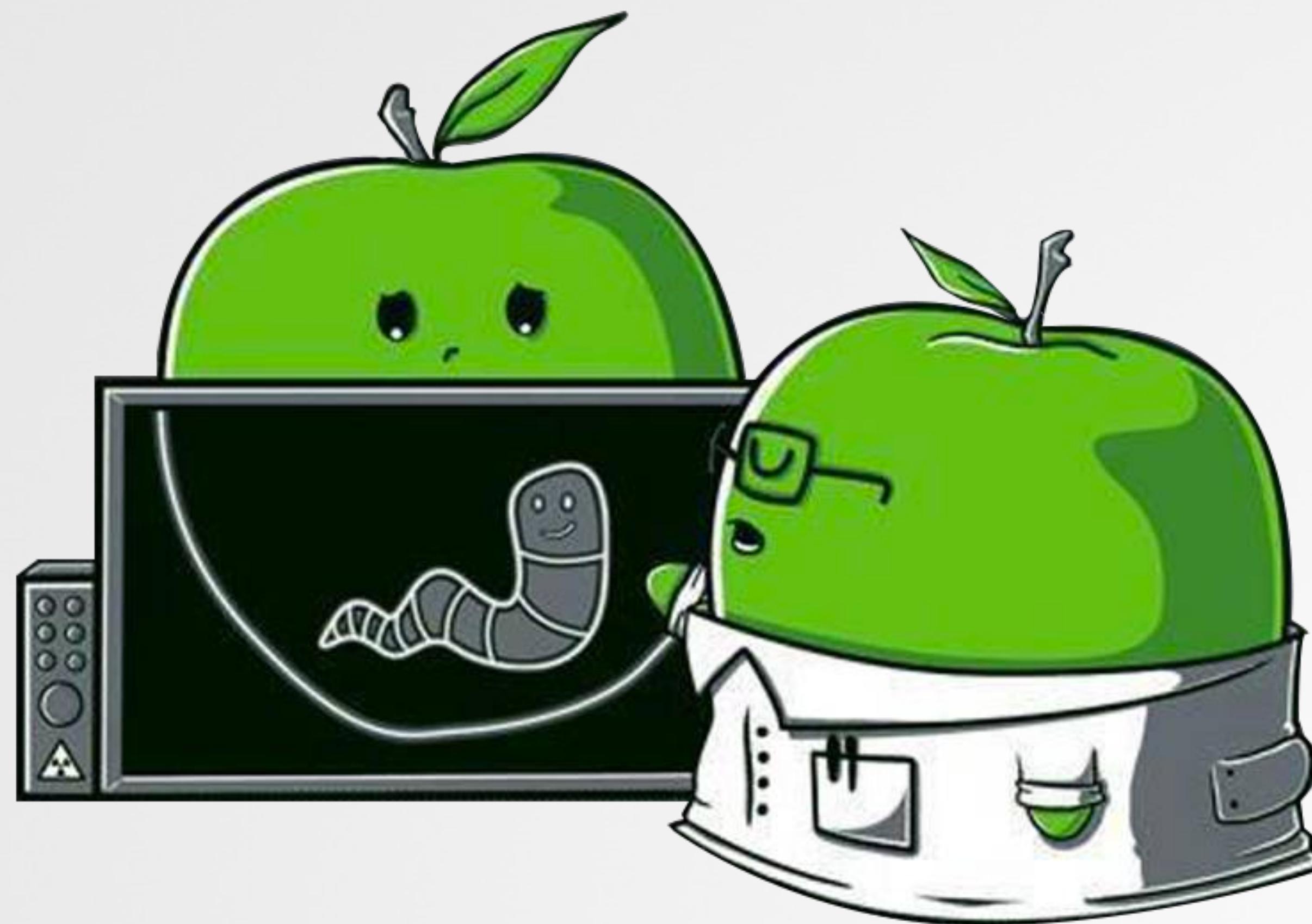


# Arm'd & Dangerous

analyzing arm64 malware targeting macOS



# WHOIS



 @patrickwardle



**Objective-See**

tools, blog, & malware collection



**#OBTS**

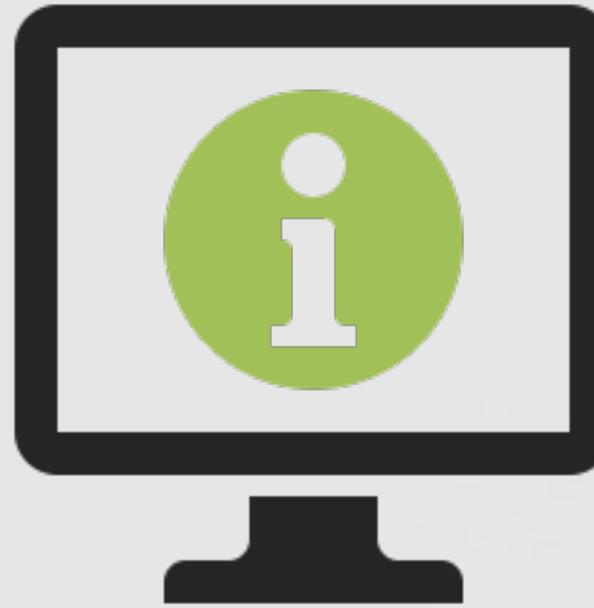
"Objective by the Sea"  
(macOS security conference)



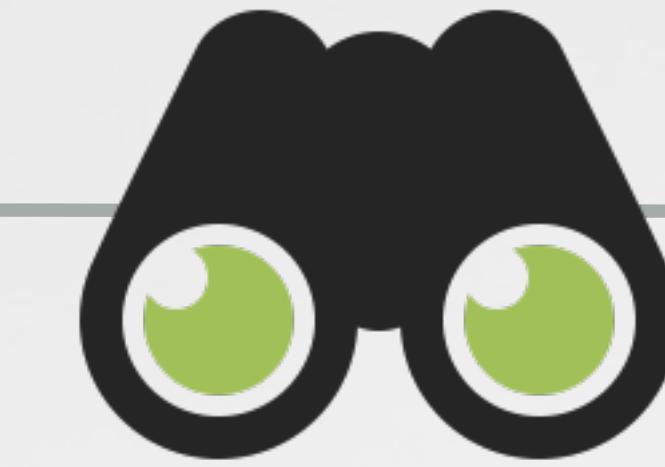
**Book(s) :**

"The Art of Mac Malware"

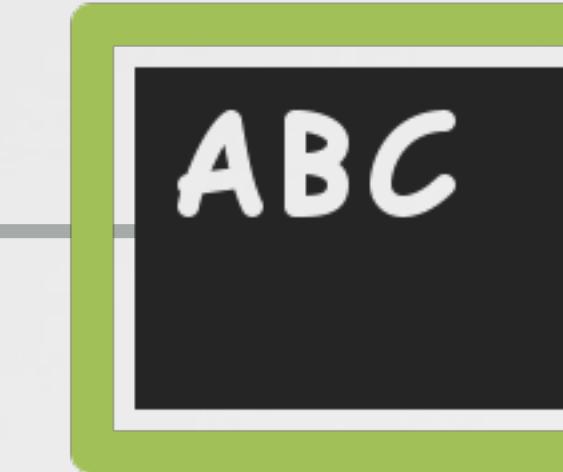
# OUTLINE



Introduction



Hunting

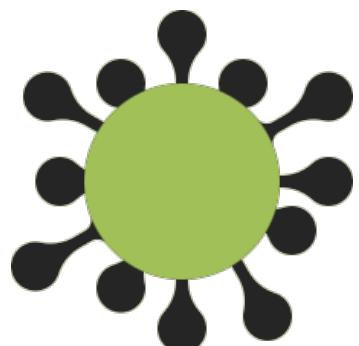


Understanding arm64



Analyzing M1 malware

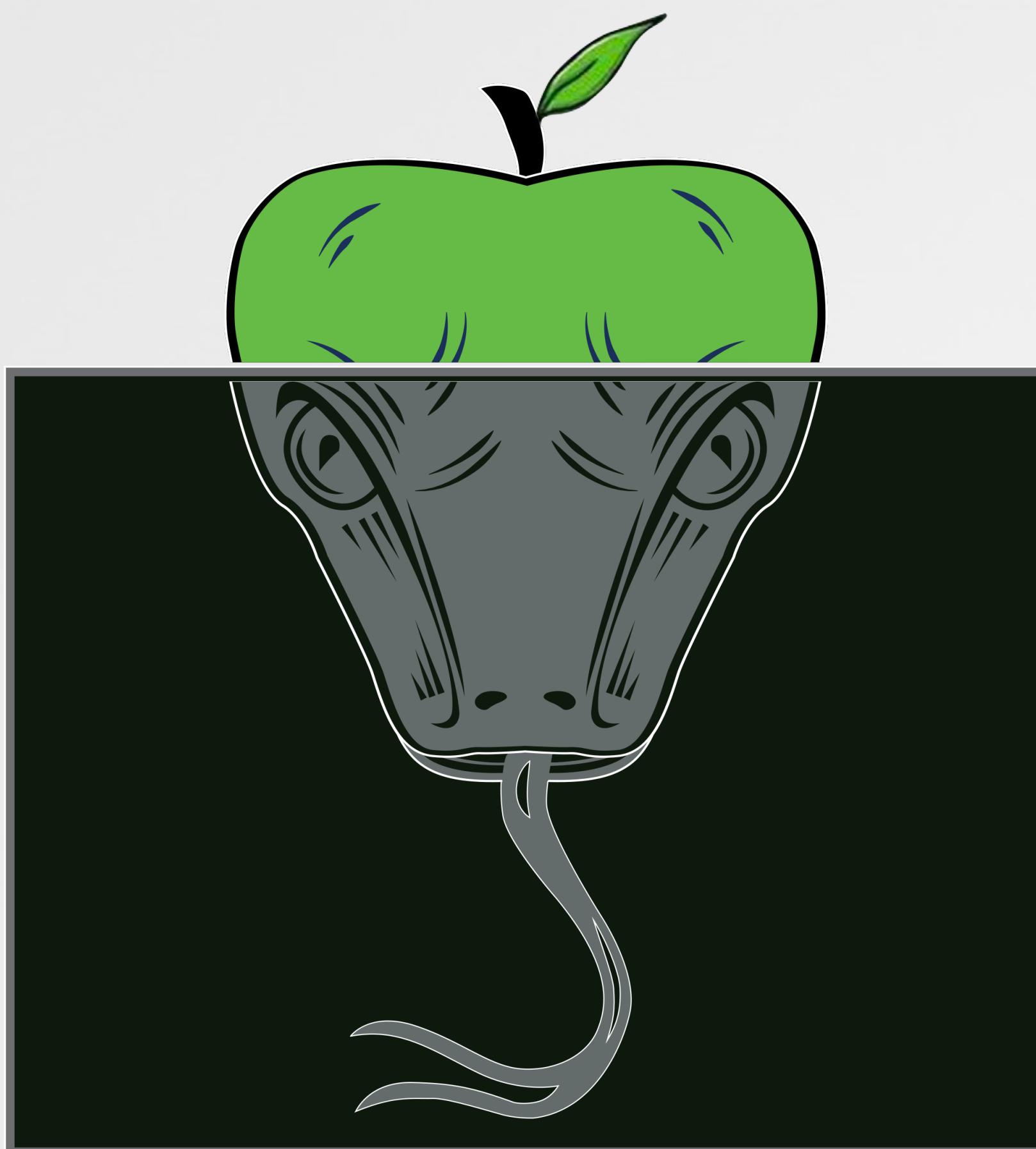
"M1 malware" defined:



Malware natively compiled to run on Apple Silicon.

More specifically, arm64 malware targeting Macs (macOS).

# Introduction



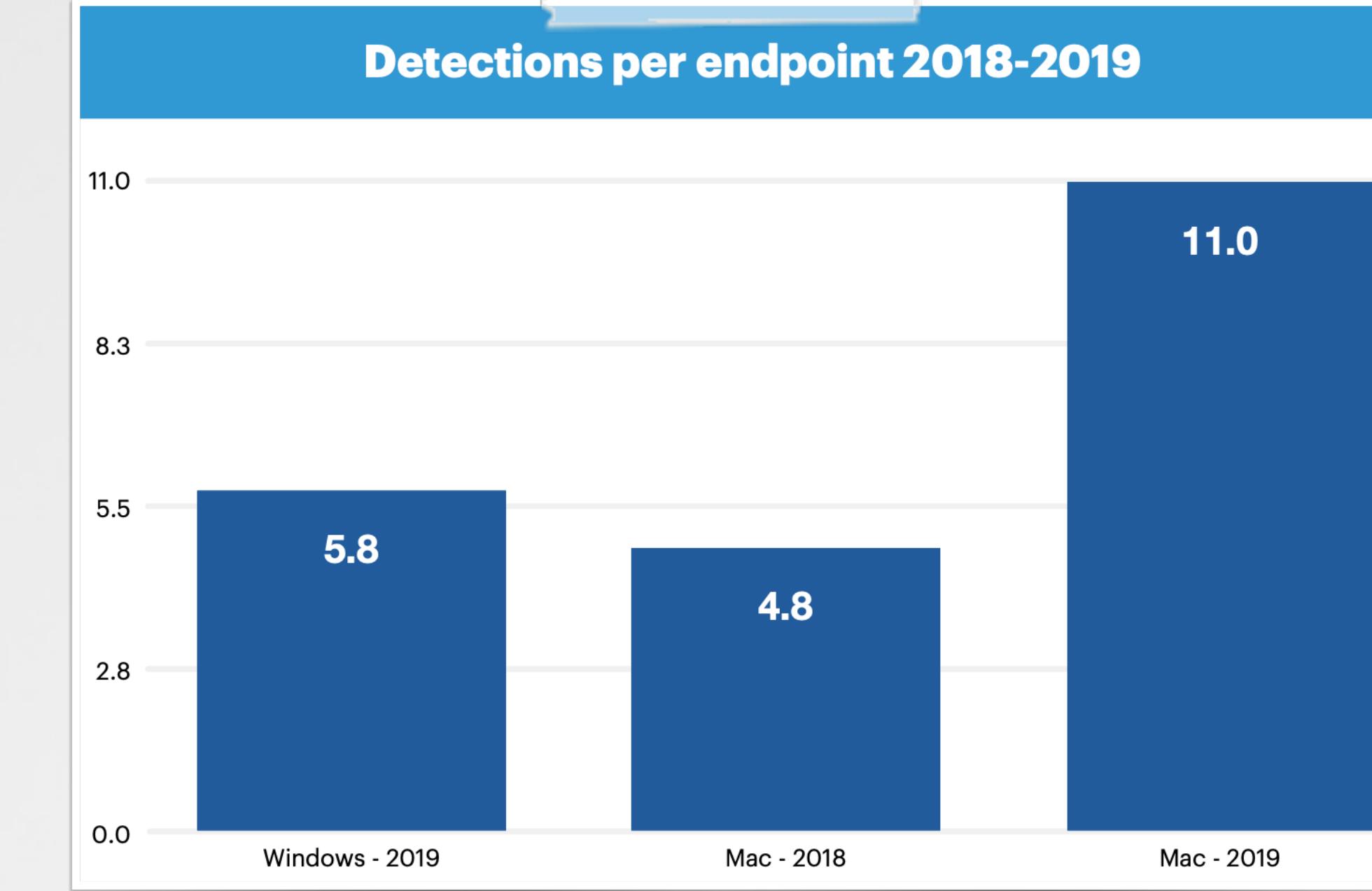
# THE GROWTH OF MACS

## ...and unsurprisingly, the growth of macOS malware

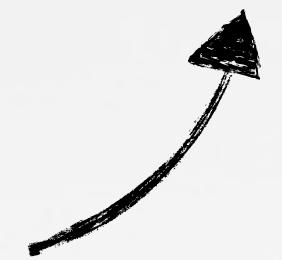
9TO5Mac ✓

IDC: Mac shipments rise  
industry-leading 49% in Q4

More Macs



more than  
Windows !?



More Mac Malware  
(credit: MalwareBytes)



As macOS becomes more prevalent, (rather obviously),  
so too does malware targeting this platform.

# THE GROWTH OF MACS

## ...driven largely by [the] M1?

Mac sales are surging thanks to Apple's new M1 processor

### Mac Sales Skyrocketing After M1 Launch

Apple's worldwide Mac shipments grew massively in the fourth quarter of 2020 after the launch of three new Macs with the [M1](#) chip, according to new PC shipping estimates [shared by Gartner](#). Apple shipped an estimated 6.9 million Macs, up from the 5.25 million it shipped at the same time in 2019, marking significant growth of 31.3 percent.

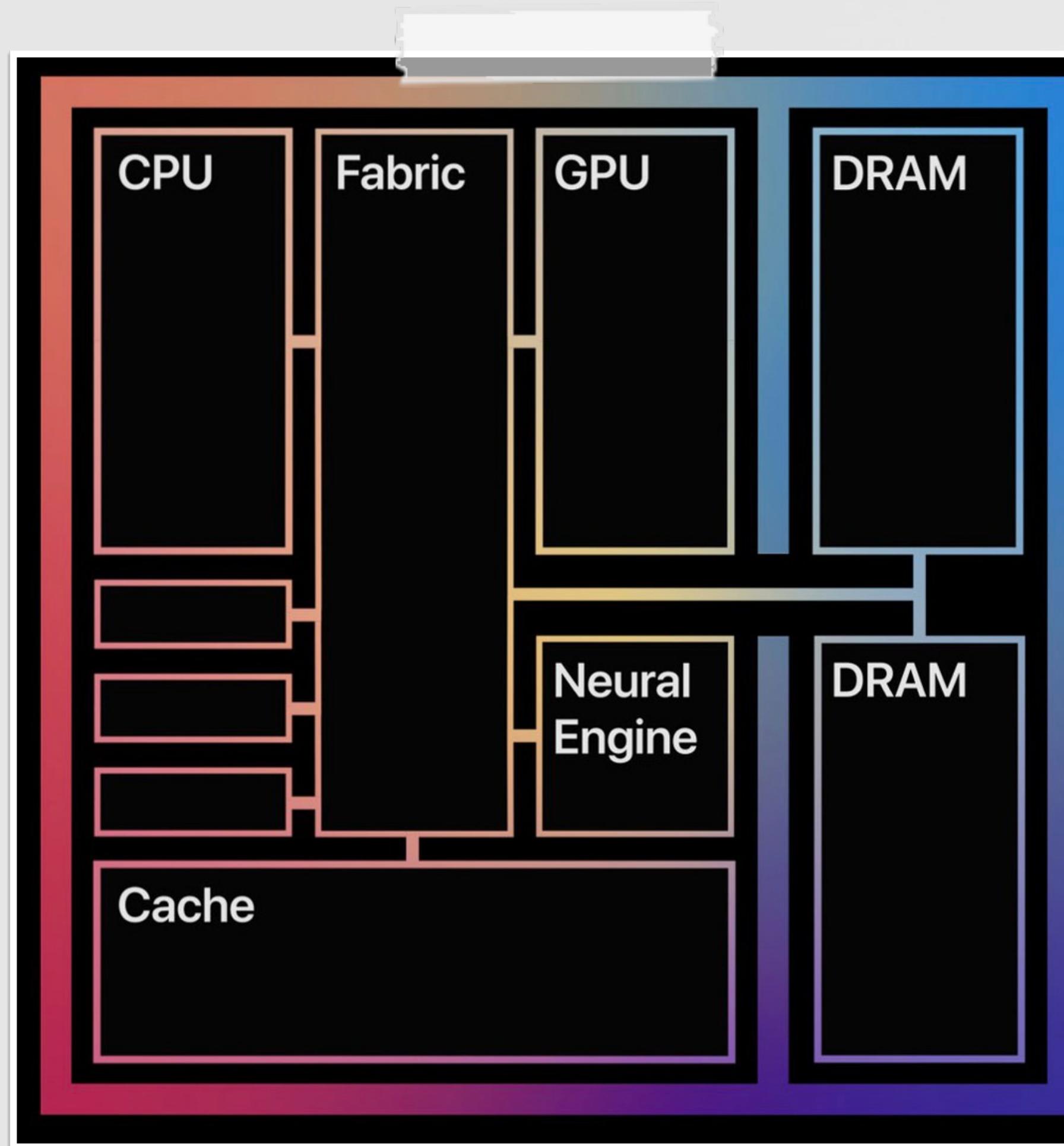
The Mac is back — and the M1 chip gets the credit

"Fueled by the M1, we set an all-time [Mac] revenue record continuing the momentum for the product category" -Tim Cook (Apple)



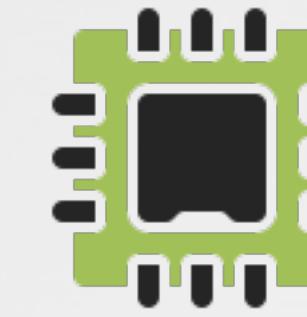
# WHAT IS M1? (AKA "APPLE SILICON")

## an arm-based system on a chip (SoC)



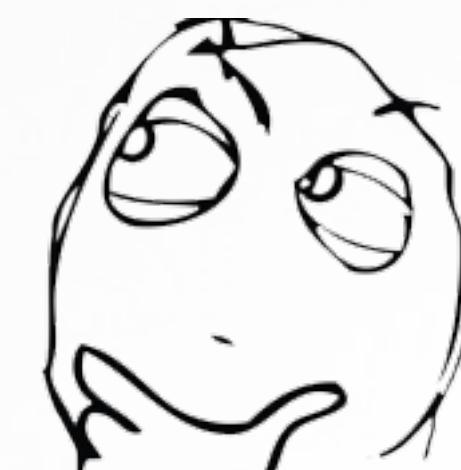
The M1 chip

Multiple technologies  
combined on a single chip



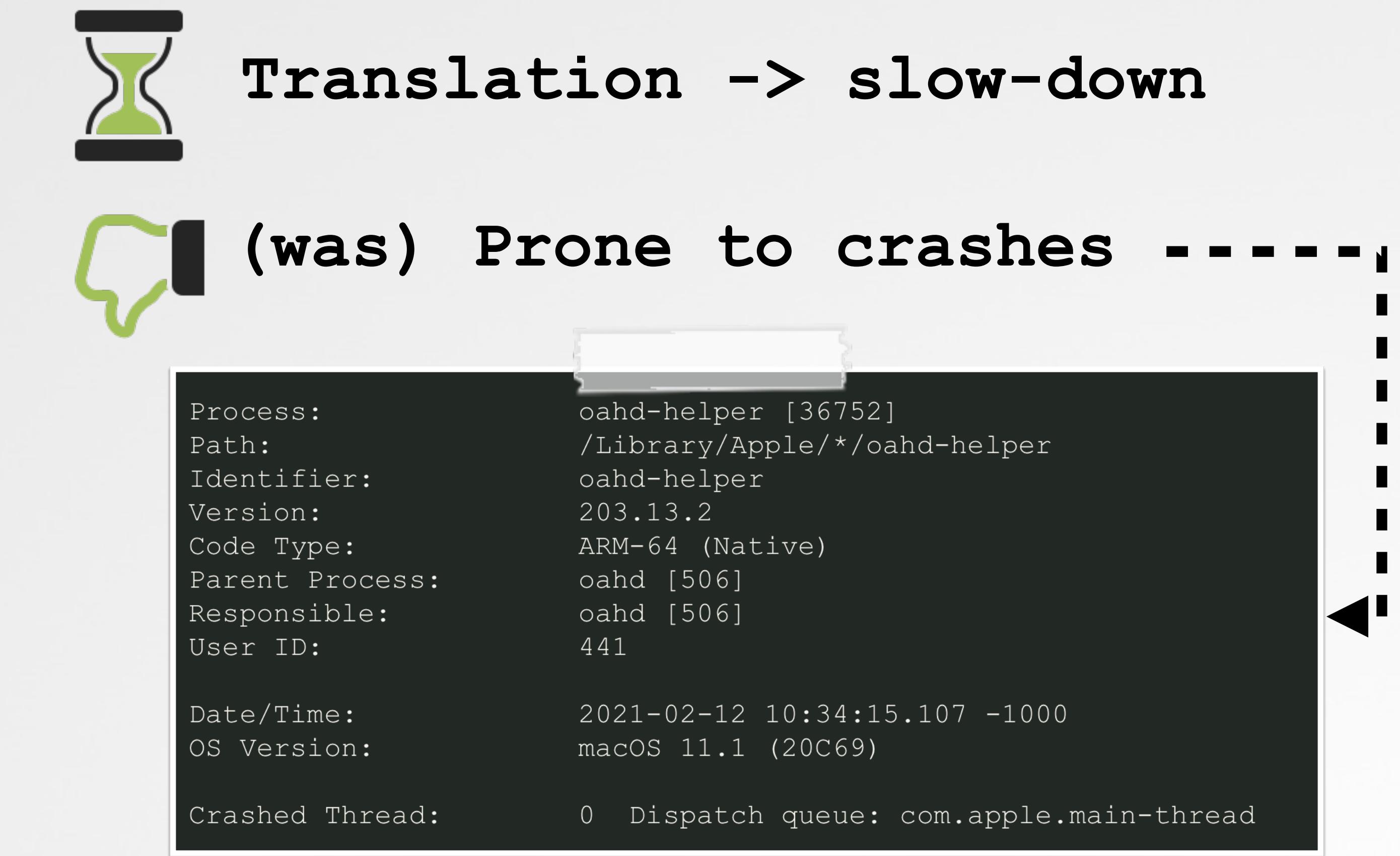
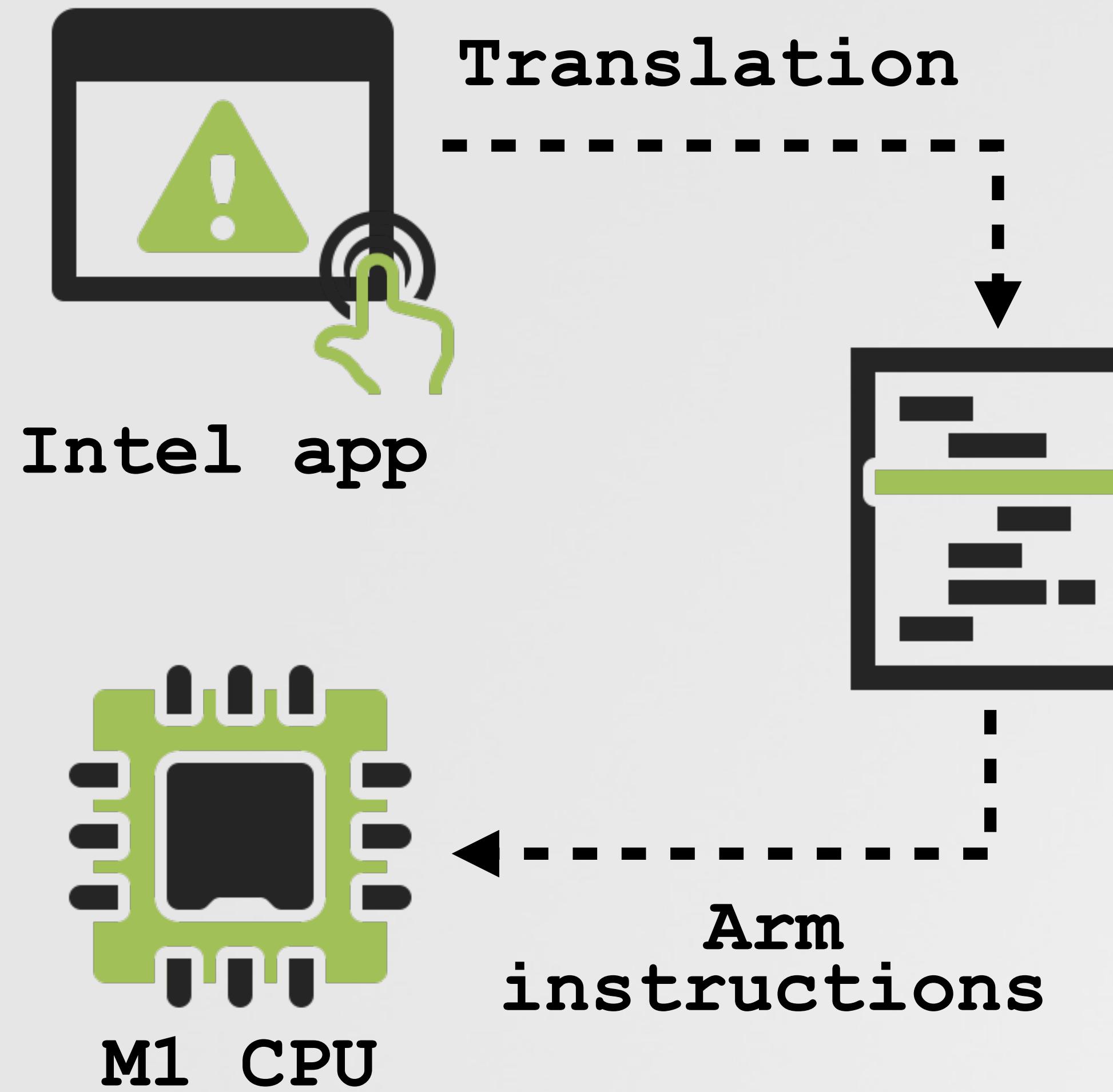
CPU:  
arm64 instruction set

...malware native to this CPU, will  
disassemble into Arm (vs. Intel).



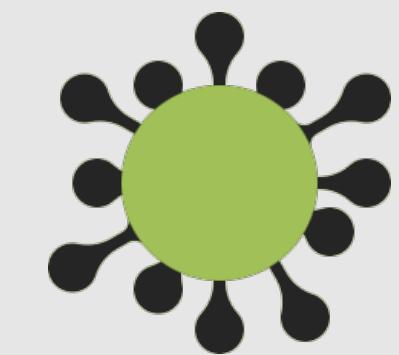
# AND ROSETTA (2)

...run intel-based apps on apple silicon



"not a substitute for creating a native version of your app" -apple

# WHY TALK ABOUT M1 MALWARE well, several important reasons!

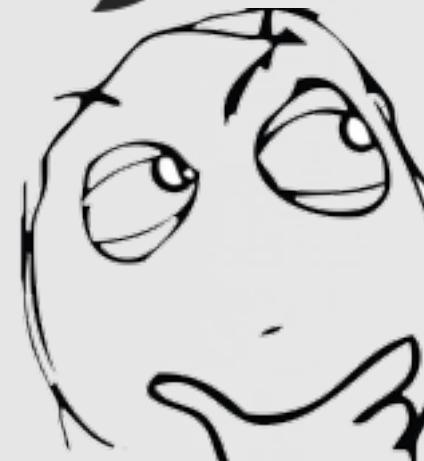


- Is (was) inevitable
  - no rosetta crashes
  - native code, faster



Disassembly is arm64

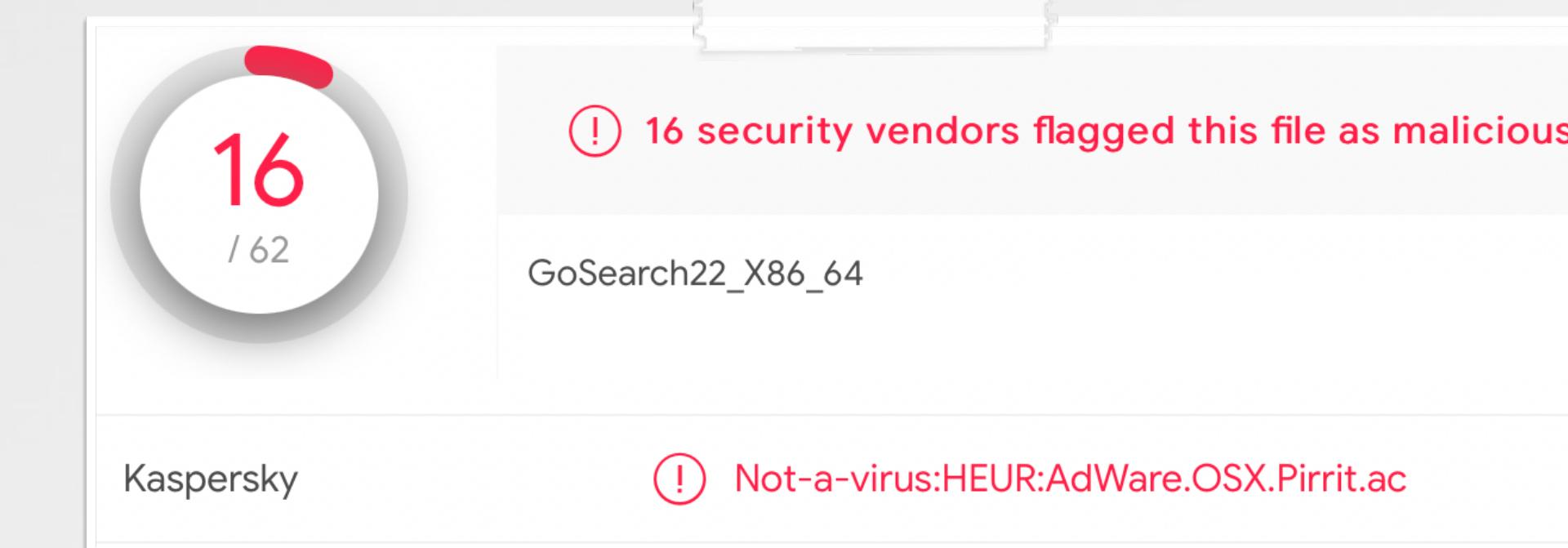
```
01 movz      x0, #0x1a
02 movz      x1, #0x1f
03 movz      x2, #0x0
04 movz      x3, #0x0
05 movz      x16, #0x0
06
07 svc       #0x80
```



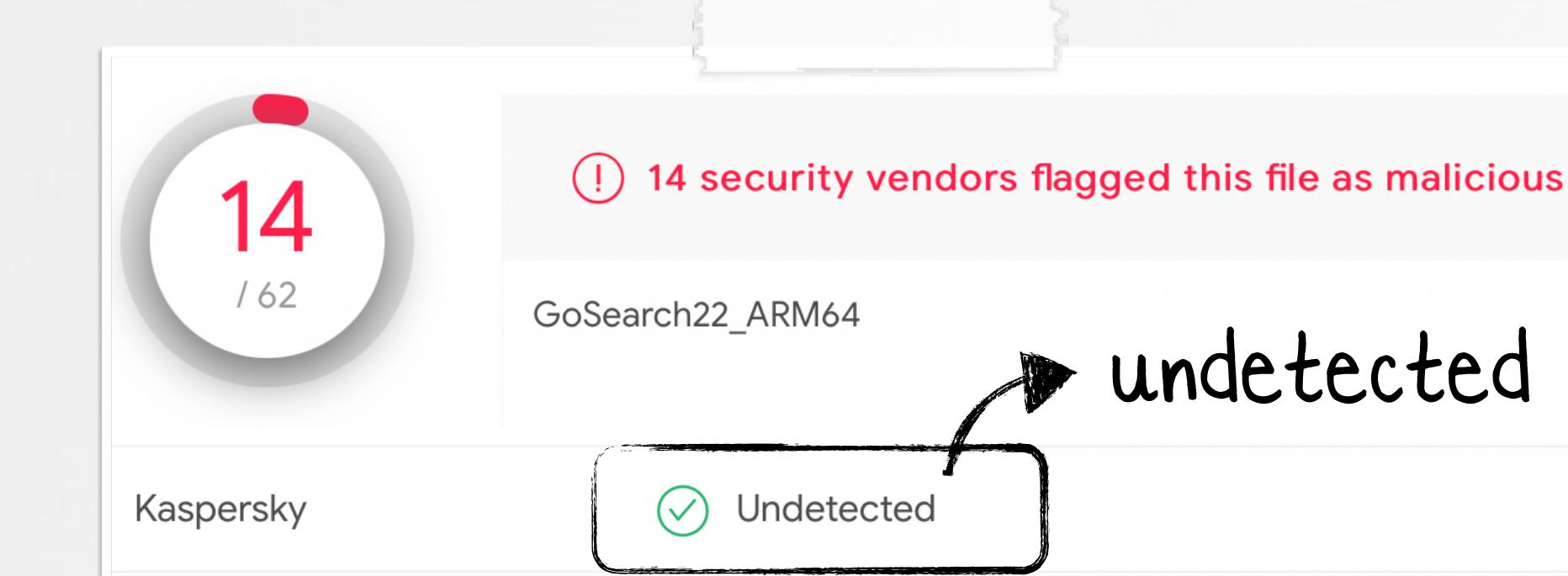
an unfamiliar instructions set!?



Missed AV detections  
(known malware: 10%+ drop)



Intel version: detected

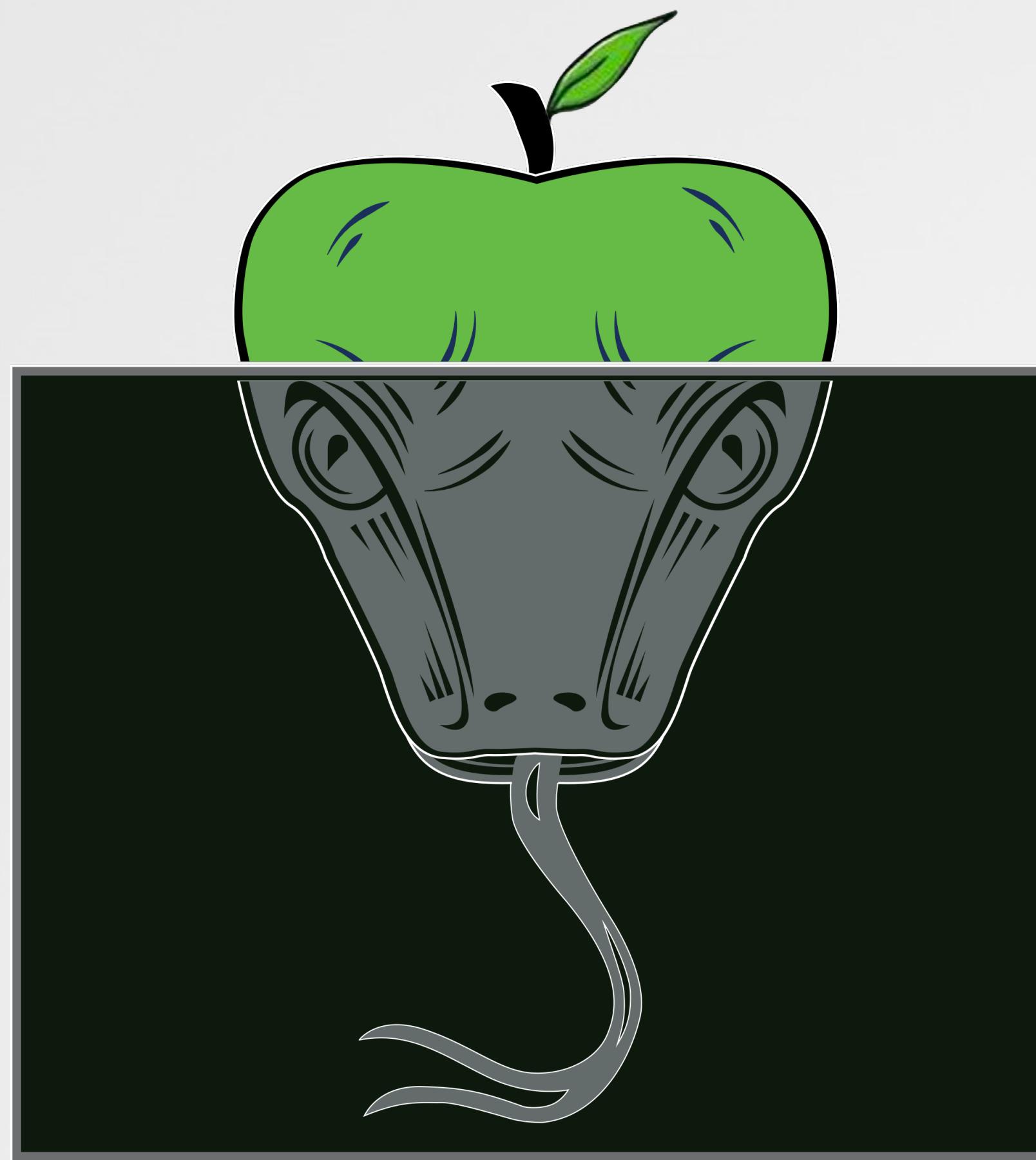


Arm version: not detected

} same malware

undetected

# Hunting for M1 malware



# HOW TO IDENTIFY M1 CODE in short, a macOS binary with arm64/e

```
% file Calculator.app/Contents/MacOS/Calculator
Mach-O universal binary with 2 architectures:
Mach-O 64-bit executable x86_64
Mach-O 64-bit executable arm64e
```

```
% lipo -archs Calculator.app/Contents/MacOS/Calculator
x86_64 arm64e
```

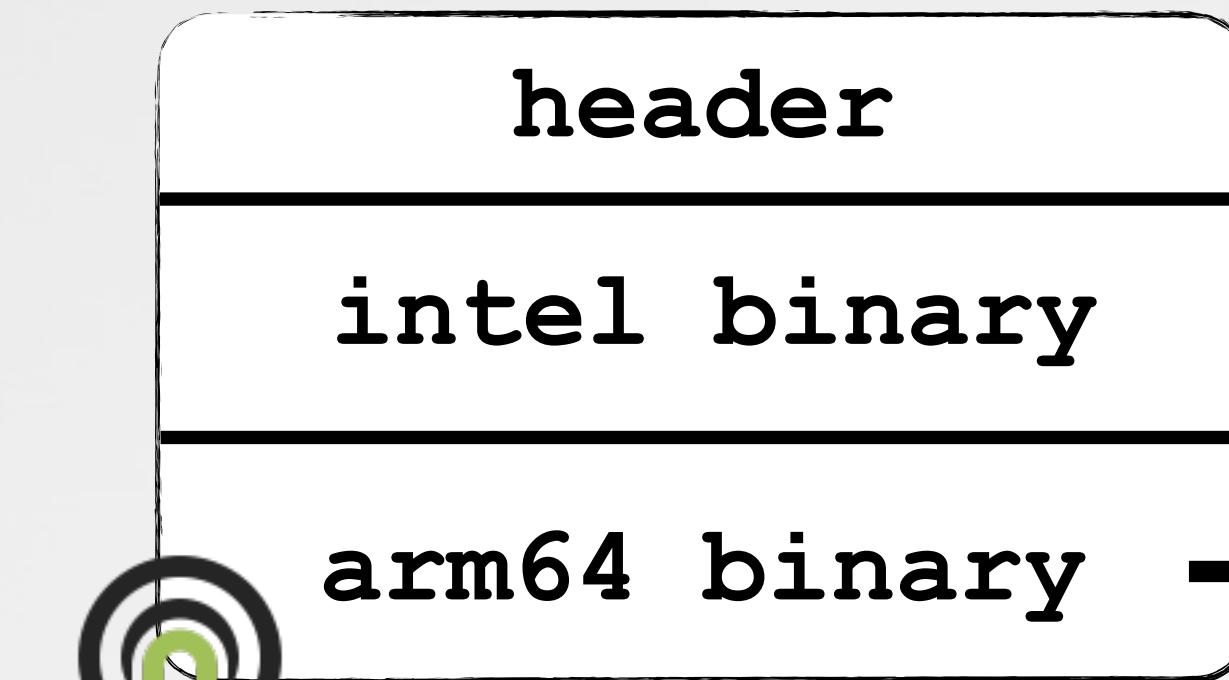
arm64/arm64e code  
(may be found in  
universal binary)

What's arm64e?



arm64 enhanced  
+pointer auth, etc.

Universal binary

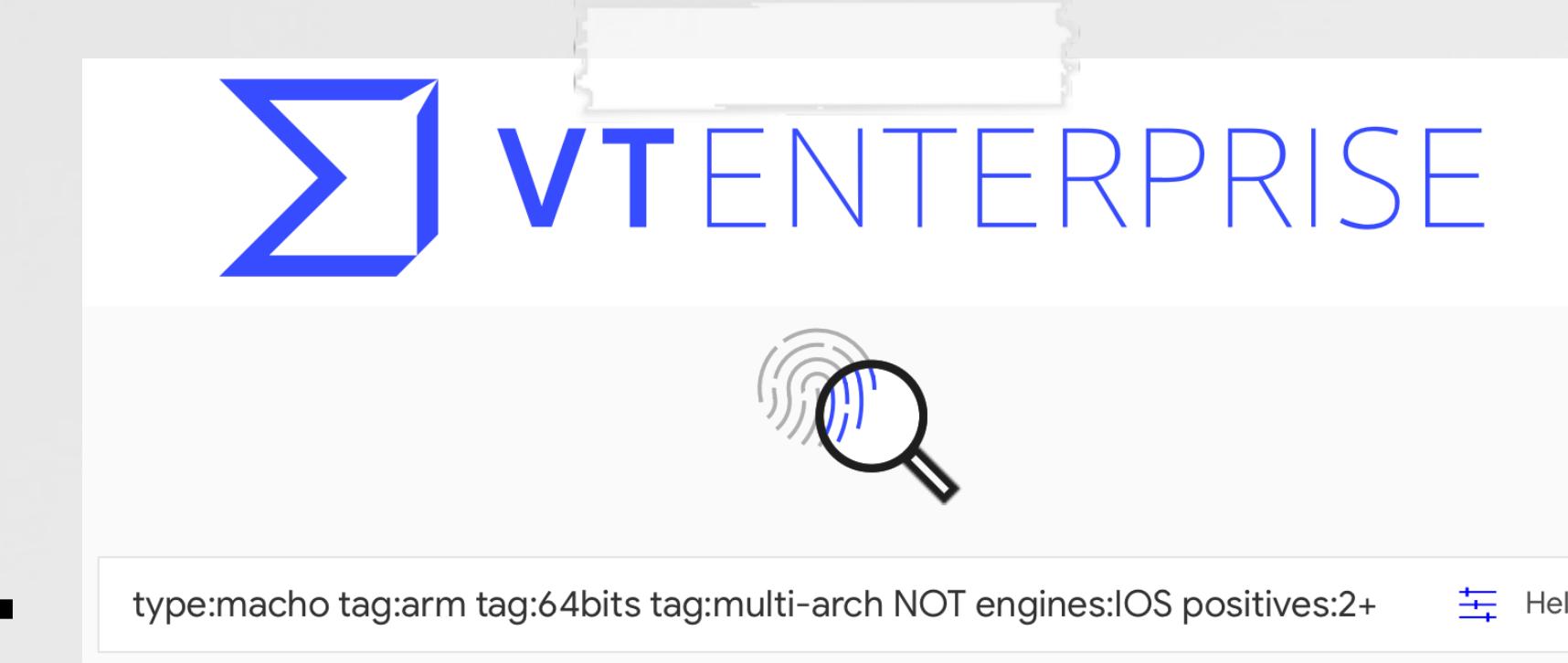


```
% otool -lv Calculator.app/Contents/MacOS/Calculator
...
Load command 10
    cmd LC_BUILD_VERSION
    cmdsize 32
    platform MACOS
    minos 11.4
    sdk 11.4
```

...built for macOS  
(also: LC\_VERSION\_MIN\_MACOSX)

# HUNTING FOR M1 MALWARE

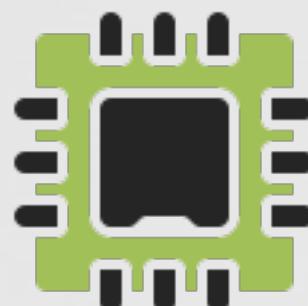
## querying virustotal for specimens



-----  
↓  
**type:macho tag:arm tag:64bits tag:multi-arch NOT engines:IOS positives:2+**



**tag:macho**  
**apple executable**



**tag:arm**  
**contains arm code**



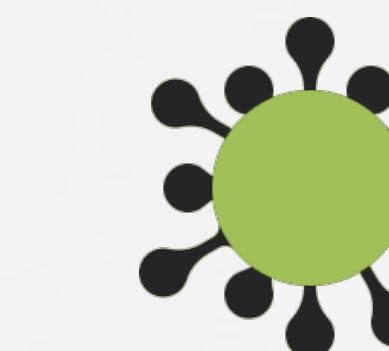
**tag:64bits**  
**contains 64bit code**



**tag: multi-arch**  
**universal binary**



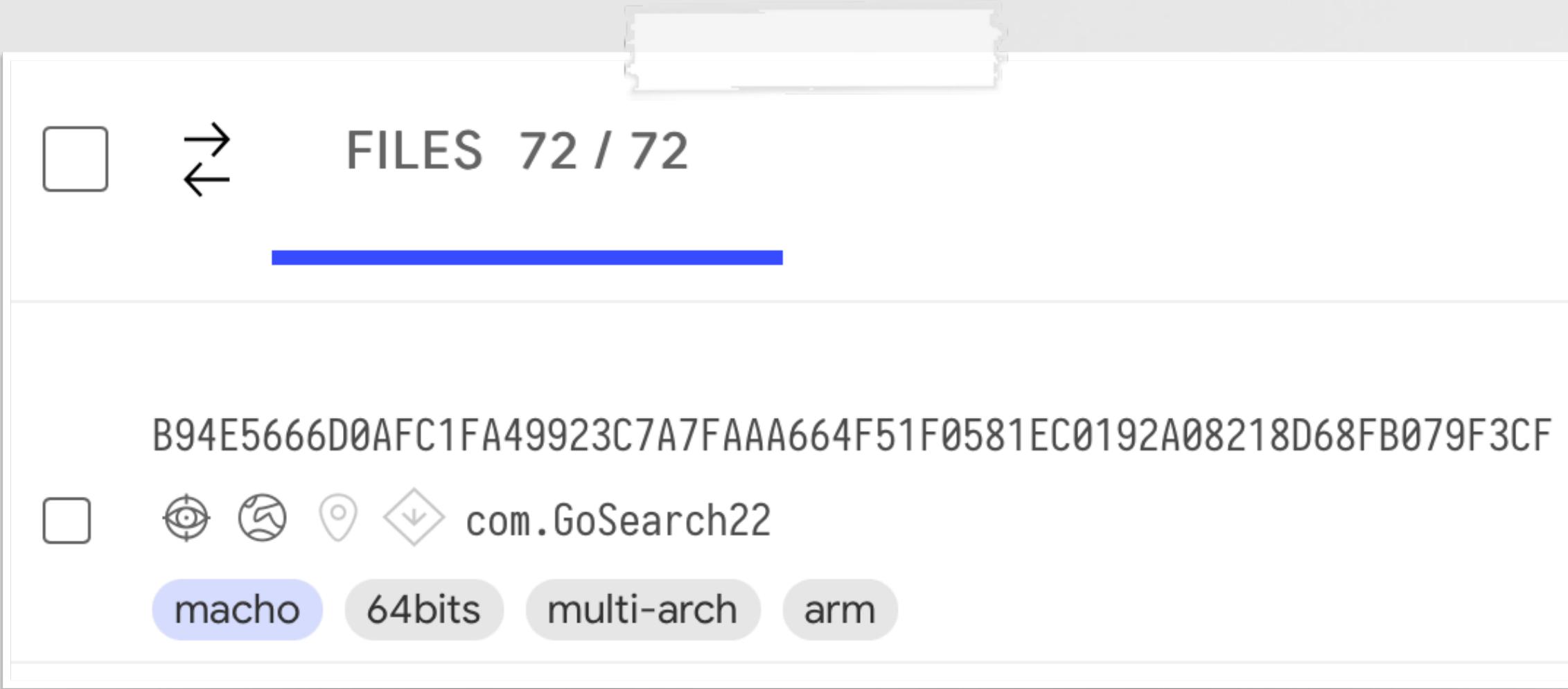
**NOT engines:IOS**  
**not an iOS binary**



**positives:2+**  
**flagged by 2+ AV engines**

# TRIAGING GoSEARCH22

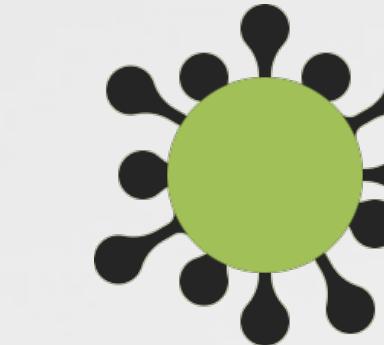
## a candidate (M1) binary



```
% file GoSearch22
Mach-O universal binary with 2 architectures:
 [arm64:Mach-O 64-bit executable arm64]
 [x86_64:Mach-O 64-bit executable x86_64]

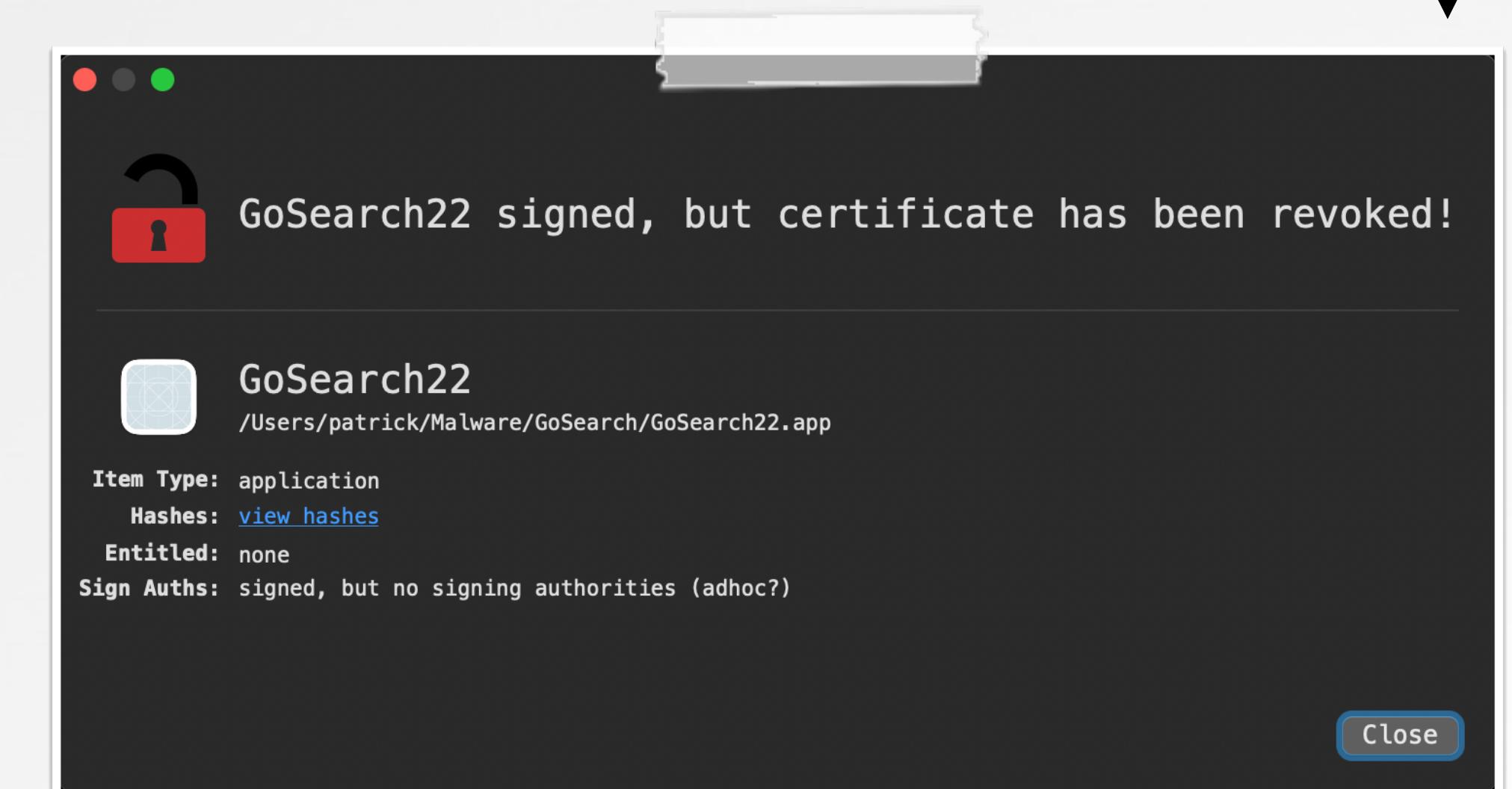
% otool -lv GoSearch22
...
Load command 9
    cmd LC_VERSION_MIN_MACOSX
    version 10.12
```

Universal macOS binary (with arm64)



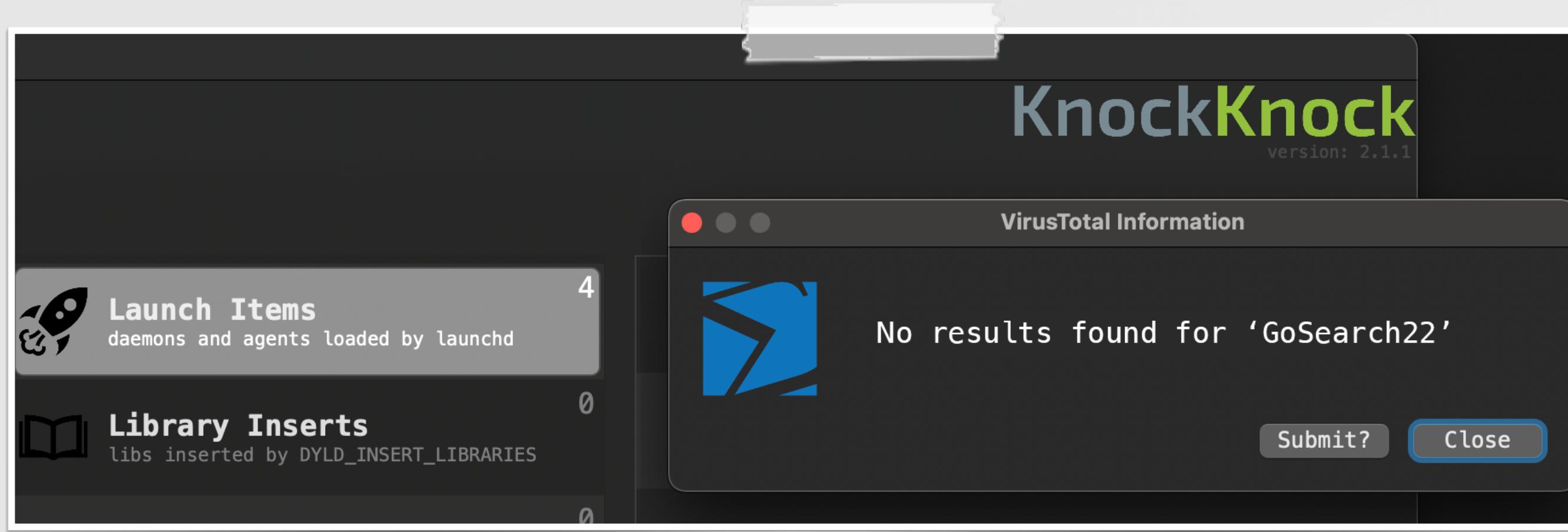
Flagged by several  
AV engines (intel code?)

+ app's cert. revoked -



Certificate revoked  
(by Apple)

# HOW DID IT END UP ON VIRUSTOTAL? ...detected and submitted via KnockKnock!



**KnockKnock detection  
(free: objective-see.com)**

A dashed arrow points from the 'VirusTotal Information' window in the first screenshot down to the 'Submissions' section of the second screenshot.

**Submissions**

Date	Name	Source
2020-12-27 23:37:51	GoSearch22	63b1639b - api

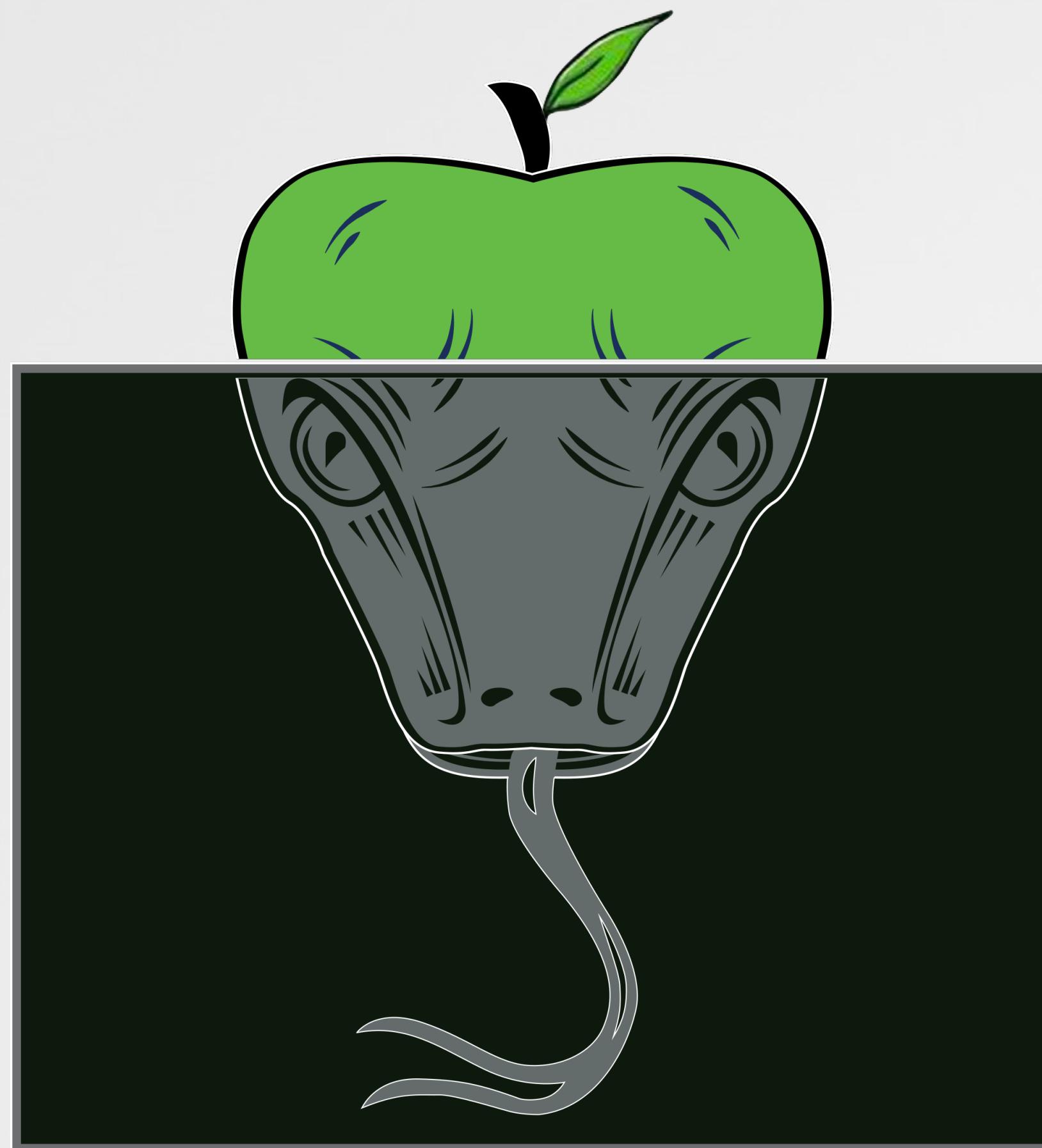
**Autostart Locations**

In-the-wild end-user machine registry keys and autostart locations where this file has been seen.

LaunchItems  
↳ GoSearch22

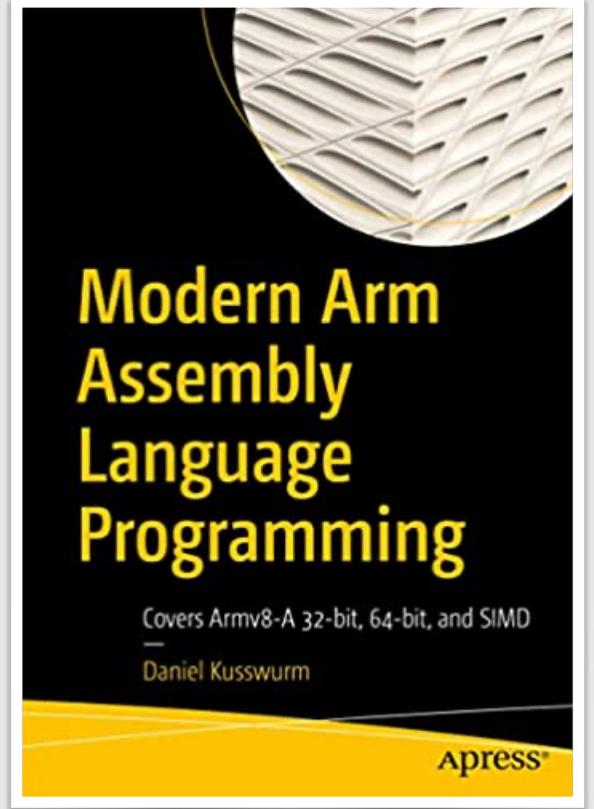
*via API  
(via KnockKnock)*

# Understanding arm64



# A BRIEF INTRODUCTION TO ARM64

## first, some most excellent resources



**"Modern Arm Assembly Language Programming"** (Daniel Kusswurm)

free, online

**"arm64 Assembly Crash Course"**

[github.com/Siguza/ios-resources/blob/master/bits/arm64.md](https://github.com/Siguza/ios-resources/blob/master/bits/arm64.md)



**"How to Read ARM64 Assembly Language"**

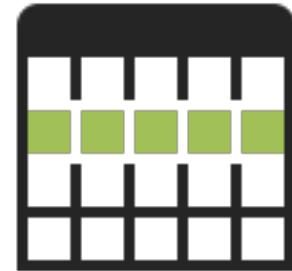
[wolchok.org/posts/how-to-read-arm64-assembly-language/](http://wolchok.org/posts/how-to-read-arm64-assembly-language/)

**"Introduction To Arm Assembly Basics"**

[azeria-labs.com/writing-arm-assembly-part-1/](http://azeria-labs.com/writing-arm-assembly-part-1/)

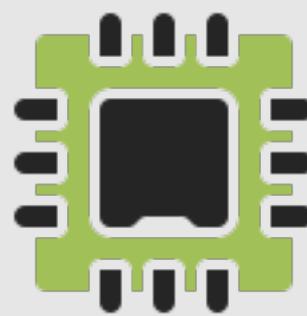
# REGISTERS and their uses

(somewhat) synonymous to variables in your fav. programming language

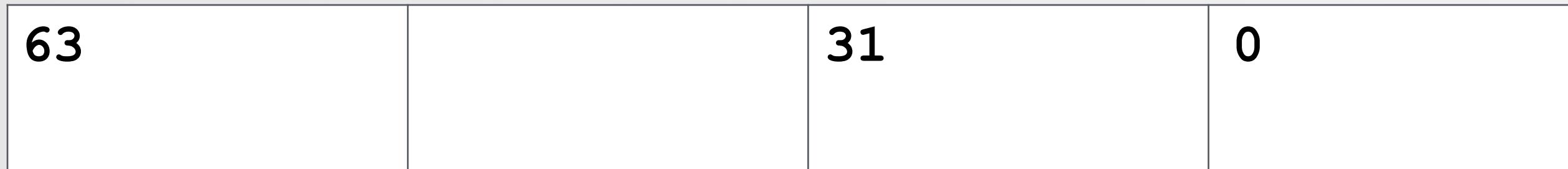


Registers: temporary storage "slots" on the CPU that can be referenced by name.

arm64



31 64-bit registers: x0 - x30  
w\* (e.g. w0)



sp: stack pointer

pc: program counter

xzr: virtual register, value: 0

PSTATE  
(processor state)

N	Z	C	V	...
---	---	---	---	-----

→ Condition flags

N: negative

Z: zero

C: carry

V: overflow

# REGISTERS

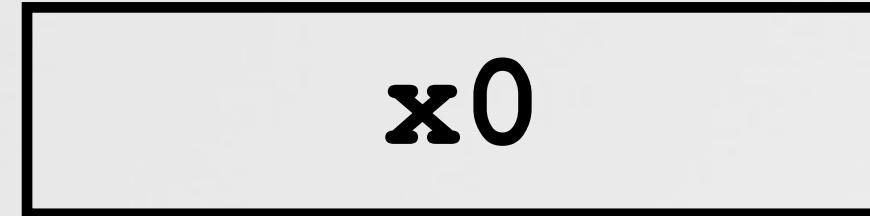
## usage, during a function call



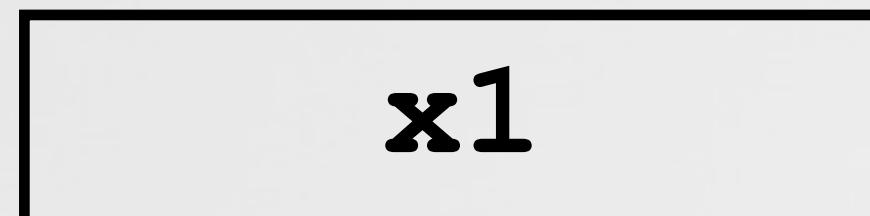
During analysis, we largely focus on api calls and their arguments.

Arguments:

arg 0 - ->

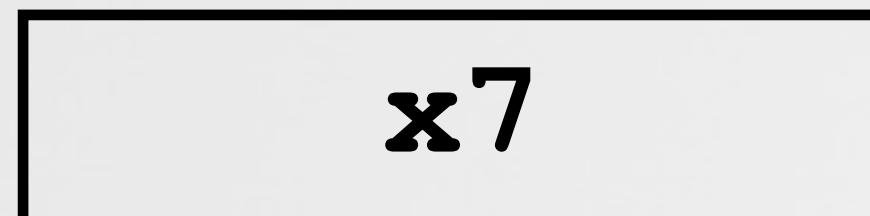


arg 1 - ->



...

arg 7 - ->



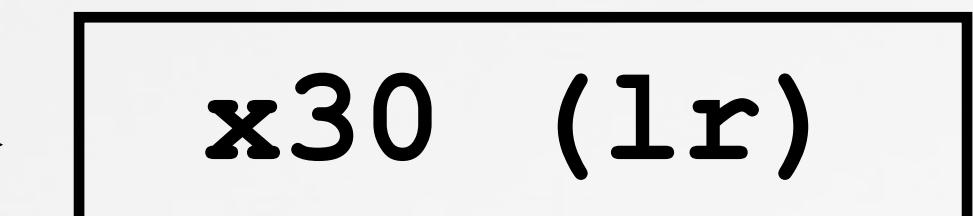
Frame pointer

- ->

x29 (fp)

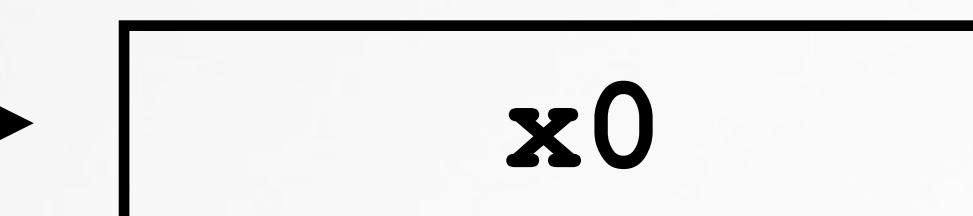
Return address:

- ->



Return value:  
(64/128 bits)

- ->

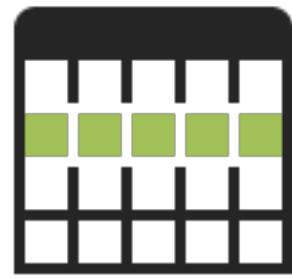


- ->



# INSTRUCTIONS

instruct the cpu what to do



Instructions: map to a specific sequence of bytes that instructs the CPU to perform an operation.

add	x1	x0	42
-----	----	----	----

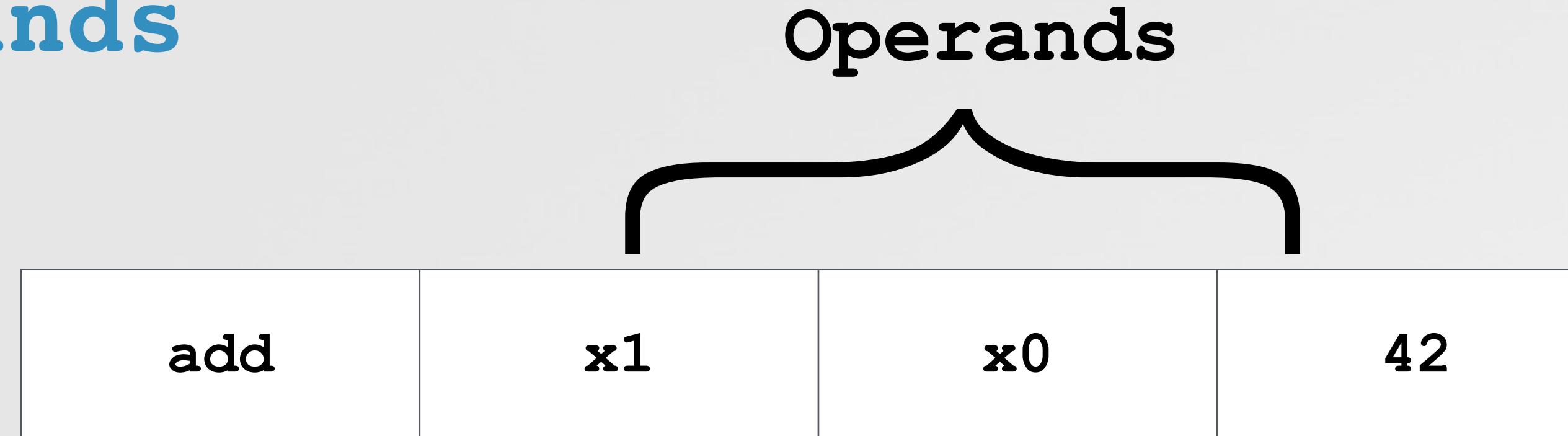
in C:  $x1 = x0 + 42;$

:  
:->

Mnemonic:

a (human-readable) abbreviation of the operation that the instructions perform.

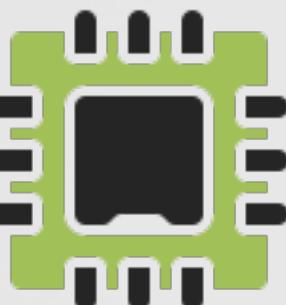
# INSTRUCTIONS the operands



Operand types:



**Immediate:**  
a constant value (e.g. 42)



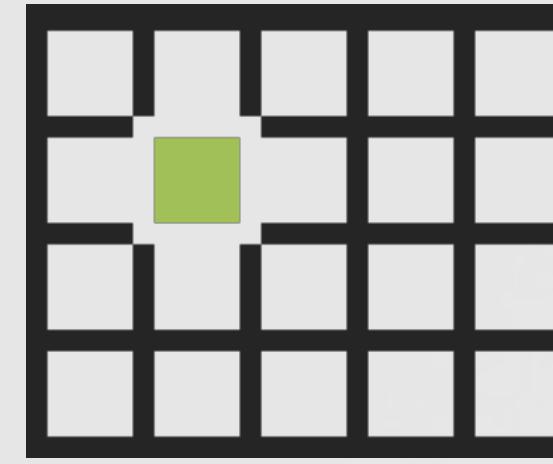
**Register:**  
a cpu register (e.g. x0, x1)



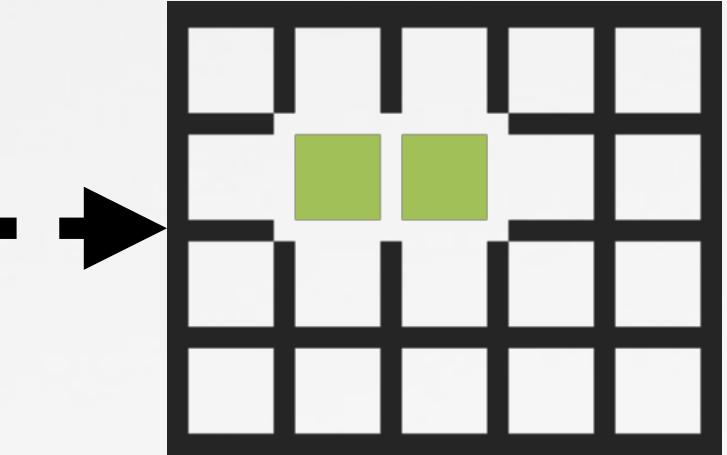
**Memory:**  
a cpu register, that points to a value in memory

# MEMORY ACCESS MODEL

arm's model is a "load & store"



- 1 Load (into register)
- 2 Perform any operation(s)
- 3 Store (into memory)



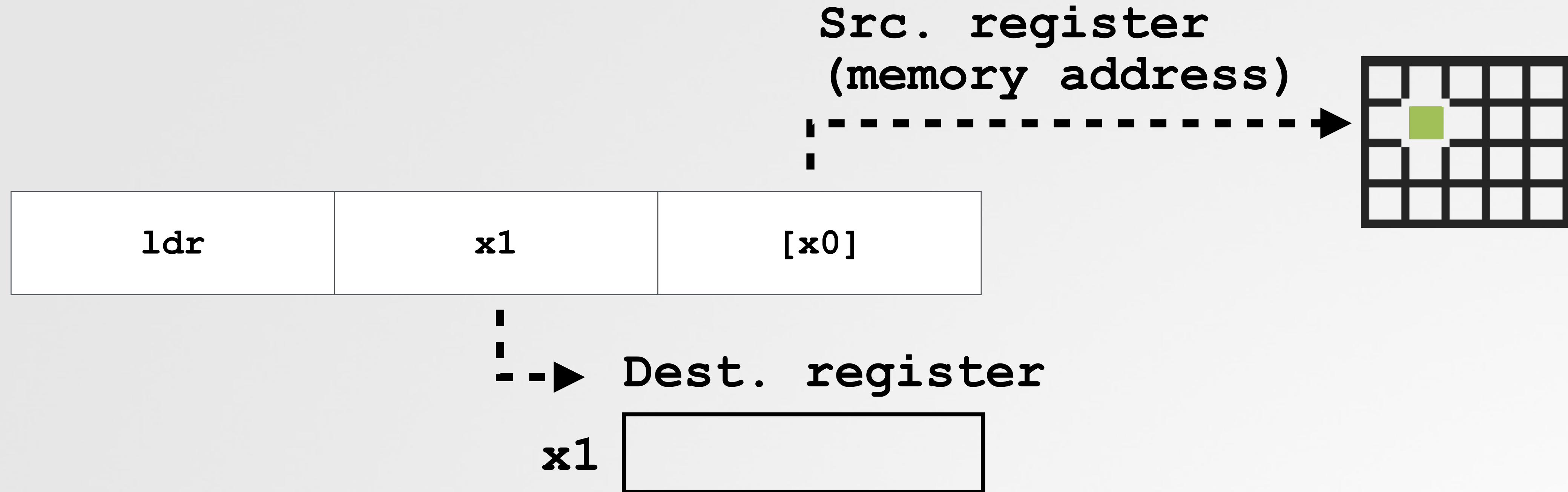
*"ARM uses a load-store model for memory access which means that only load/store (LDR and STR) instructions can access memory.*

*...on ARM data must be moved from memory into registers before being operated on" -Maria Markstedter (Azeria Labs)*

...a few other variants, ldp/stp

# MEMORY ACCESS MODEL

load via the `ldr` instruction (+ variants)

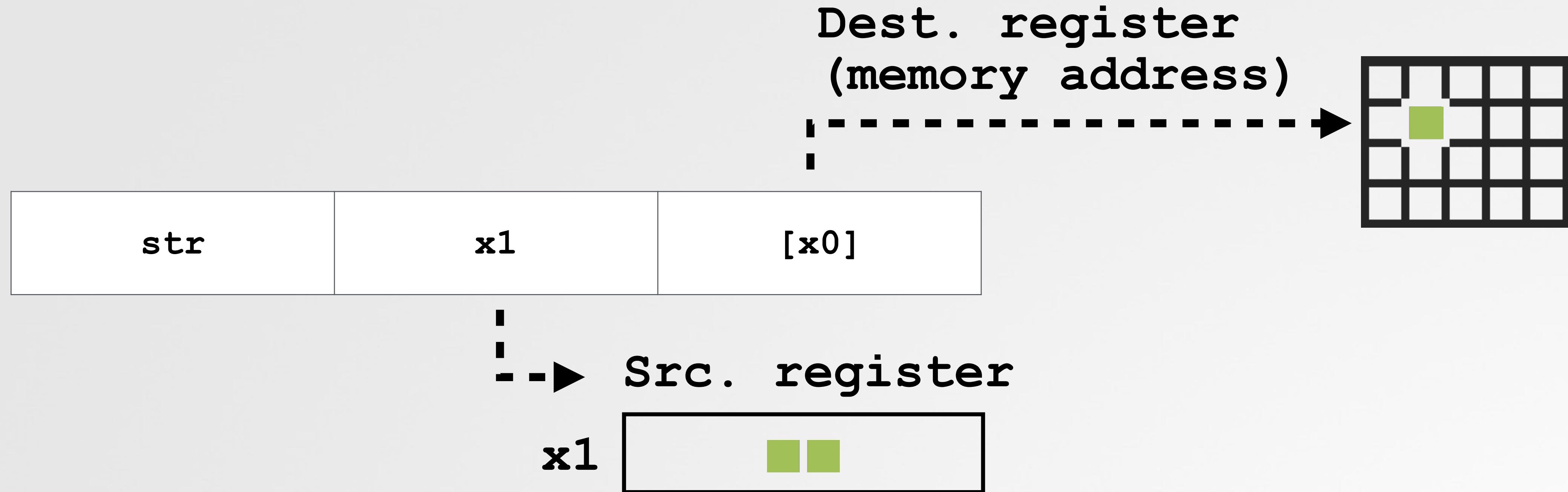


Analogous statement (in C):

`x1 = *x0;`

# MEMORY ACCESS MODEL

store via the `str` instruction (+ variants)



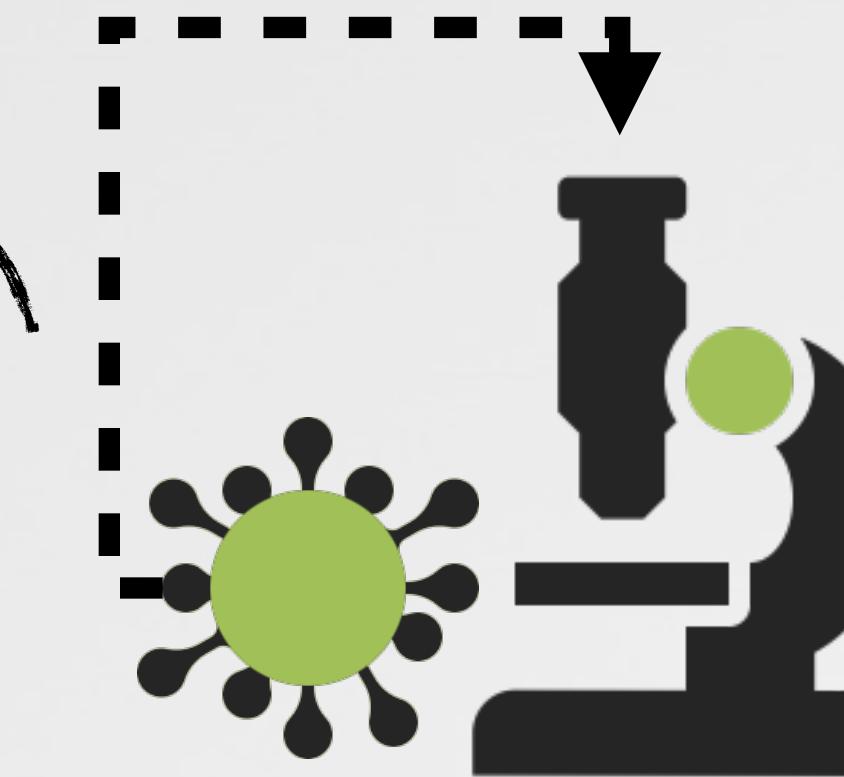
Analogous statement (in C):

`*x0 = x1;`

# CONDITIONS

set via `cmp`, etc...

`if( isDebugged )  
 exit`



cmp	x0	42
-----	----	----

;  
--> (discarded) subtract  
updates PSTATE flags

e.g. x0 is 42?

z flag is set

PSTATE  
(processor state)

N	Z	C	V	...
---	---	---	---	-----

--> Condition flags

# CONDITIONS

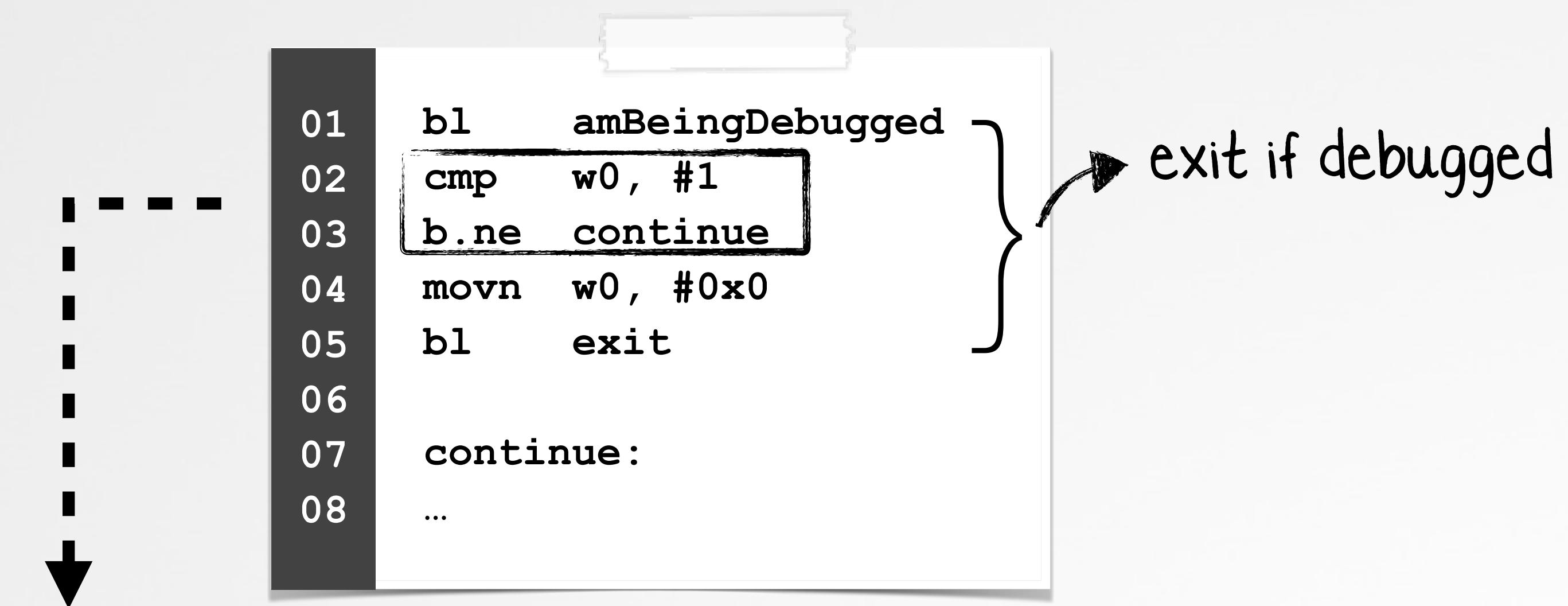
## condition codes



Once (condition) flags have been set, subsequent instructions can act upon them using condition codes

Name	Meaning
EQ	equal
NE	not equal
GE	greater or equal
GT	greater than
LE	lesser or equal
LT	less than
...	

Condition codes



b.ne

label

Branch (jump), if z not set

# BRANCHES

alter control flow of a program

1

b/br	imm/register
------	--------------

Branch (unconditionally)

2

b.cond	imm
--------	-----

Branch (if condition met)

3

bl/blr	imm/register
--------	--------------

Branch (store address of next instruction in x30 (lr))

e.g. function call

```
01 bl amBeingDebugged  
02 cmp w0, #1  
03 b.ne continue  
04 movn w0, #0x0  
05 bl exit  
06  
07 continue:  
08 ...
```

```
01 bl amBeingDebugged  
02 cmp w0, #1  
03 ...
```

&next instruction  
(cmp w0, #1)

ret
-----

Branch back to x30 (lr)

x30 (lr)

# REVERSING "HELLO, WORLD!"

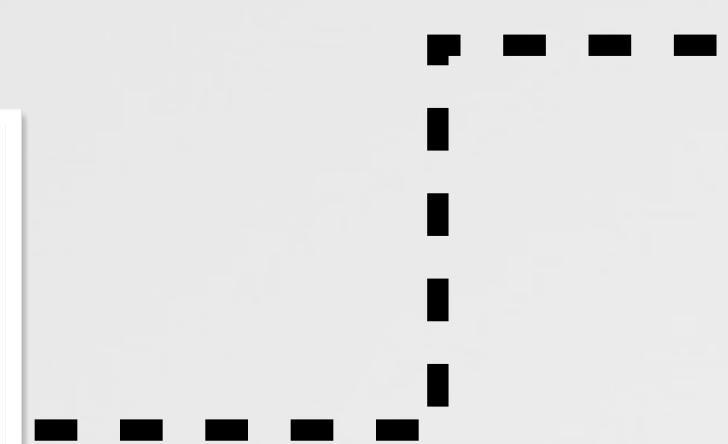
macOS arm64 version

```
01 int main(int argc, char * argv[]) {  
02     @autoreleasepool {  
03         NSLog(@"Hello, World!");  
04     }  
05     return 0;  
06 }
```

(Apple's) "Hello, World!"

Note, @autoreleasepool:

- 1 `objc_autoreleasePoolPush`
- +
- 2 `objc_autoreleasePoolPop`



```
01 main:  
02     sub    sp, sp, #0x30  
03     stp    x29, x30, [sp, #0x20]  
04     add    x29, sp, #0x20  
05     movz   w8, #0x0  
06     stur   wzr, [x29, #-0x4]  
07     stur   w0, [x29, #-0x8]  
08     str    x1, [sp, #0x10]  
09     str    w8, [sp, #0xc]  
10    bl     objc_autoreleasePoolPush  
11    adrp   x9, #0x0000000100004000  
12    add    x9, x9, #0x8 ; @"Hello, World!"  
13    str    x0, [sp]  
14    mov    x0, x9  
15    bl     NSLog  
16    ldr    x0, [sp]  
17    bl     objc_autoreleasePoolPop  
18    ldr    w0, [sp, #0xc]  
19    ldp    x29, x30, [sp, #0x20]  
20    add    sp, sp, #0x30  
21    ret
```

"Hello, World!"  
disassembled

# REVERSING "HELLO, WORLD!"

## function prologue

```
01 main:  
02     sub    sp, sp, #0x30      ----->  
03     stp    x29, x30, [sp, #0x20] ----->  
04     add    x29, sp, #0x20      ----->  
05     movz   w8, #0x0          |  
06     stur   wzr, [x29, #-0x4]  } ----->  
07     stur   w0, [x29, #-0x8]  }----->  
08     str    x1, [sp, #0x10]    |----->  
09     str    w8, [sp, #0xc]     |----->  
10    bl     objc_autoreleasePoolPush  
11    adrp   x9, #0x0000000100004000  
12    add    x9, x9, #0x8 ; @"Hello, World!"  
13    str    x0, [sp]  
14    mov    x0, x9  
15    bl     NSLog  
16    ldr    x0, [sp]  
17    bl     objc_autoreleasePoolPop  
18    ldr    w0, [sp, #0xc]  
19    ldp    x29, x30, [sp, #0x20]  
20    add    sp, sp, #0x30  
21    ret
```

-----> Subtract 0x30 from stack pointer  
-----> Store x29 & x30 at SP + 0x20  
-----> Set frame pointer (x29) to sp + 0x20  
-----> Save registers/init local variables

	Offset	Value
sp	----->	0x30
	0x28	x30
x29	----->	0x20
	...	x29
sp	----->	0x00



Function prologue makes space on the stack for saving registers, local variables, and init's frame pointer

# REVERSING "HELLO, WORLD!"

## invoking `objc_autoreleasePoolPush`

```
01 main:  
02     sub    sp, sp, #0x30  
03     stp    x29, x30, [sp, #0x20]  
04     add    x29, sp, #0x20  
05     movz   w8, #0x0  
06     stur   wzr, [x29, #-0x4]  
07     stur   w0, [x29, #-0x8]  
08     str    x1, [sp, #0x10]  
09     str    w8, [sp, #0xc]  
10     bl     objc_autoreleasePoolPush -  
11     adrp   x9, #0x0000000100004000  
12     add    x9, x9, #0x8 ; @"Hello, World!"  
13     str    x0, [sp] -  
14     mov    x0, x9  
15     bl     NSLog  
16     ldr    x0, [sp]  
17     bl     objc_autoreleasePoolPop  
18     ldr    w0, [sp, #0xc]  
19     ldp    x29, x30, [sp, #0x20]  
20     add    sp, sp, #0x30  
21     ret
```

Branch to (call)  
`objc_autoreleasePoolPush`

address of next instruction,  
stored in link register (x30)

Return value (x0: pool object)  
saved to local variable



`objc_autoreleasePoolPush` takes no arguments  
...returns a pointer to a pool object (in x0).

# REVERSING "HELLO, WORLD!" invoking NSLog with the "Hello, World!" string

```
01 main:
02 sub    sp, sp, #0x30
03 stp    x29, x30, [sp, #0x20]
04 add    x29, sp, #0x20
05 movz   w8, #0x0
06 stur   wzr, [x29, #-0x4]
07 stur   w0, [x29, #-0x8]
08 str    x1, [sp, #0x10]
09 str    w8, [sp, #0xc]
10 bl     objc_autoreleasePoolPush
11 adrp   x9, #0x0000000100004000
12 add    x9, x9, #0x8 ; @"Hello, World!
13 str    x0, [sp]
14 mov    x0, x9
15 bl     NSLog
16 ldr    x0, [sp]
17 bl     objc_autoreleasePoolPop
18 ldr    w0, [sp, #0xc]
19 ldp    x29, x30, [sp, #0x20]
20 add    sp, sp, #0x30
21 ret
```

- - - → Initialize address to "Hello World!" (string) object
- - - → Initialize 1st argument with address of string object
- - - → Branch to (call) NSLog function



Here, NSLog is invoked with a single argument, the address of the string object to print (passed in x0).

# REVERSING "HELLO, WORLD!"

## invoking `objc_autoreleasePoolPop` with pool object

```
01 main:  
02     sub    sp, sp, #0x30  
03     stp    x29, x30, [sp, #0x20]  
04     add    x29, sp, #0x20  
05     movz   w8, #0x0  
06     stur   wzr, [x29, #-0x4]  
07     stur   w0, [x29, #-0x8]  
08     str    x1, [sp, #0x10]  
09     str    w8, [sp, #0xc]  
10     bl     objc_autoreleasePoolPush  
11     adrp   x9, #0x0000000100004000  
12     add    x9, x9, #0x8 ; @"Hello, World!"  
13     str    x0, [sp]  
14     mov    x0, x9  
15     bl     NSLog  
16     ldr    x0, [sp] - - - - -  
17     bl     objc_autoreleasePoolPop - - - - -  
18     ldr    w0, [sp, #0xc]  
19     ldp    x29, x30, [sp, #0x20]  
20     add    sp, sp, #0x30  
21     ret
```

Initialize 1st argument with  
address pool object (previous  
stored on the stack)

Branch to (call)  
`objc_autoreleasePoolPop` function



`objc_autoreleasePoolPop` takes a single argument, the  
address of the pool object to release (passed in `x0`).

# REVERSING "HELLO, WORLD!"

## function epilogue

```
01 main:  
02     sub    sp, sp, #0x30  
03     stp    x29, x30, [sp, #0x20]  
04     add    x29, sp, #0x20  
05     movz   w8, #0x0  
06     stur   wzr, [x29, #-0x4]  
07     stur   w0, [x29, #-0x8]  
08     str    x1, [sp, #0x10]  
09     str    w8, [sp, #0xc]  
10     bl     objc_autoreleasePoolPush  
11     adrp   x9, #0x0000000100004000  
12     add    x9, x9, #0x8 ; @"Hello, World!"  
13     str    x0, [sp]  
14     mov    x0, x9  
15     bl     NSLog  
16     ldr    x0, [sp]  
17     bl     objc_autoreleasePoolPop  
18     ldr    w0, [sp, #0xc] -----'  
19     ldp    x29, x30, [sp, #0x20] -----'  
20     add    sp, sp, #0x30 -----'  
21     ret    -----'
```

Initialize return value  
(previously set to zero)

----- → Restore x29/x30 registers

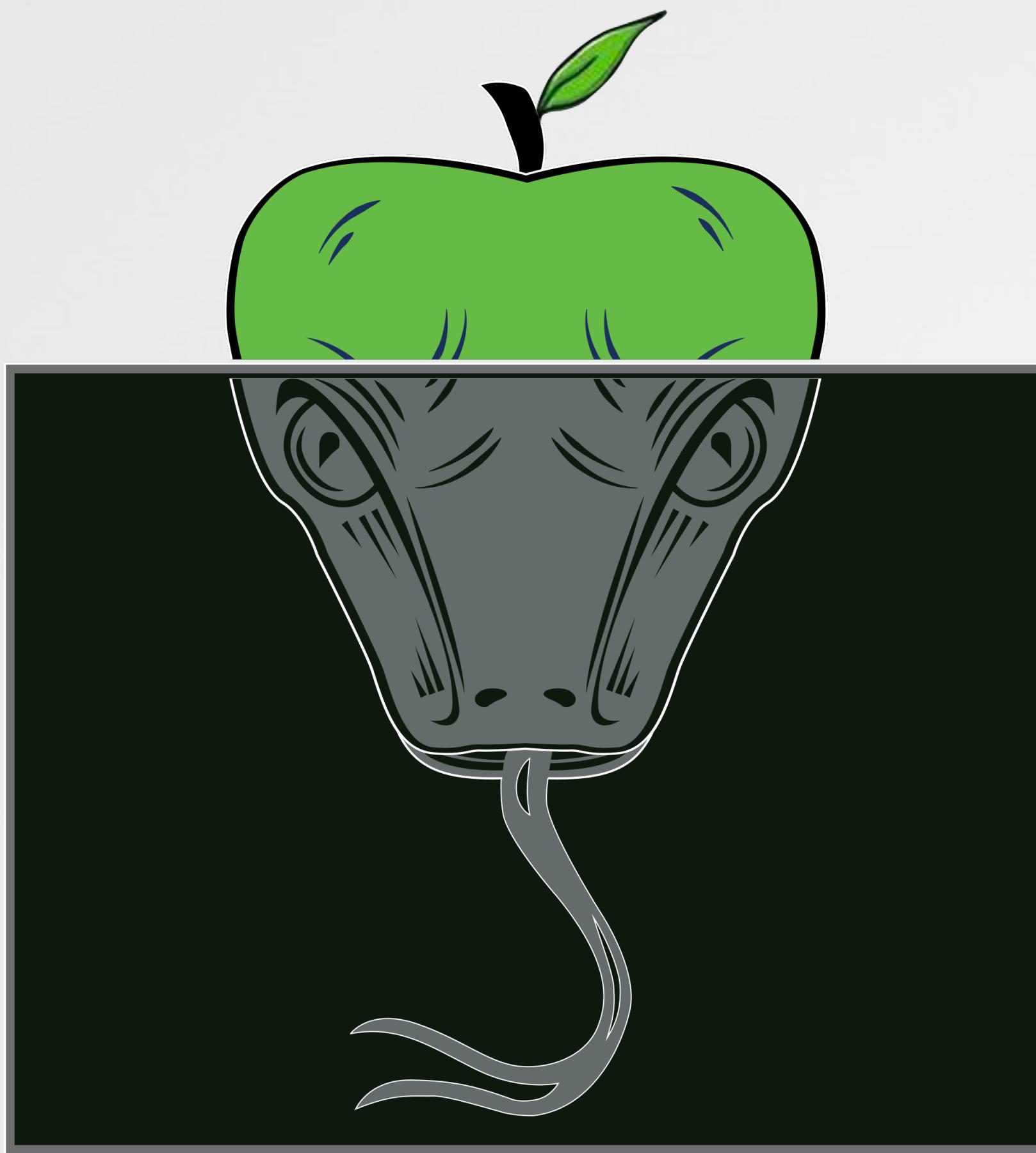
----- (re)Adjust stack

----- ↑ Return to caller (lr/x30)



A function epilogue restores saved registers, and (re)adjusts the stack. Return, branches to lr (x30).

# Practical M1 Malware Analysis



# A FUNDAMENTAL UNDERSTANDING OF ARM64 SUFFICE? ...often, yes!

```
01 main:  
02     sub    sp, sp, #0x30  
03     stp    x29, x30, [sp, #0x20]  
04     add    x29, sp, #0x20  
05     movz   w8, #0x0  
06     stur   wzr, [x29, #-0x4]  
07     stur   w0, [x29, #-0x8]  
08     str    x1, [sp, #0x10]  
09     str    w8, [sp, #0xc]  
10     bl     objc_autoreleasePoolPush  
11     adrp   x9, #0x0000000100004000  
12     add    x9, x9, #0x8 ; @"Hello, World!"  
13     str    x0, [sp]  
14     mov    x0, x9  
15     bl     NSLog  
16     ldr    x0, [sp]  
17     bl     objc_autoreleasePoolPop  
18     ldr    w0, [sp, #0xc]  
19     ldp    x29, x30, [sp, #0x20]  
20     add    sp, sp, #0x30  
21     ret
```



```
01 int main(int arg0, int arg1) {  
02     var_20 =objc_autoreleasePoolPush();  
03     NSLog(@"Hello, World!");  
04     objc_autoreleasePoolPop(var_20);  
05     return 0x0;  
06 }
```

"Hello, World!" decompiled

## Dynamic Analysis Tools:



File monitor



Process monitor

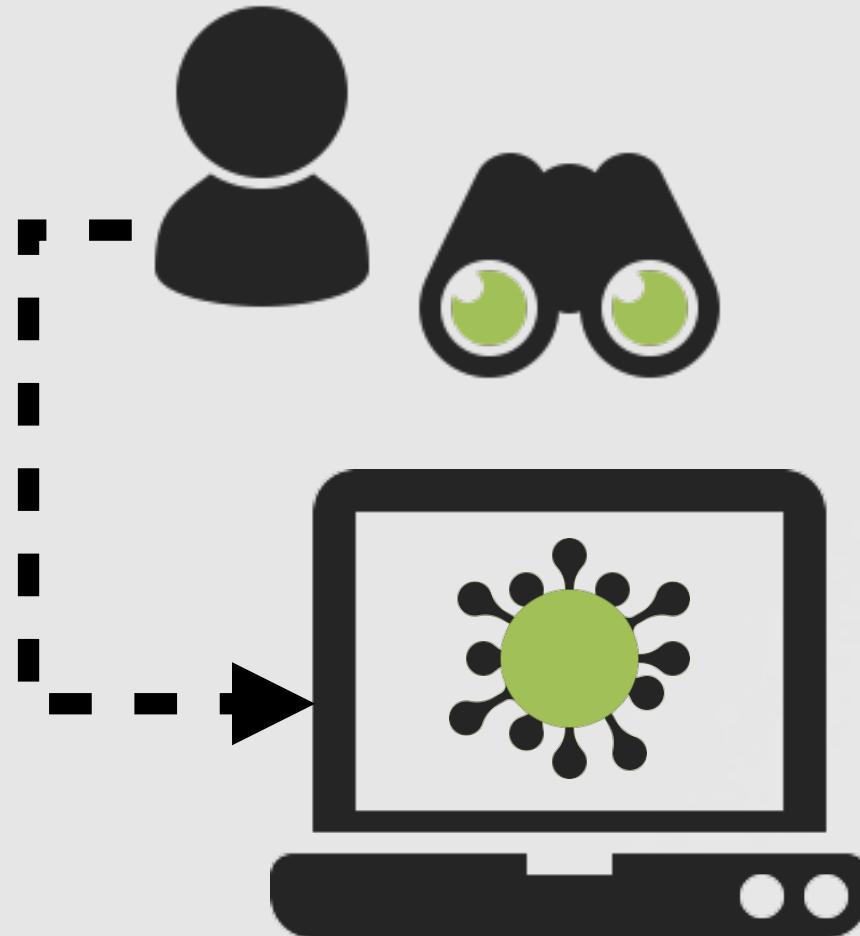


Network monitor



By leveraging a decompiler and dynamic analysis tools,  
often a fundamental understanding of arm64 will suffice!

# DYNAMIC ANALYSIS TOOLS may (trivially) reveal malware's capabilities



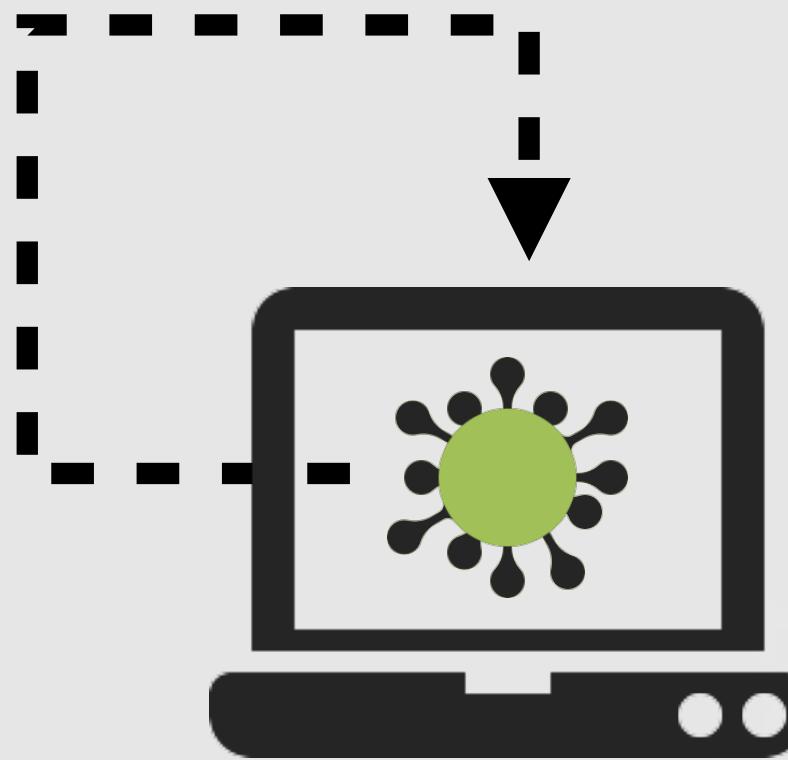
Analysis  
machine

```
# FileMonitor.app/Contents/MacOS/FileMonitor -pretty
{
    "event" : "ES_EVENT_TYPE_NOTIFY_CREATE",
    "file" : {
        "destination" : "/Users/user/Library/LaunchAgents/mdworker.plist",    ↗ launch agent persistence
        "process" : {
            "uid" : 501,
            "arguments" : [
                "/bin/sh",
                "-c",
                "/Users/user/Desktop/eTrader.app/Contents/Utils/mdworker"
            ],
            "path" : "/Users/user/Desktop/eTrader.app/Contents/Utils/mdworker",
            "name" : "mdworker"
        }
    }
}
```

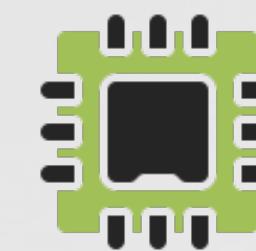
Uncovering persistence  
(via a file monitor)

# ANTI-ANALYSIS LOGIC

## aim to thwart (dynamic) analysis environments/tools



### Introspection



Am I being debugged?



Am I in a virtual machine?

simply terminates :(

```
% lldb GoSearch22.app
(lldb) target create "GoSearch22.app"

(lldb) c
Process 654 resuming
Process 654 exited with status = 45 (0x0000002d)
```

### GoSearch22 vs. debugger



One must identify and bypass anti-analysis mechanisms before comprehensive analysis of a malicious sample can commence!

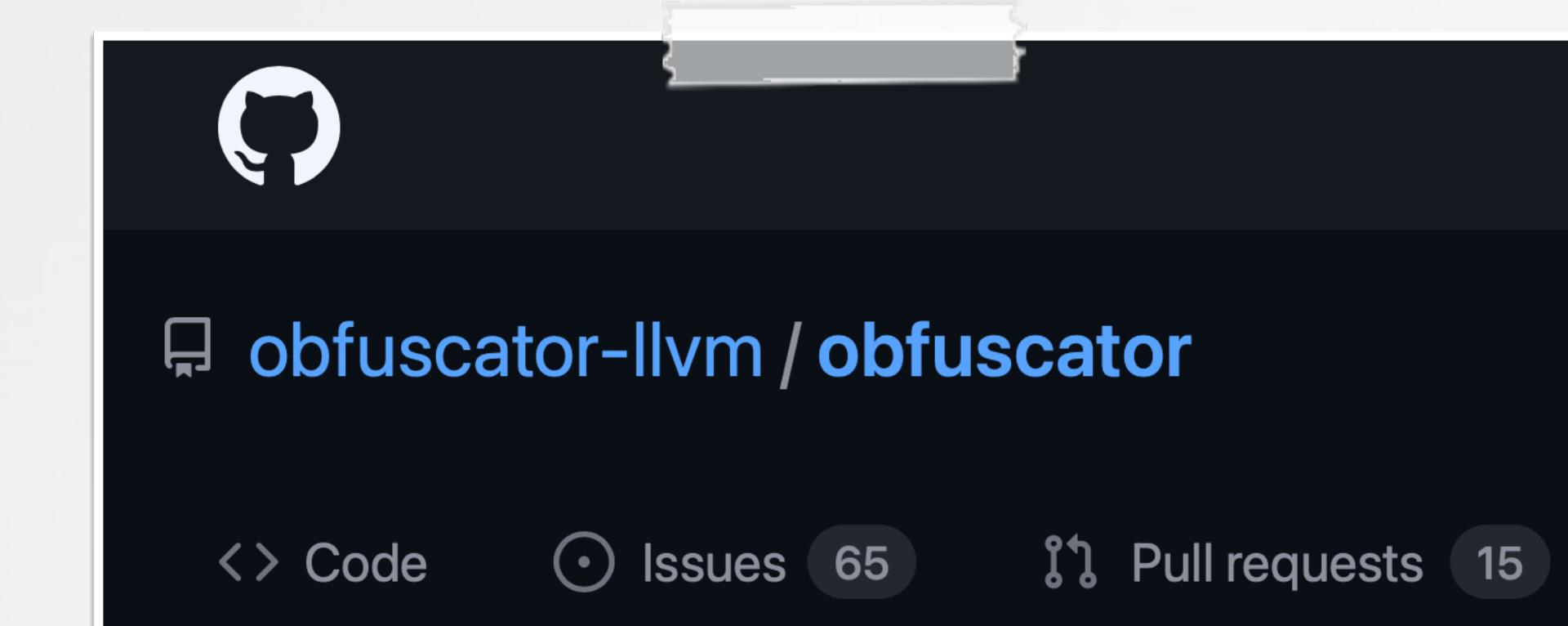
# GoSEARCH22

## ...also contains static analysis obfuscations

```
*(int128_t *)0x100071a30 = *(int128_t *)0x100071a30;
*(int8_t *)0x100071a40 = *(int8_t *)0x100071a40 ^ 0xcd;
*(int128_t *)0x100071a50 = *(int128_t *)0x100071a50;
*(int8_t *)0x100071a60 = *(int8_t *)0x100071a60 ^ 0xa;
*(int8_t *)0x100071a61 = *(int8_t *)0x100071a61 ^ 0x3a;
*(int8_t *)0x100071a62 = *(int8_t *)0x100071a62 ^ 0xaf;
*(int8_t *)0x100071a63 = *(int8_t *)0x100071a63 ^ 0x48;
*(int128_t *)0x100071a70 = *(int128_t *)0x100071a70;
*(int8_t *)0x100071a80 = *(int8_t *)0x100071a80 ^ 0xc9;
*(int128_t *)0x100071a90 = *(int128_t *)0x100071a90;
*(int8_t *)0x100071aa0 = *(int8_t *)0x100071aa0 ^ 0xfffffffffffffb;
*(int8_t *)0x100071aa1 = *(int8_t *)0x100071aa1 ^ 0xea;
*(int8_t *)0x100071aa2 = *(int8_t *)0x100071aa2 ^ 0xd5;
*(int8_t *)0x100071aa3 = *(int8_t *)0x100071aa3 ^ 0x51;
*(int128_t *)0x100071ab0 = *(int128_t *)0x100071ab0;
*(int8_t *)0x100071ac0 = *(int8_t *)0x100071ac0 ^ 0xd4;
*(int128_t *)0x100071ad0 = *(int128_t *)0x100071ad0;
*(int8_t *)0x100071ae0 = *(int8_t *)0x100071ae0 ^ 0xed;
*(int8_t *)0x100071ae1 = *(int8_t *)0x100071ae1 ^ 0xf2;
*(int8_t *)0x100071ae2 = *(int8_t *)0x100071ae2 ^ 0xa;
*(int8_t *)0x100071ae3 = *(int8_t *)0x100071ae3 ^ 0xffffffffffff8f;
v0 = v0 ^ v1 ^ v1 ^ v1 ^ v1 ^ v1 ^ v1 ^ v1;
```

```
var_70 = var_78;
var_98 = dlsym(dlopen(0x0, 0xa), 0x100071b40);
dlsym(dlopen(0x0, 0xa), 0x100071b28);
dlsym(dlopen(0x0, 0xa), 0x100071b10);
dlsym(dlopen(0x0, 0xa), 0x100071af0);
dlsym(dlopen(0x0, 0xa), 0x100071ad0);
dlsym(dlopen(0x0, 0xa), 0x100071ab0);
dlsym(dlopen(0x0, 0xa), 0x100071a90);
dlsym(dlopen(0x0, 0xa), 0x100071a70);
```

## Spurious function calls



## Garbage instructions?

## Popular obfuscator



See, "Using LLVM to Obfuscate Your Code During Compilation" ([www.apriorit.com](http://www.apriorit.com))

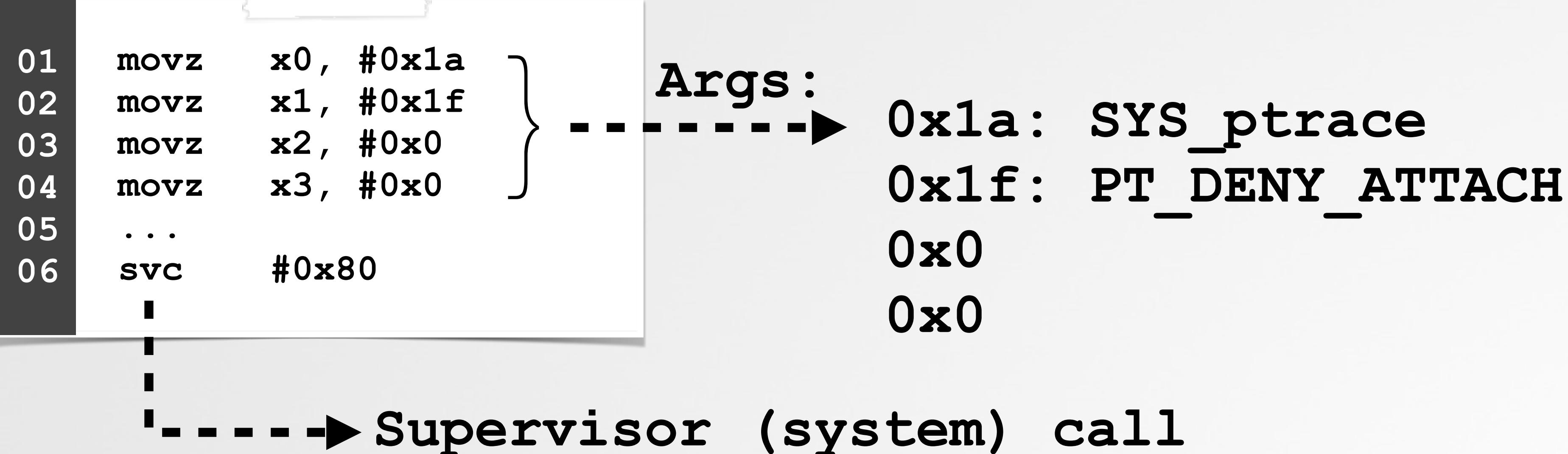
# GoSearch22's ANTI-ANALYSIS LOGIC

## debugger detection via ptrace/PT\_DENY\_ATTACH

```
% lldb GoSearch22.app  
Process 654 exited with status = 45 (0x0000002d)
```

45 (02xd)  
ENOTSUP (from PT\_DENY\_ATTACH)

### GoSearch22 vs. debugger



ptrace() + PT\_DENY\_ATTACH, prevents future attachments or terminates (with 45) if a debugger is currently attached.

# BYPASSING ANTI-ANALYSIS LOGIC

## once detected and identified, trivial to bypass

```
01 0x00000001000541f4 movz    x3, #0x0
02 0x00000001000541f8 movz    x16, #0x0
03 0x00000001000541fc svc    #0x80
04
05 0x0000000100054200 movz    w11, #0x6b8f
```

simply skip over  
ptrace system call :)

```
% lldb GoSearch22.app

(lldb) b 0x00000001000541fc
Breakpoint 1: address = 0x00000001000541fc

(lldb) Process 1486 stopped
* thread #1, queue = 'com.apple.main-thread'
  stop reason = breakpoint 1.1:
-> 0x00000001000541fc svc    #0x80  modify PC

(lldb) reg write $pc 0x100054200
```

Modify pc register

# GoSEARCH22's ANTI-ANALYSIS LOGIC

## system integrity protection (sip) status detection

```
01  ldr    x8, [sp, #0x190 + var_120] - - - , : -----> Two arguments  
02  ldr    x0, [sp, #0x190 + var_100] } - - - + -  
03  ldr    x1, [sp, #0x190 + var_F8] } - - -  
04  blr    x8  <-----| ...but what's the branch target?
```

```
% lldb GoSearch22.app  
...  
(lldb) x/i $pc  
-> 0x1000538dc: 0xd63f0100 blr x8  
  
(lldb) reg read $x8  
x8 = 0x0000000193a5f160 libobjc.A.dylib`objc_msgSend
```

Debugger introspection



As we've identified (and thwarted) the malware's anti-debugging logic, we can now fully leverage the debugger!

# GoSEARCH22's ANTI-ANALYSIS LOGIC

## system integrity protection (sip) status detection

### objc\_msgSend

Sends a message with a simple return value to an instance of a class.

#### Declaration

```
void objc_msgSend(void);
```

#### Parameters

##### self

A pointer that points to the instance of the class that is to receive the message.

##### op

The selector of the method that handles the message.

##### ...

A variable argument list containing the arguments to the method.

-----> Arg 0: self  
object method is invoked upon  
-----> Arg 1: op  
selector of method

```
% lldb GoSearch22.app
...
(lldb) po $x0
<NSConcreteTask: 0x1058306c0>

(lldb) x/s $x1
0x1e9fd4fae: "launch"
```

Debugger introspection  
[NSTask launch];

# GoSEARCH22's ANTI-ANALYSIS LOGIC

## SIP status detection

Class

### NSTask

An object representing a subprocess of the current process.

Instance Property

#### launchPath

Sets the receiver's executable.

Instance Property

#### arguments

Sets the command arguments that should be used to launch the executable.

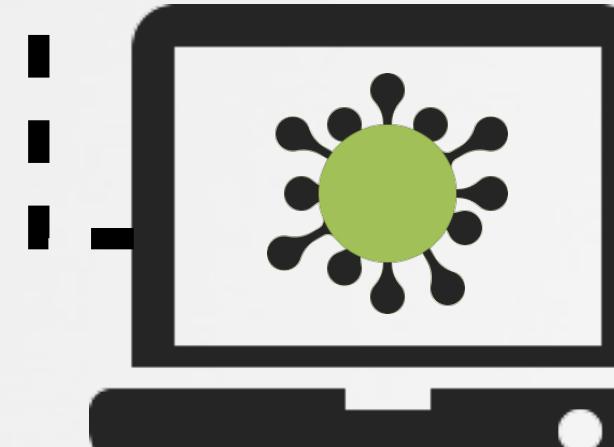
NSTask  
+ it's properties

```
(lldb) po [$x0 launchPath]  
/bin/sh  
  
(lldb) po [$x0 arguments]  
<__NSArrayI 0x10580dfd0>(  
-C,  
command -v csrutil > /dev/null && csrutil status |  
grep -v "enabled" > /dev/null && echo 1 || echo 0  
)
```

**SIP status detection  
(via "csrutil status")**

```
% csrutil status  
System Integrity Protection status: disabled.
```

**SIP: disabled?  
(malware exists!)**



analysis machine

# GoSearch22's Anti-Analysis Logic

## virtual machine detection

```
01 ldr    x8, [sp, #0x190 + var_120]  
02 ldr    x0, [sp, #0x190 + var_100]  
03 ldr    x1, [sp, #0x190 + var_F8]  
04 blr    x8-----:--> (another) call  
          to obj_msgSend
```

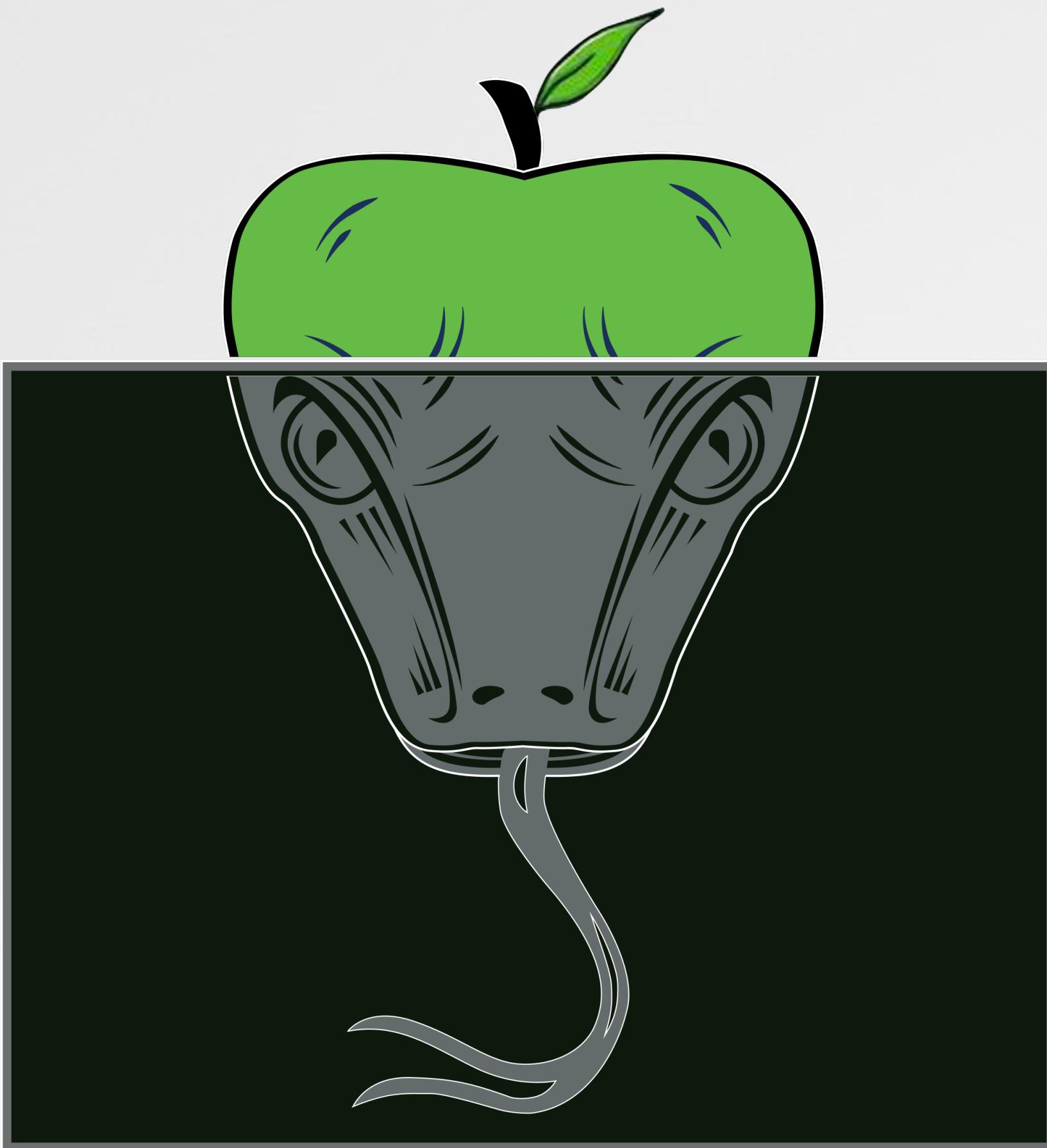
```
(lldb) po $x0  
<NSConcreteTask: 0x1058306c0>
```

```
(lldb) po [$x0 launchPath]  
/bin/sh
```

```
(lldb) po [$x0 arguments]  
<__NSArrayI 0x10580c1f0> (  
-c,  
readonly VM_LIST="VirtualBox\|Oracle\|VMware\|Parallels\|qemu";is_hwmodel_vm(){ ! sysctl -n hw.model|grep  
"Mac">/dev/null;};is_ram_vm(){((($($($sysctl -n hw.memsize)/ 1073741824))<4));};is_ped_vm(){ local -r ped=$  
(ioreg -rd1 -c IOPlatformExpertDevice);echo "${ped}"|grep -e "board-id" -e "product-name" -e "model"|grep  
-qi "${VM_LIST}";}|echo "${ped}"|grep "manufacturer"|grep -v "Apple">/dev/null;};is_vendor_name_vm(){ ioreg  
-l|grep -e "Manufacturer" -e "Vendor Name"|grep -qi "${VM_LIST}";};is_hw_data_vm(){ system_profiler  
SPHardwareDataType 2>&1 /dev/null|grep -e "Model Identifier"|grep -qi "${VM_LIST}";};is_vm()  
{ is_hwmodel_vm||is_ram_vm||is_ped_vm||is_vendor_name_vm||is_hw_data_vm; };main(){ is_vm&&echo 1||echo  
0;};main "${@}" )
```

looks for artifacts from  
various virtualization products

# Conclusions



# ... AND MORE ! notarization, infection numbers, etc...

ConfiantIntel  
@ConfiantIntel

OSX/Bundlore Loader found (0 detections in VT)  
compiled for ARM (targeting the new M1 MacBook!) , and  
notarized by Apple ... 🙌 It was delivered through  
malvertising. Downloads Bundlor and an unknown  
payload

Developer ID Application: Bobbie Miller (PX3WCCP368)

```
spctl -a
accepted
```

```
[test@ /Volumes/Install 3]: spctl -a -vv /Volumes/Install\ 3/Installer.app/
/Volumes/Install 3/Installer.app/: accepted
source=Notarized Developer ID
origin=Developer ID Application: Bobbie Miller (PX3WCCP368)
```

notarized by Apple

Mashable

Tech Apple

## New malware 'Silver Sparrow' is targeting both Intel and M1 Macs

Nearly 30,000 Macs (and counting?) have been infected.

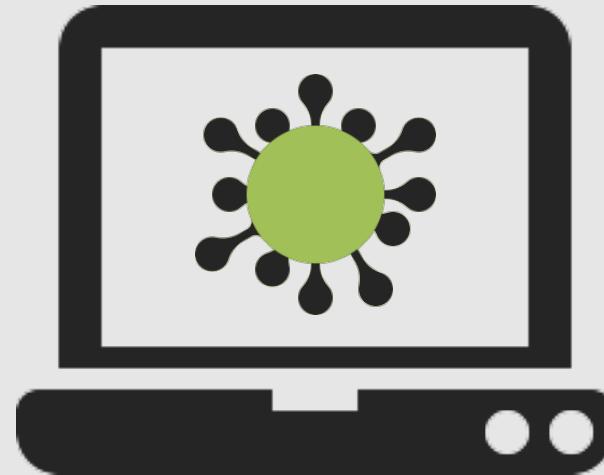
# OSX.SilverSparrow (30k+ infections!)

## OSX.Hydromac (notarized!)

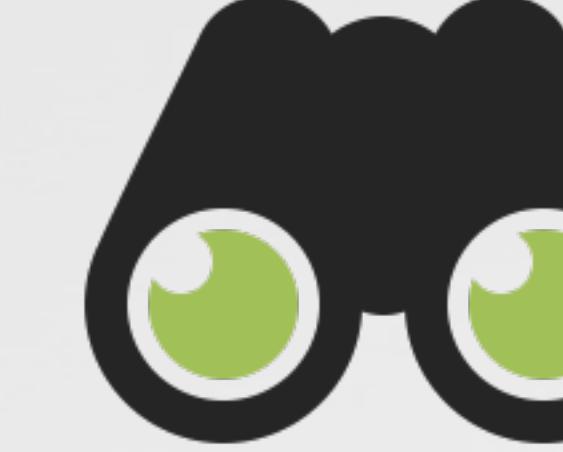
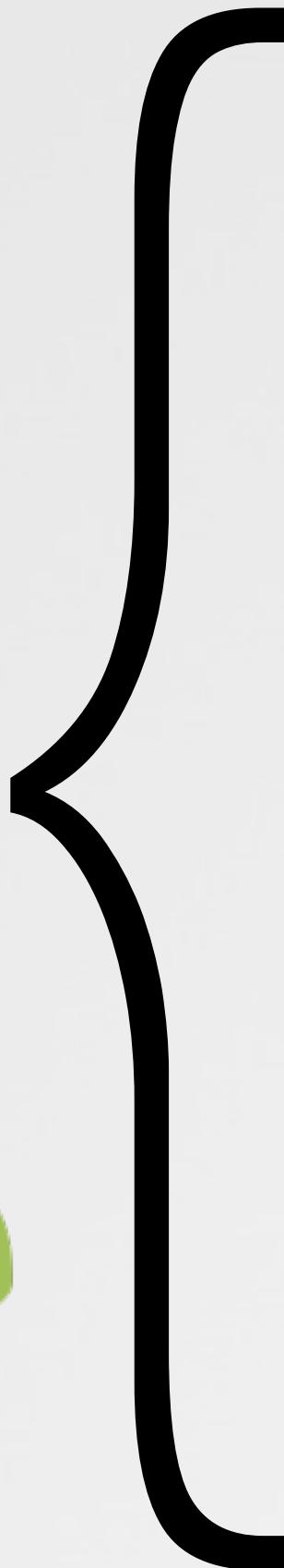


"OSX/Hydromac: New Mac adware, leaked from a flashcards app"  
(Taha Karim (@lordx64) [objective-see.com/blog/blog\\_0x65.html](http://objective-see.com/blog/blog_0x65.html))

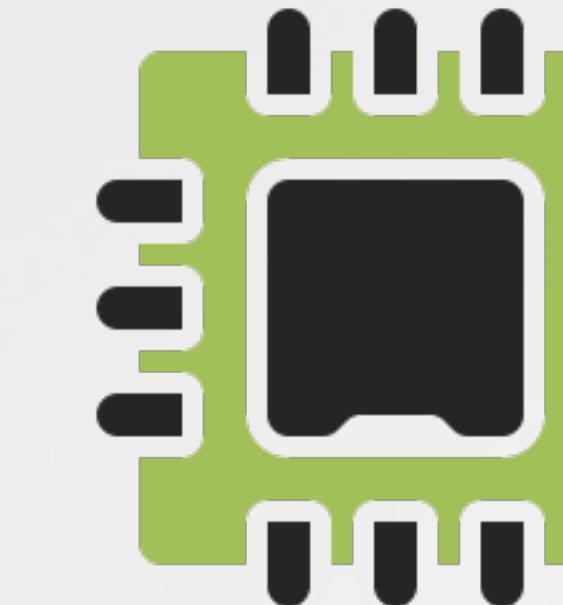
# KEY TAKEAWAYS



M1 malware  
is here to stay



Hunting for  
native M1 malware



Understanding arm64



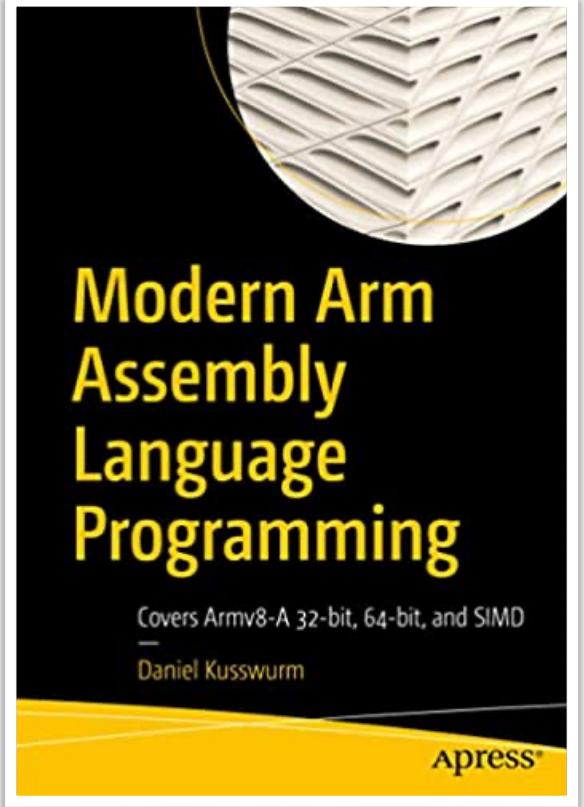
Practical M1  
malware analysis



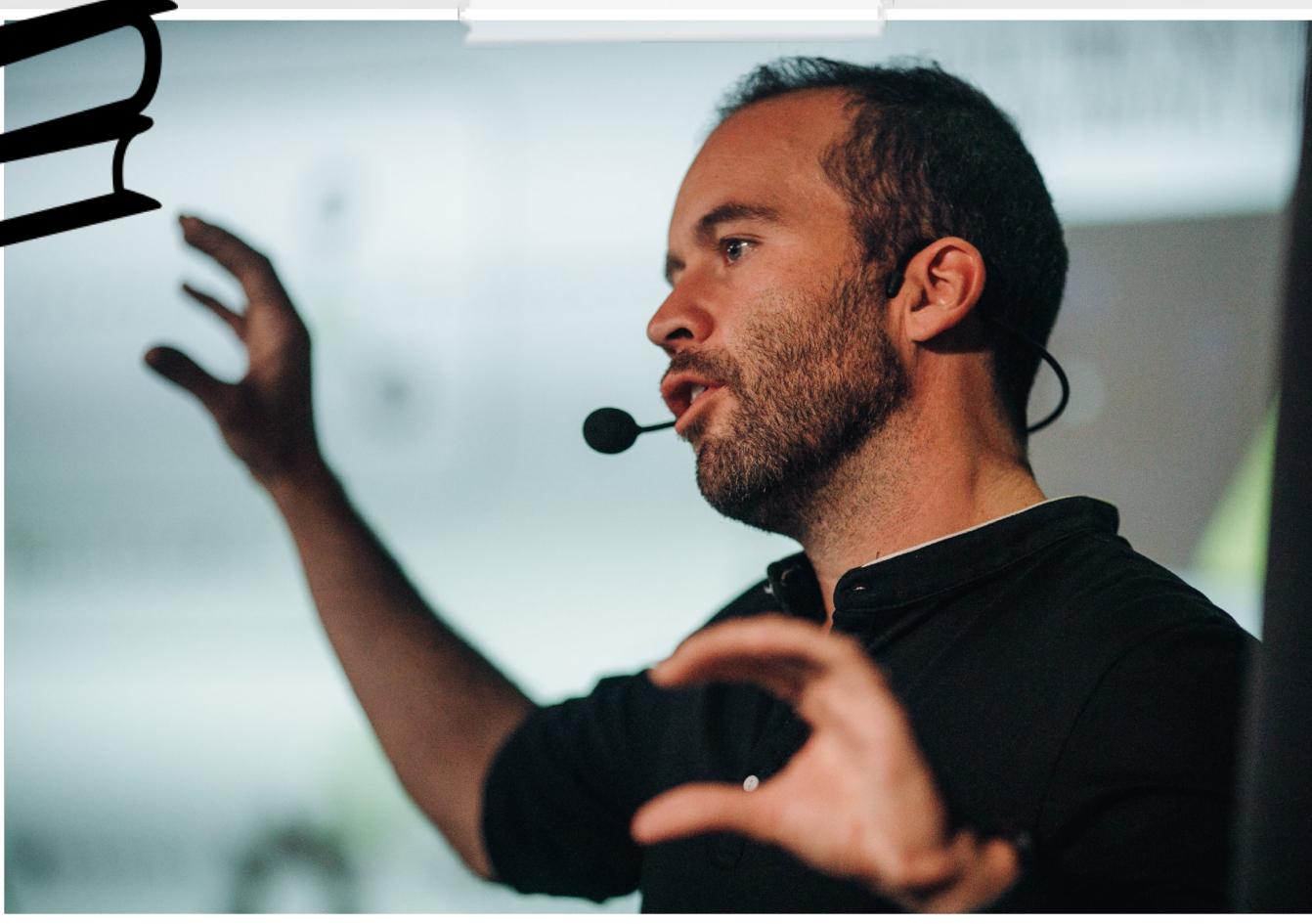
Armed with the topics presented here today, you're well on the way to becoming a proficient analyst of m1 malware!

# LEARN MORE?

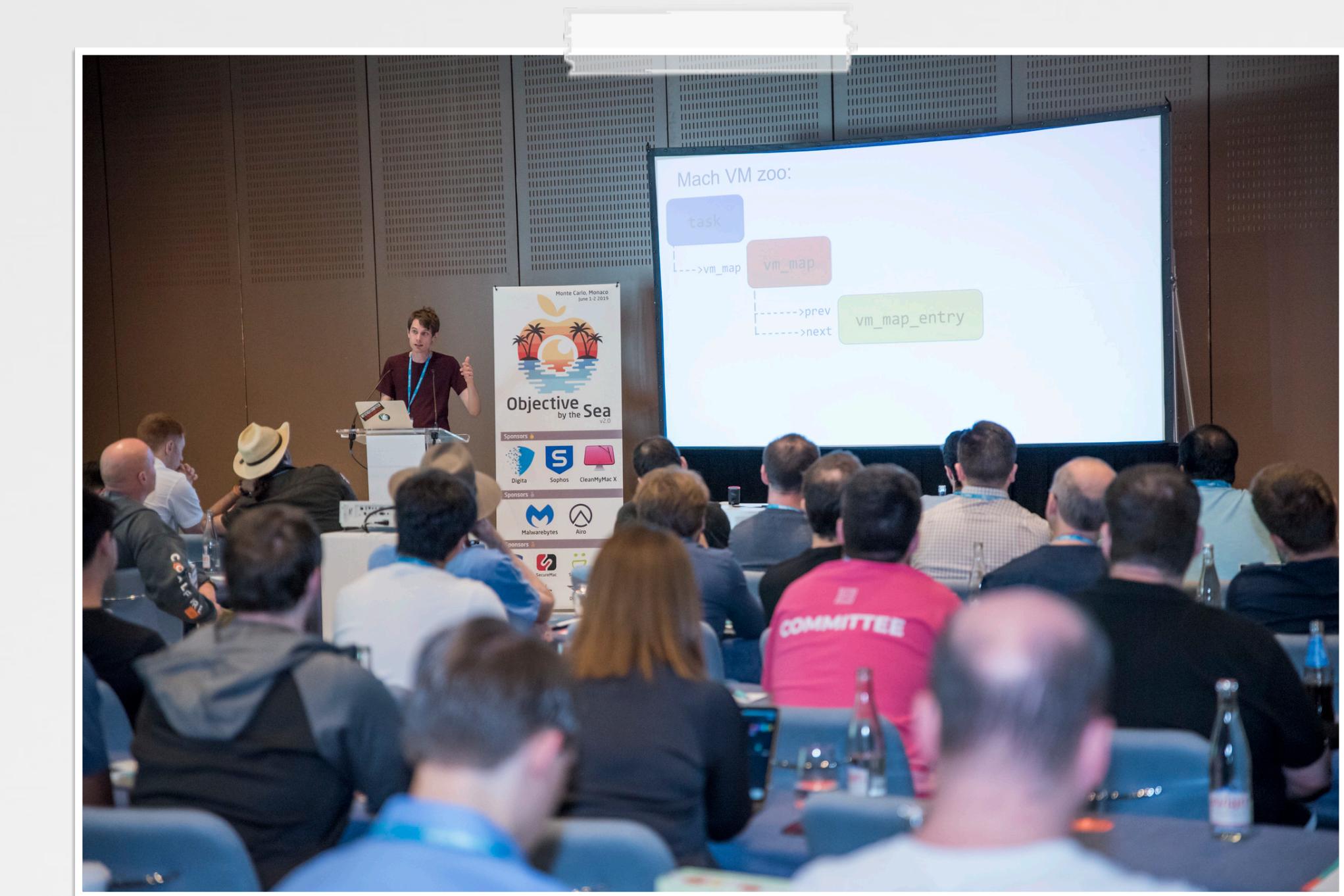
arm64, malware analysis, macOS security topics



## "Modern Arm Assembly Language Programming"



"The Art of Mac Malware"  
[taomm.org](http://taomm.org)



"Objective by the Sea"

Sept 30/Oct 1

Maui, Hawaii, USA

[ObjectiveByTheSea.com](http://ObjectiveByTheSea.com)

# MAHALO!

"Friends of Objective-See"



MOSYLE



SmugMug



Guardian Mobile Firewall



SecureMac



iVerify



Halo Privacy

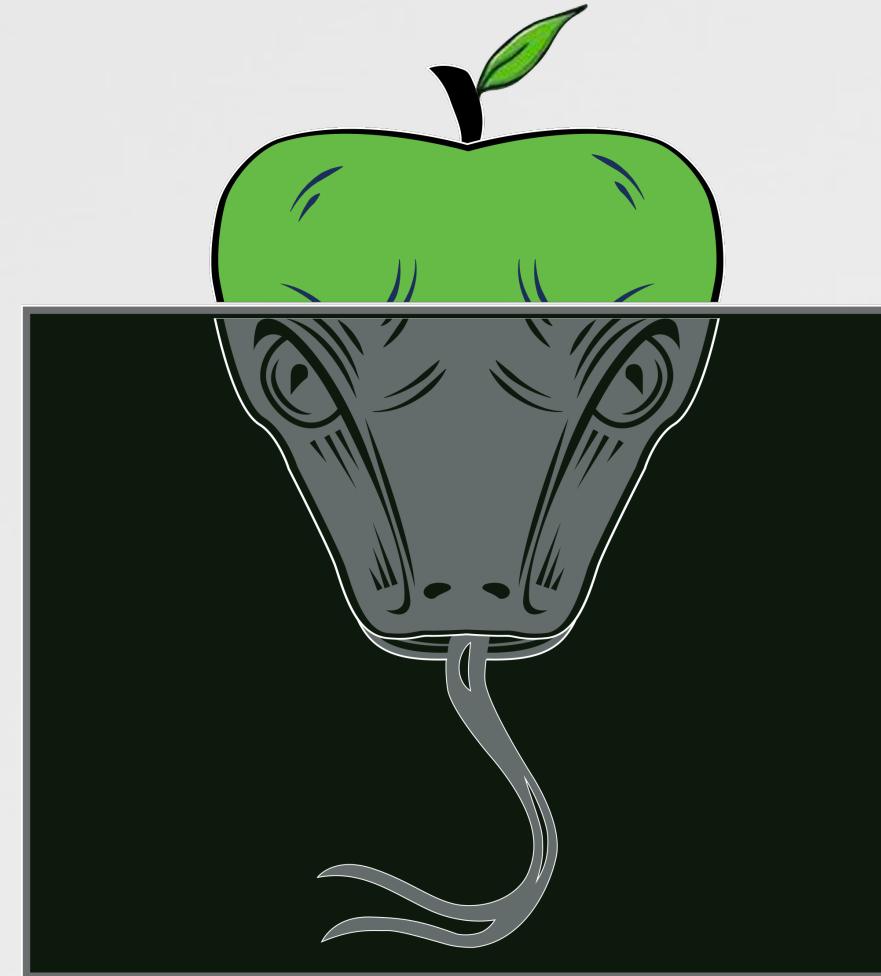


uberAgent



Grab the Slides:  
[speakerdeck.com/patrickwardle](https://speakerdeck.com/patrickwardle)

# Arm'd & Dangerous



## RESOURCES :

**"Modern Arm Assembly Language Programming"**  
[www.apress.com/gp/book/9781484262665](http://www.apress.com/gp/book/9781484262665)

**"arm64 Assembly Crash Course"**  
[github.com/Siguza/ios-resources/blob/master/bits/arm64.md](https://github.com/Siguza/ios-resources/blob/master/bits/arm64.md)

**"How to Read arm64 Assembly Language"**  
[wolchok.org/posts/how-to-read-arm64-assembly-language/](http://wolchok.org/posts/how-to-read-arm64-assembly-language/)

**"Introduction To Arm Assembly Basics"**  
[azeria-labs.com/writing-arm-assembly-part-1/](http://azeria-labs.com/writing-arm-assembly-part-1/)