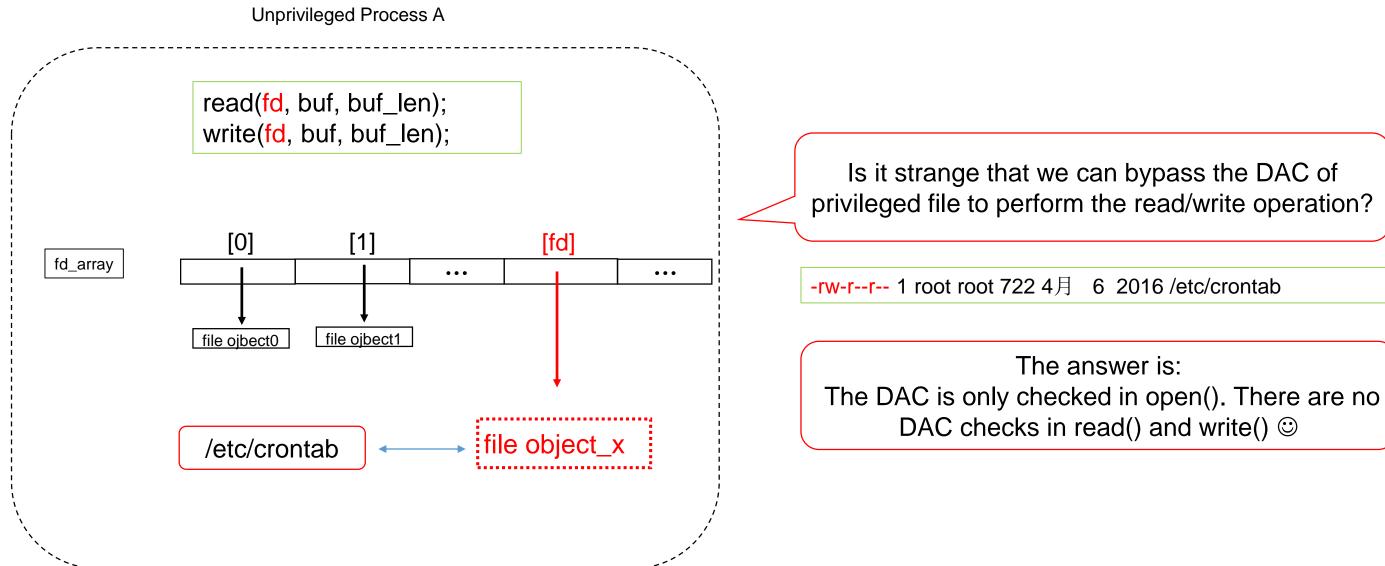


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So what if the released file object get reused by some other privileged processes when opening a privileged file?

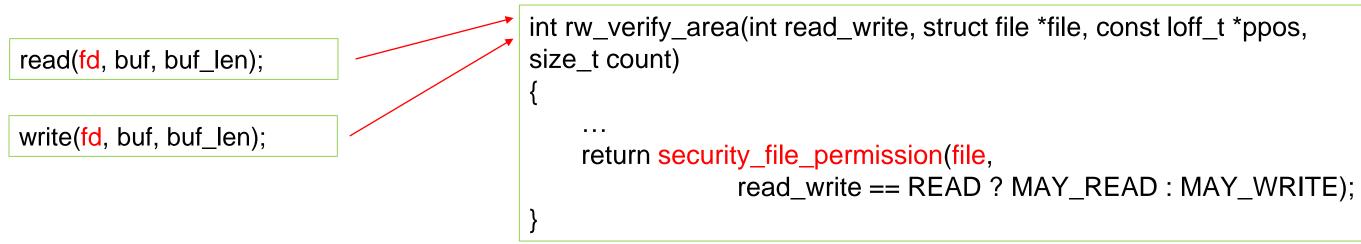


If the SELinux is disabled, the unprivileged process will have the ability to read/write the "stolen" privileged file:



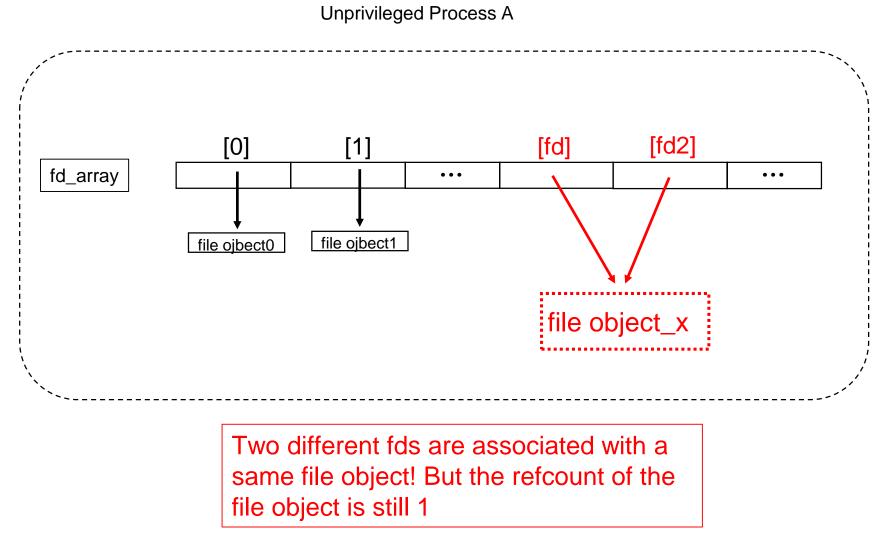
The exploitation method of "stealing" privileged file from others has been mentioned by Mathias Krause here, but this won't work on Android.

On Android, the unprivileged process cannot read/write the "stolen" privileged file because of SELinux 🙁

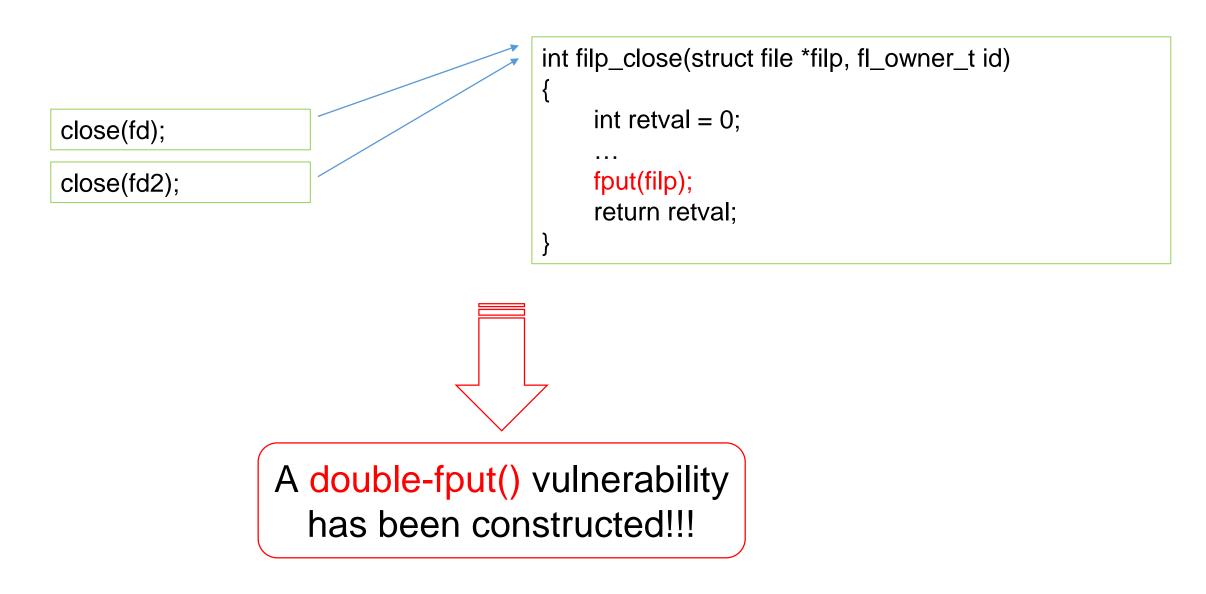


Let's find some other way out!

What if the released file object gets reused in the same process?

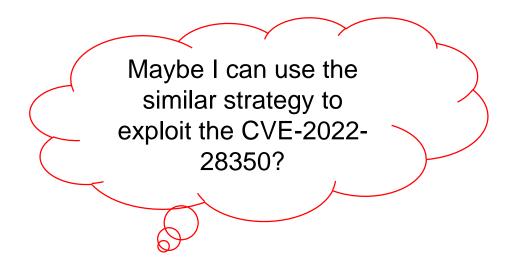


What happens if we close both fd and fd2?

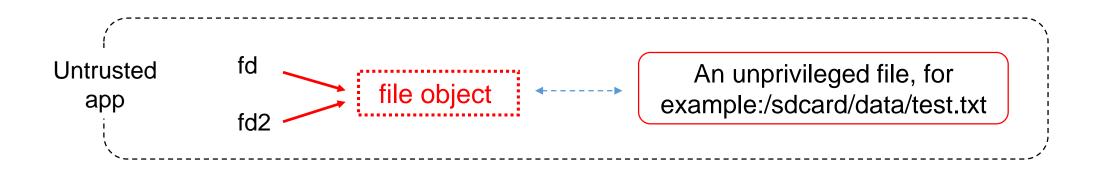


What can we do with a double-fput() vulnerability?

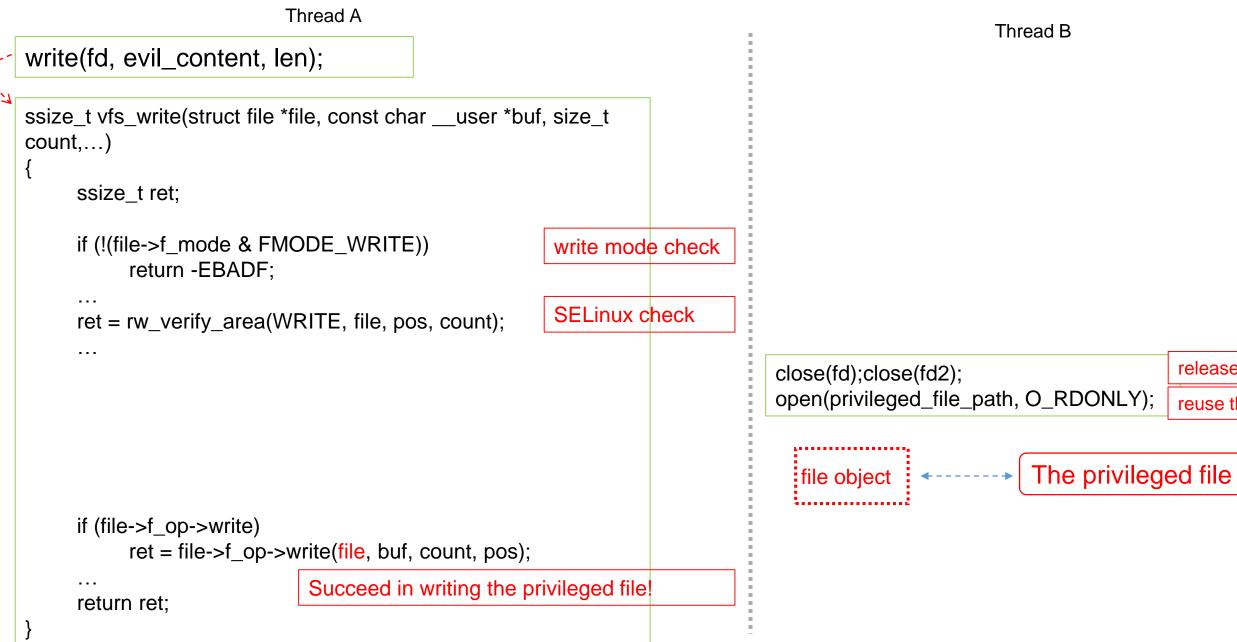
Jann Horn from Google Project Zero has given an answer to this question <u>here</u>, he showed how to write a privileged file from a unprivileged process with a double-fput() vulnerability!



Step1: Construct the scene with CVE-2022-28350



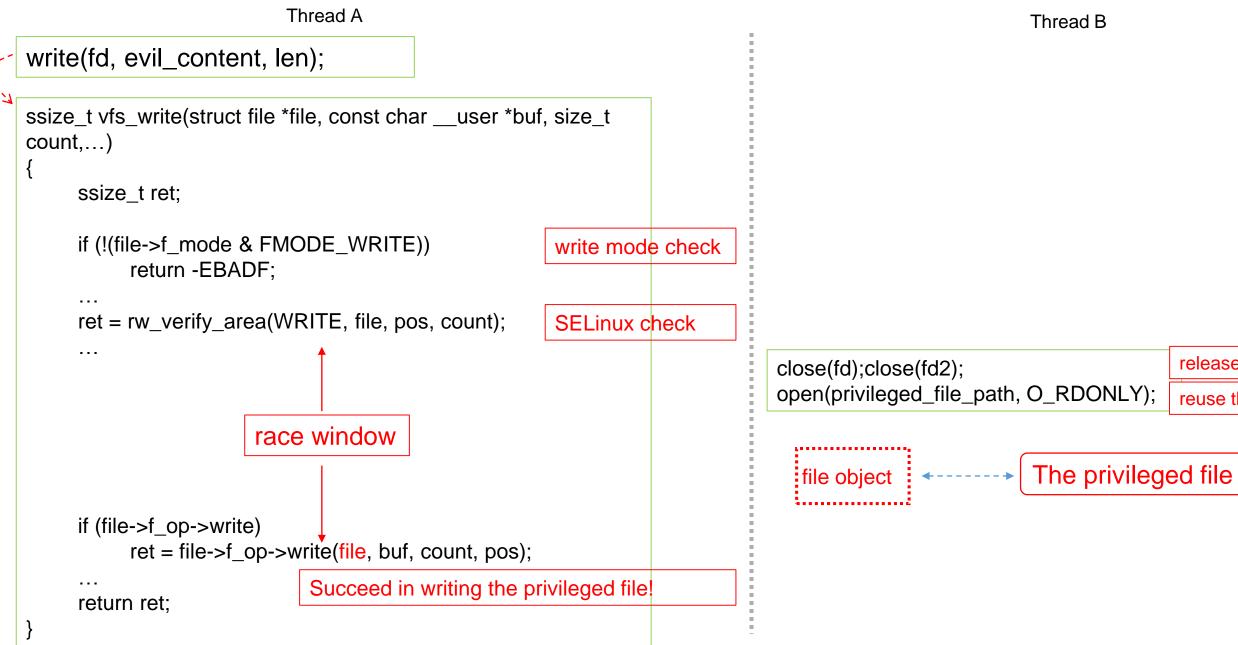
Step2: try to write the privileged file in a race condition



release the file object

reuse the file object

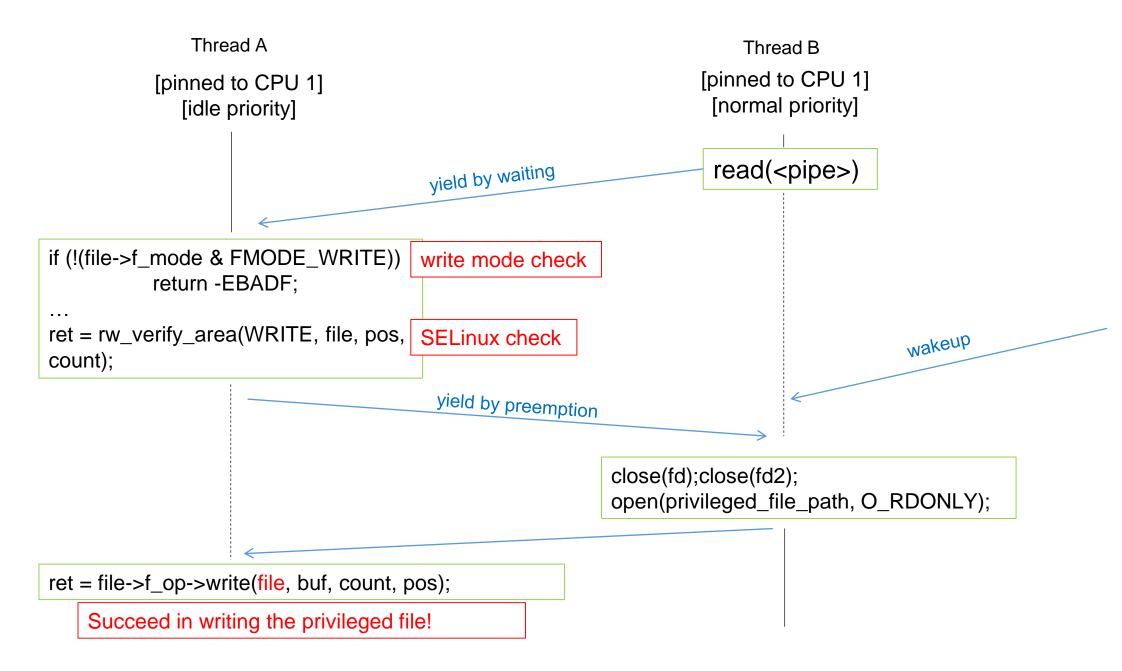
The tiny race window is still a challenge:



release the file object

reuse the file object

Try to widen the race window with the <u>method</u> given by Jann Horn:



Thread C [pinned to CPU 2] [normal priority]

write(<pipe>)

The exploit will succeed in a big chance \odot :

Tested on an affected Android 12 device



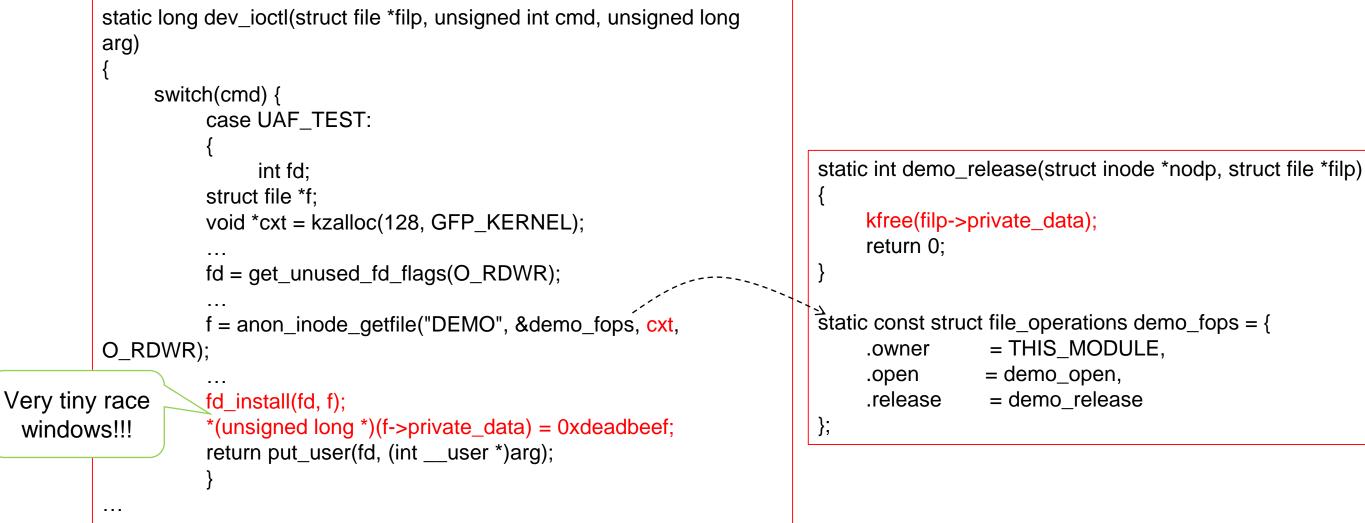
Supplement

• Exploit of CVE-2022-28350

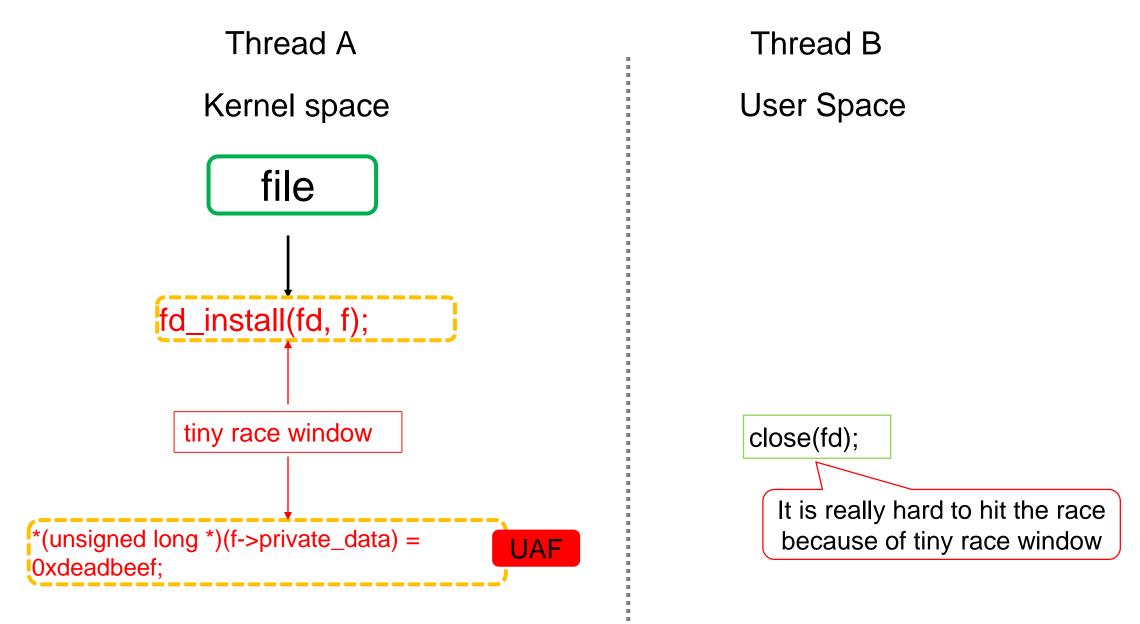
• Small race windows can be exploitable!

- UAF caused by race condition in fd export operation
- Fd type confusion caused by race condition in fd import operation •

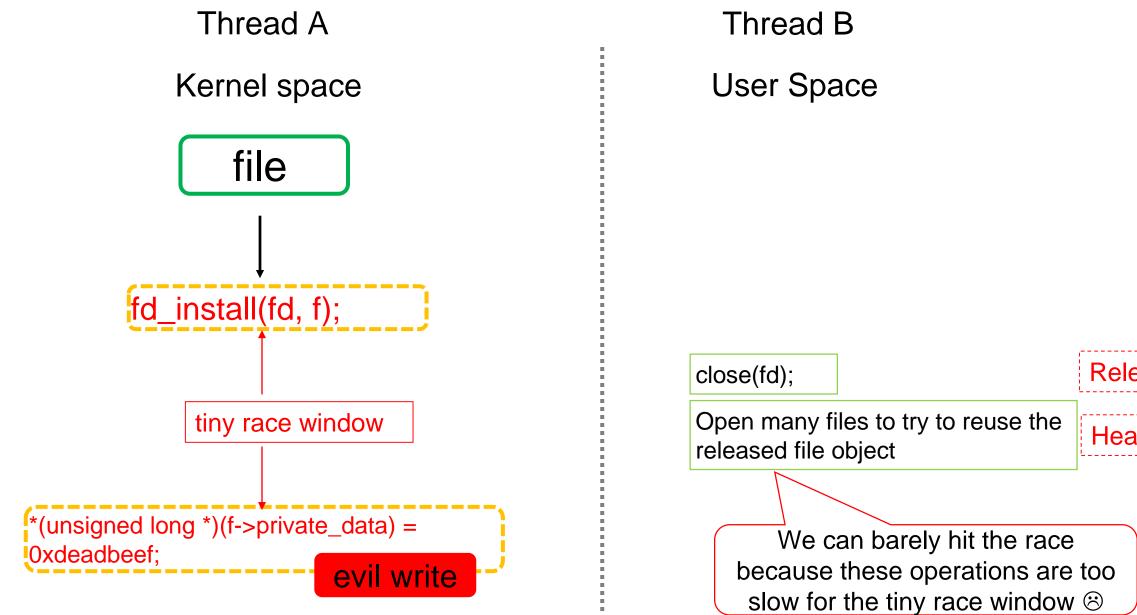
A typical issue with a tiny race window:



Try to trigger the UAF:



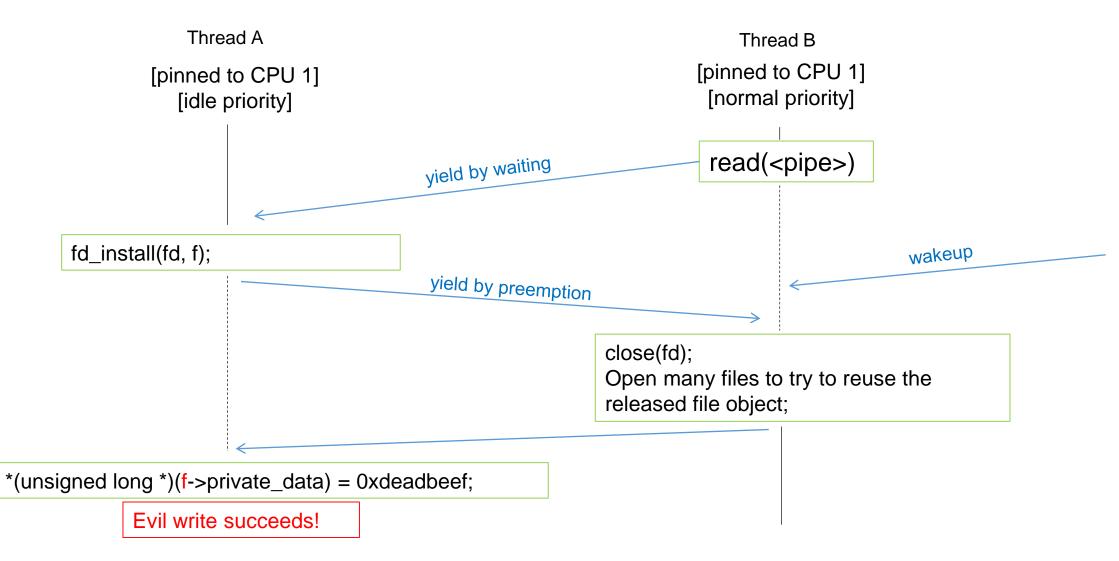
If we want to exploit the issue:



Release file object

Heap spray

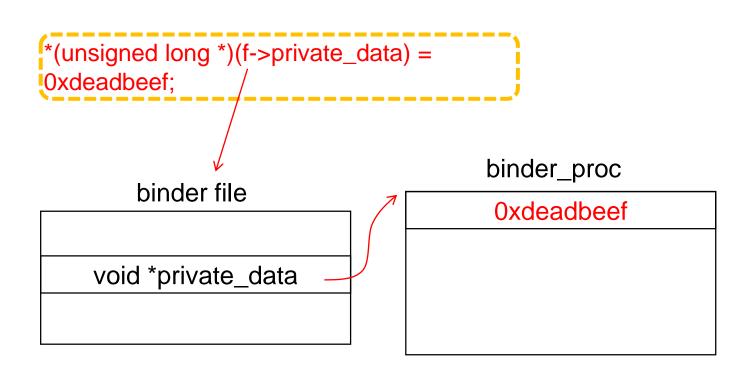
Try to widen the race window with the <u>method</u> given by Jann Horn:



Thread C [pinned to CPU 2] [normal priority]

write(<pipe>)

We have a big chance to hit the race and turn the issue to a memory corruption:



<pre>[T289] Unable to handle kernel paging request at v: [T289] Mem abort info: [T289] ESR = 0x96000045 [T289] EC = 0x25: DABT (current EL), IL = 32 bits [T289] SET = 0, FnV = 0 [T289] EA = 0, S1PTW = 0 [T289] Data abort info: [T289] Data abort info: [T289] ISV = 0, ISS = 0x00000045 [T289] CM = 0, WnR = 1 [T289] user pgtable: 4k pages, 39-bit VAs, pgdp=000 [T289] [0000000deadbef7] pgd=0000000000000000, p40</pre>
[T289] Hardware name: Raven DVT (DT)
[T289] Workqueue: events binder deferred fun
[T289] pstate: 20c00005 (nzCv daif +PAN +UA0
<pre>[T289] pc : binder deferred release+0x74/0x1</pre>
<pre>[T289] lr : binder deferred func+0x2bc/0x454</pre>
[T289] sp : ffffffc014593c80
[T289] x29: ffffffc014593cc0 x28: ffffff88e3
[T289] x27: ffffffd27e88f230 x26: ffffff88e3
[T289] x25: 000000000000003 x24: fffffd27e
[T289] x23: ffffff88e364b400 x22: ffffffd27e
<pre>[T289] x21: ffffff8810b04a00 x20: ffffff88133</pre>
<pre>[T289] x19: ffffff8810b04a00 x18: ffffffc012</pre>
[T289] x17: 000000000000002 x16: 0000000000
[T289] x15: 00000000000000 x14: 000000000
[T289] x13: ffffffd27bfb3a00 x12: 0000000000
[T289] x11: 000000080200000 x10: fffffff23
[T289] x9 : ffffff88e6963400 x8 : 00000000de
[T289] x7 : 0000000000000001
[T289] x6 : 000000000000000
[T289] x5 : 000000080200000 x4 : fffffff23
[T289] x3 : 000000080200000 x2 : ffffff8810
[T289] x1 : 00000000000000 x0 : 0000000000
[T289] Call trace:
<pre>[T289] binder_deferred_release+0x74/0x103c</pre>
[T289] binder_deterred_tunc+0x2bc/0x454
[T289] process_one_work+0x248/0x820
[T289] worker_Thread+0x438/0xbd8
[T289] kthread+0x150/0x200
[T289] ret_from_fork+0x10/0x30

irtual address 00000000deadbef7

5

0000095c9d3000 d=000000000000000, pud=0000000000000000

c.cfi_jt -TCO BTYPE=--.03c 64b450 64b400 88f000 6e9000 254160 ffb028)0000b 000018 002d4c 6d9620 adbeef 6d9620 04a00 000000

Supplement

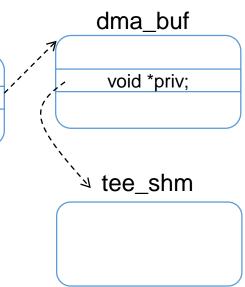
• Exploit of CVE-2022-28350

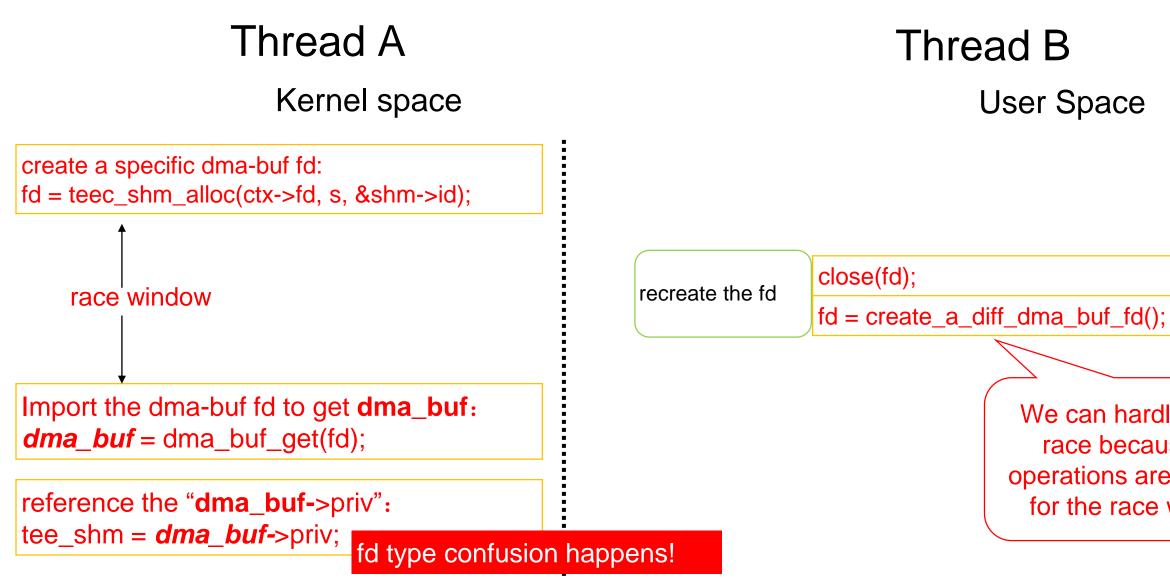
Small race windows can be exploitable!

- UAF caused by race condition in fd export operation
- Fd type confusion caused by race condition in fd import operation

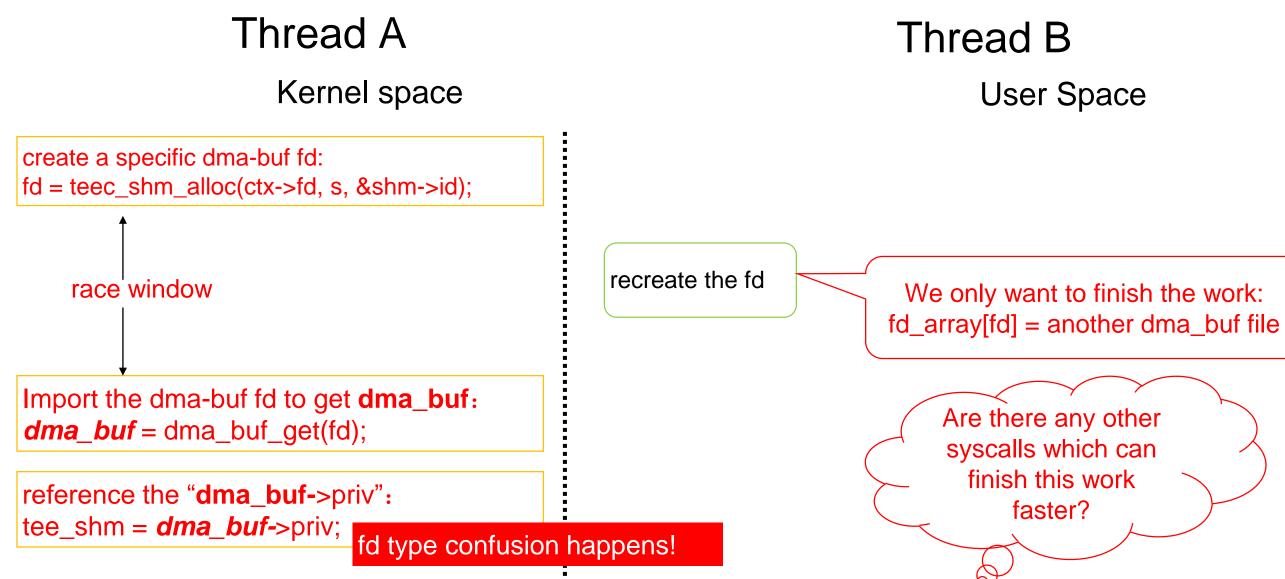
```
CVE-2022-21772
```

```
TEEC_Result TEEC_RegisterSharedMemory(struct TEEC_Context *ctx,
                                           struct TEEC SharedMemory *shm)
        int fd;
        size ts;
        struct dma buf *dma buf;
        struct tee shm *tee shm;
                                                                                                         file
        fd = teec shm alloc(ctx->fd, s, &shm->id);
                                                   create a specific dma-buf fd
                                                                                                  void *private_data;
        dma_buf = dma_buf_get(fd);
                                                   import the dma-buf fd to get dma buf
        . . .
        tee_shm = dma_buf->priv;
                                                   reference the "dma buf->priv" as tee shm
         . . .
        shm->shadow_buffer = tee_shm->kaddr;
        return TEEC_SUCCESS;
```

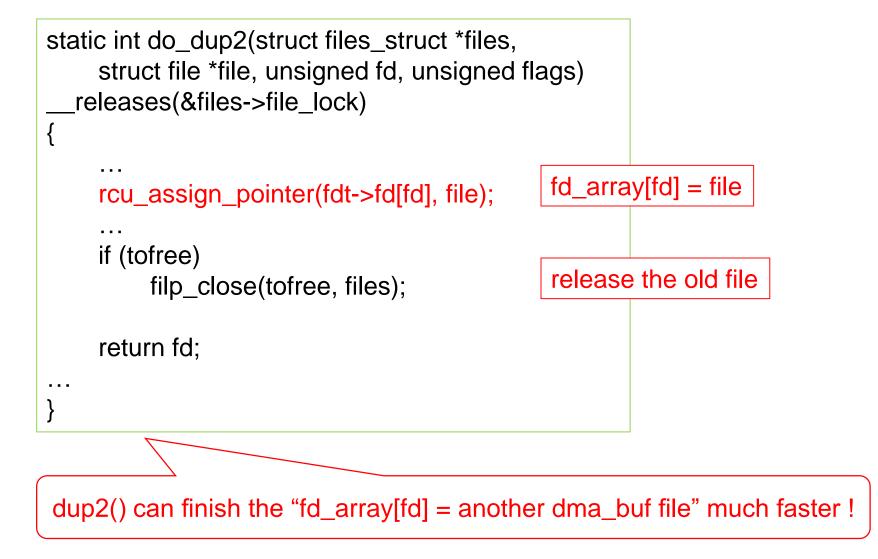


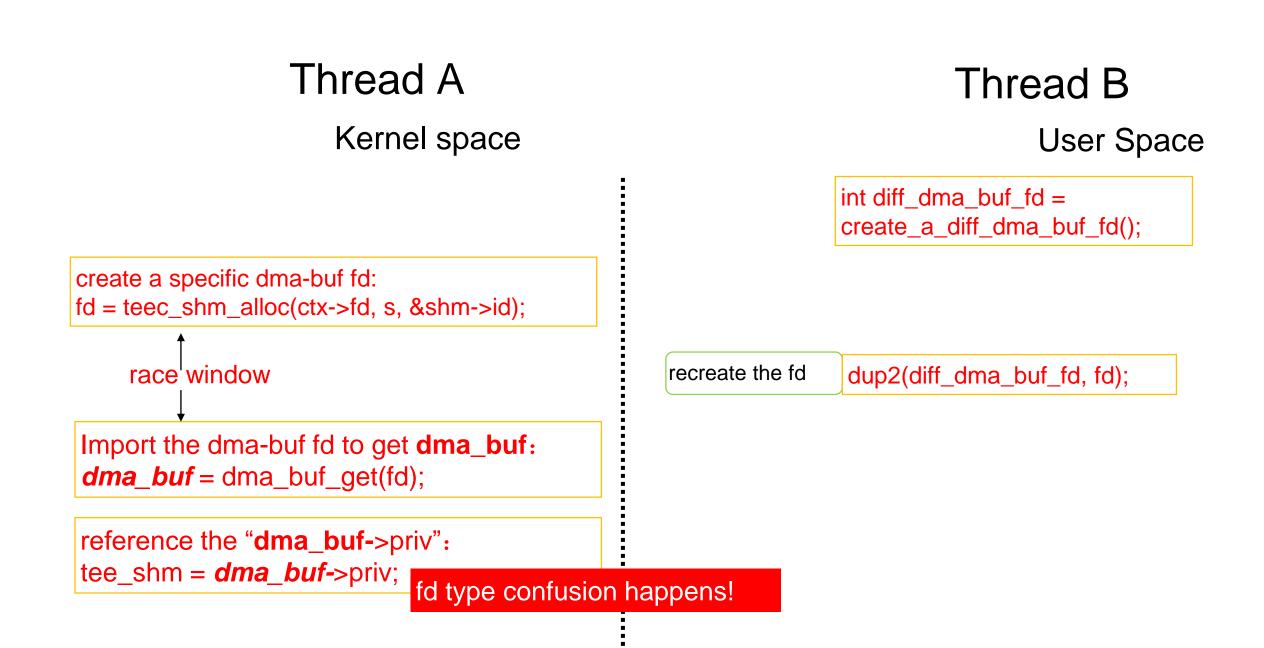


We can hardly hit the race because the operations are too slow for the race window



Syscall:dup2(int oldfd, int newfd)





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We have a big chance to hit the race and turn the issue to a memory corruption:

