

AES Wireless Keyboard

- Template Attack for Eavesdropping -

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http://www.sntworks.kr

March 22th, 2018

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- Kwonyoup Kim, Founder, CEO of SNT Works Inc.
- Tae Hyun Kim, PhD., CTO of SNT Works Inc.

- SNT Works
 - Specialized in Security Analysis on Embedded Systems
 - ✓ Reverse Engineering for Security Evaluation, Assessment
 - ✓ Side Channel Attack & Fault Attack
 - ✓ Patent infringement Investigation

Previous Related Works of Wireless Keyboard & Mouse

- KeyKeriki v1.0 and v2.0 by Dreamlab Technologies, 2010
- Promiscuity is the nRF24L01+'s Duty, Travis Goodspeed, 2011
- KeySweeper, Samy Kamkar, 2015
- MouseJack, Bastille Networks Internet Security, 2016
- Deeg & Klostermeier, Of Mice and Keyboards : On the Security of Modern Wireless Desktop Sets, Hacktivity 2016

Motivations

- Deeg & Klostermeier's Results, Hacktivity 2016
 - Found Security Vulnerabilities
 - 1. Insufficient protection of code (firmware) and data (cryptographic key)
 - 2. Unencrypted and unauthenticated data communication
 - 3. Missing protection against replay attacks
 - 4. Insufficient protection against replay attacks
 - 5. Cryptographic issues

Manufacturer Feedback

"As you called out in your email, given each wireless desktop set has different cryptographic key which makes this attack not generic at all. <u>It also requires physical access to the keyboard and</u> <u>sniffer to capture packets to decrypt with obtained key. If you can open keyboard and dump flash</u> from it you can as well change the whole board. Hence, this doesn't meet security servicing <u>bugbar</u>. We have opened a bug in the next version of the product for the core team to evaluate." (concerning insufficient protection of code and data)

Our Goals

- Eavesdropping of AES Wireless Keyboard without physical access
 - Reverse Engineering of AES Wireless Keyboard
 - ✓ Construction of SCA(Side Channel Attacks) environment
 - Side Channel Attacks of AES Wireless Keyboard
 - ✓ EM Side Channel Attacks
 - ✓ Recover secret 128 bits AES keys
 - Building templates using the recovered key
 - Attacking the other keyboard with different unknown key





AES Wireless Keyboard

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- Reverse Engineering -

Firmware & NVRAM Acquisition

• Keyboard (nRF24LE1)

Static Analysis

- Disassembling firmware (IDA Pro)
- Finding Encryption functions

Dynamic Analysis

- Debugging Codes (nRFgo Studio)
- Verifying Encryption functions (IV, Salt, Key, Enc-Mode, ...)

Firmware Modification

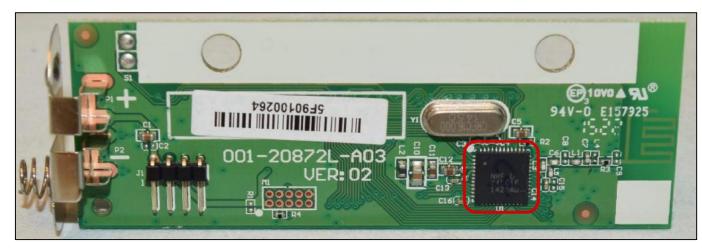
- Modifying Codes for Repeated execution
- Reprogramming of Keyboard firmware





Firmware Acquisition – Keyboard







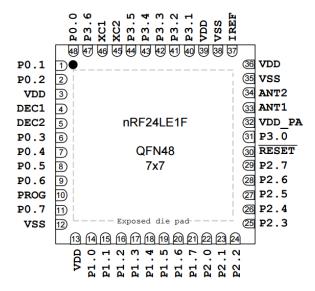
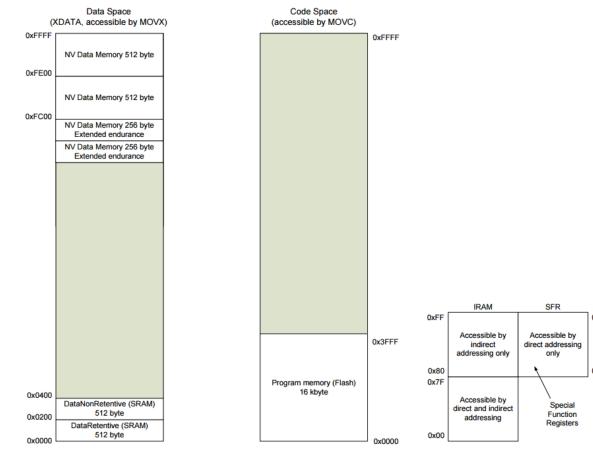


Figure 4. nRF24LE1F pin assignment (top view) for a QFN48 7×7 mm package

Firmware Acquisition – nRF24LE1

• Flash memory : features



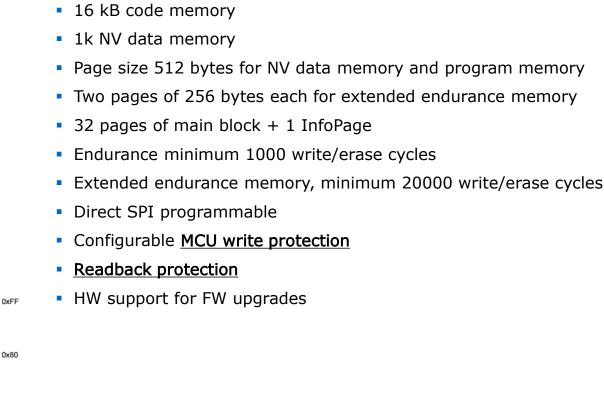
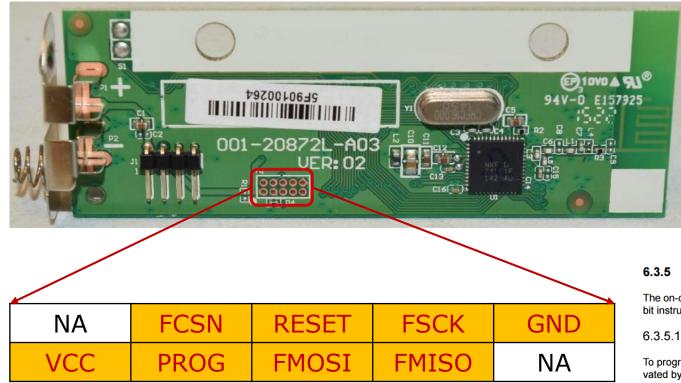
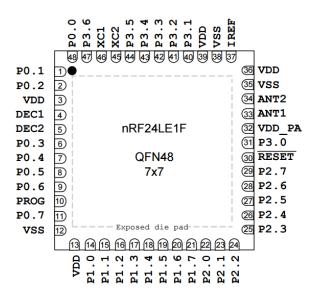


Figure 30. Memory map

Firmware Acquisition – Pinout





6.3.5 Flash programming through SPI

The on-chip flash is designed to interface a standard SPI device for programming. The interface uses an 8bit instruction register and a set of instructions/commands to program and configure the flash memory.

5.3.5.1 SPI slave interface

To program the memory the SPI slave interface is used. SPI slave connection to the flash memory is activated by setting pin PROG = 1 while the reset pin is kept inactive. After the PROG pin is set to 1, apply a pulse on the RESET pin (Pull RESET pin low for a minimum of 0.2 µs and return to high). Selected nRF24LE1 GPIO pins are automatically configured as a SPI slave as shown in <u>Table 33</u>. Further information on SPI slave timing can be found in <u>chapter 18 on page 147</u>.

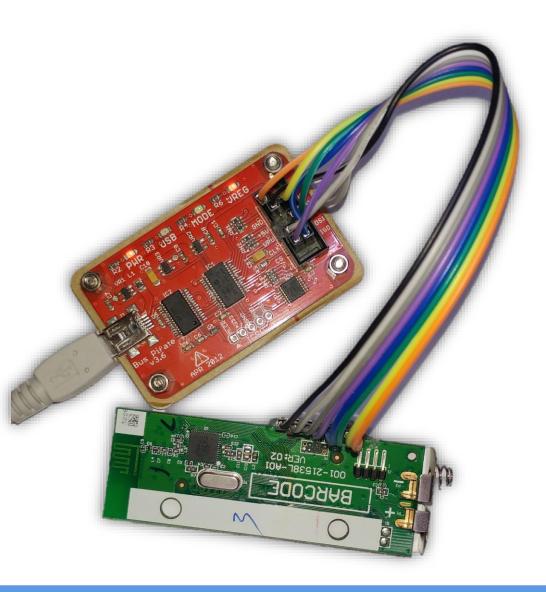
	24pin-4×4	32pin-5×5	48pin-7×7
FCSN	P0.5	P1.1	P2.0
FMISO	P0.4	P1.0	P1.6
FMOSI	P0.3	P0.7	P1.5
FSCK	P0.2	P0.5	P1.2

Table 33. Flash SPI slave physical interface for each nRF24LE1 package alternative

Firmware Acquisition – nRF24LE1

Flash memory Dump

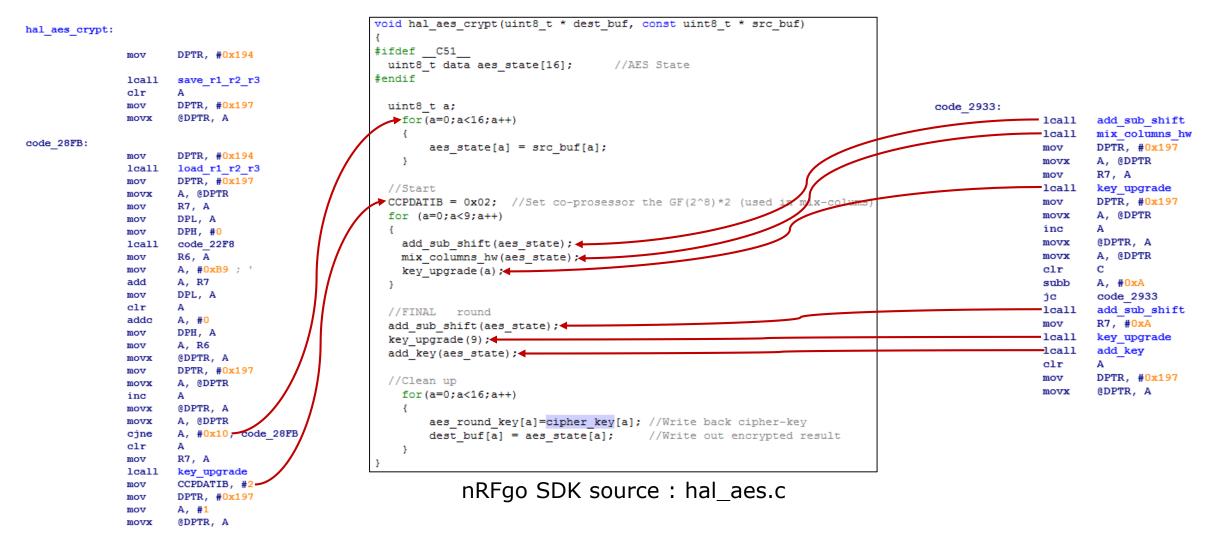
\$./nrfprog /dev/ttyUSB0 dump.bin Opening the Bus Pirate UART Setting the Bus Pirate to Binary Mode Configuring SPI mode Backing up the info page Starting read operation Reading from the device to dump.bin Starting read operation Putting the Bus Pirate back in normal operating mode Closing the Bus Pirate \$ ls -alh dump.bin info page.dat -rwxrwxrwx 1 root root 16K April 28 19:32 dump.bin -rwxrwxrwx 1 root root 256 April 28 19:32 info page.dat



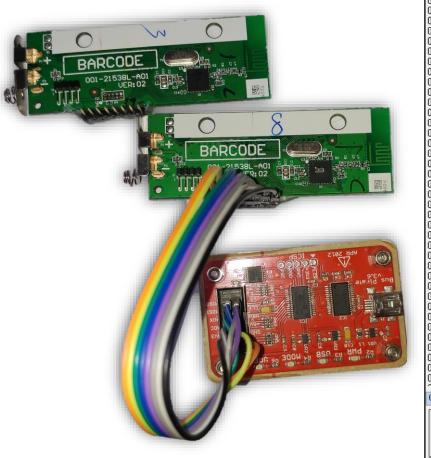
Static Analysis – Disassembling

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dump, bin 🗷 🧹 🤄	
▼ Edit As: Hex ▼ Run Script ▼ Run Template ▼	Library function 📗 Data 🔜 Regular function 📕 Unexplored 📕 Instruction 📒 External symbol
0 1 2 3 4 5 6 7 8 9 A B C D E F 0123456789ABCDER	
0000h: 02 28 56 02 00 11 C2 B9 D2 1A 32 02 00 12 D2 1F .(VÂ ¹ Ò.2Ò.	Function name A CODE: 00000000 ;
0010h: 32 32 32 02 00 19 AF D7 22 32 32 02 2F 66 32 32 222 × "22./f22	code:00000000 ; +
0020h: 32 32 FF 02 00 1A FF FF FF FF FF 02 00 1E 7C 00 22ÿÿÿÿÿÿ	• Code: 000000000 ; This file has been generated by The Interactive Disassembler
0030h: 7D F9 90 00 F9 E4 F5 3C E0 62 3C A3 DF FA E5 3C }ùùäð<àb<£ßúå<	code:000000000; Copyright (c) 2015 Hex-Rays, <support@hex-rays.com></support@hex-rays.com>
0040h: F4 F0 22 02 00 1F FF FF FF FF FF 02 00 06 FF FF 68"ÿÿÿÿÿÿ	
0050h: FF FF FF 02 00 20 FF FF FF FF 62 00 21 FF FF ÿÿÿ ÿÿÿÿÿ !ÿ	Y Z Image: Second seco
0060h: FF FF FF 02 2F FF FF FF FF FF 62 00 0E 7F 00 ÿÿÿ./ÿÿÿÿÿ	· TF1
0070h: 90 01 84 74 01 F0 A3 74 00 F0 75 34 08 90 01 84,t.&£t.&u4,	" code:00000000;
0080h: E0 FE A3 E0 F5 82 8E 83 E4 F0 90 01 84 75 F0 01 ap£aõ,Žfäð"uð.	<pre> code:00000000 ; code:00000000 ; code:00000000 ; lnput SHA256 : CC5519D4FCC41729BCE751A3A24F3E88C3B55F19B514688FD0 code:00000000 ; Input MD5 : 660E822C449E859AFCBC9567D36D489D code:000000000 ; Input MD5 : 660E822C449E859AFCBC9567D36D489D code:00000000 ; Input MD5 : 660E822C449E859AFCBC9567D36D489D code:00000000 ; Input MD5 : 660E822C449E859AFCBC9567D36D489D code:00000000 ; Input MD5 : 660E82C449E859AFCBC9567D36D489D code:00000000 ; Code:00000000 ; Input MD5 : 660E82C449E859AFCBC9567D36D489D code:0000000000 ; Input MD5 : 660E82C449E859AFCBC9567D36D489D code:000000000000000000000000000000000000</pre>
0090h: 12 23 59 D5 34 E7 7E 01 90 01 84 74 01 F0 A3 74 .#YÕ4ç~,t.&£t	code:000000000; Input MD5 : 660E822C449F859AFCBC9567D36D489D t TF2_EXF2 code:000000000; Input CRC32 : D6C11742
00A0h: 01 F0 75 32 01 75 33 73 75 39 06 E4 F5 3A F5 37 .ðu2.u3su9.äõ:õ7	7 If code 2E
00B0h: 75 34 10 05 33 E5 33 70 02 05 32 F5 82 85 32 83 u43å3p2õ,2j	f f code_C51
00C0h: E0 FF 70 03 02 03 43 75 17 06 8F 22 20 12 06 20 àÿpCu"	f code-F62
00D0h: 13 03 02 03 37 E5 33 24 F0 F5 31 E5 32 34 FF F57å3\$ðő1å24ÿč	f code: 00000000 : Base Address: 00000 = 4000b = 4000b loaded length: 4000b
00E0h: 30 85 31 82 F5 83 E0 90 37 6A 93 F5 36 64 F0 70 01.öfà.7j~č6dðr	f code_1490
00F0h: 23 7F B0 7E FF 12 35 3B 8F 3B E5 3B B4 FF 15 20 #.°~ÿ.5;.;å;'ÿ.	f code: 00000000 ; Progesson ; 8032 [PM-256 POM-0 FPPOM-0]
0100h: OC 12 85 31 82 85 30 83 E4 F0 85 33 82 85 32 831,0fäð3,2j	if if code_182F code:000000000; Target assembler: ASMI if code:19A8 code:000000000; Target assembler: ASMI
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0120h: E0 F5 38 75 35 0F 74 64 25 35 F5 82 E4 34 01 F5 àõ8u5.td%5õ,ä4.č	code:00000000
0130h: 83 E4 F0 74 74 25 35 F5 82 E4 34 01 F5 83 E4 F0 fäðtt%5ő,ä4.őfäð	Code: 00000000 ;
0140h: D5 35 E3 20 13 0C 90 01 64 E5 38 F0 90 01 74 74 05ãdå8ðtt	code 1540
0150h: 08 F0 90 01 02 74 4C F0 E4 A3 F0 A3 F0 A3 F0 A3tL8ä£8£8£8£	Code: 0000000 ; segment type: Pure code
0160h: F0 A3 F0 90 01 00 74 05 F0 75 3A 03 F5 37 02 03 8£8t.8u:.õ7	code:UUUUUUUU ;.segment code
0170h: 43 E5 36 C3 94 04 50 12 30 13 03 02 03 37 85 33 Cå6Ã".P.07	f code 21E7
0170n: 43 E5 36 C3 94 04 50 12 30 13 03 02 03 37 85 35 C40A .P.0 0180h: 82 85 32 83 74 08 F0 02 03 37 E5 36 C3 94 3A 50 ,2ft.ð7å6Ã":H	f code 21EE
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01A0n: 23 59 15 39 90 01 84 E0 FE A3 E0 F5 82 8E 83 E5 #1.9,apta0,2ja 01B0h: 36 F0 20 13 09 85 33 82 85 32 83 74 08 F0 05 3A 683,2ft.8.	
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01E0h: E4 75 F0 01 12 23 59 15 39 90 01 84 E0 FE A3 E0 äuð#Y.9"àþ£ð	à Python
Pos: 16384 [4000h] Size: 16384 ANSI LIT W OV	VR AU: idle Down Disk: 602GB
	N. 1010 DOWN DISK. 0020D

Static Analysis – Finding Encryption functions



- Static Analysis Finding AES key
 - Flash Dump from another keyboard



- Static Analysis Finding AES key
 - NVRAM are not the same values
 - Whether or not the AES key is validated using debugging

dumpA, bin 🗃		$\rightarrow = \mp$
★ 0123456789ABCDEF	0123456789ABCDEF 🛏 🚖 0 1 2 3 4 5 6 7 8 9 A B C D E F 0123456789ABCDE	F
3F50h: FF		. ^
3F60h: FF		
3F70h: FF		
3F80h: FF		
3F90h: FF	3F90h: FF	•
3FAOh: FF	or non-	•
3FBOh: FF	OI DOIN	•
3FCOh: FF	3FCOh: FF	•
3FDOh: FF	or boint in the second s	•
3FEOh: FF		•
SFFUR: DO DE OL EG LO DO DO DO LO EE OF DO DO EL OF	3FFUR:	7
		<u> </u>
4020h: 02 00 C0 21 01 00 00 00 00 00 00 00 00 00 E4 00 4030h: 01 CD 30 D5 02 AA 1A 03 13 4E 44 FF 00 03 43 04		
403011 FF F	403011 FF F	
4050h: FF	404011: FF	
4060h: FF		
4070h: FF		
4080h: FF	4080h: FF	
4090h: FF		
40A0h: FF	40A0h: FF	
40B0h: FF	40B0h: FF	
40COh: FF	40C0b; FF	
40D0h: FF	40DOh: FF	
40E0h: FF		. 🗸
40F0b: FF		. 🗆 📼
	Pos: 17178 [431Ah] Val: 255 FFh 11111111b Size: 18432 ASCII LIT W OV	VR 🔡

Dynamic Analysis – Debugging codes

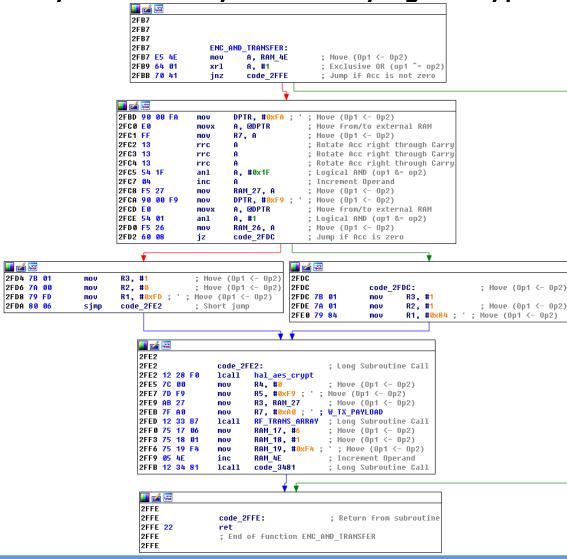
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Register	Value	<mark>⊏⟩C:0x2856</mark>	78FF	MOV	RO,#0xFF	^	Address: X:0x00			<u> </u>
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r0	0x4f	C:0x2859	F6	MOV	@RO,A		X:0x000000: 4B AD AC BC DC 06 64 24 83 8E 0C			
<mark>r1</mark>	0×0f	C:0x285A	D8FD	DJNZ	R0,C:2859		X:0x000010: 2D 6E 3F 5F EE F0 DB D7 57 3F 05	82 BC	F5 50) EF
r2	0x21	C:0x285C	900000	MOV	DPTR,#0x0000		X:0x000020: 29 E1 EF 37 95 57 D0 B5 F1 93 5B	5E B1	FA 93	81
r3 r4	0x86 0xd1	C:0x285F	7F00	MOV	R7,#0x00		X:0x000030: FD 4B 0E E8 F0 F0 B5 80 1F 2A A0	87 BD	F5 AE	: D0
r4 r5	0×01 0×92	C:0x2861	7E04	MOV	R6,#0x04		X:0x000040: 2B 47 46 09 67 32 D2 B9 0E 06 0E			
r6	0xa1	C:0x2863	E4	CLR	A		X:0x000050: 17 0D 43 4B 38 28 B1 B6 0F CF 85			
	0×70	C:0x2864	FO	MOVX	@DPTR,A		X:0x000060: 6C 48 8B EE DB 84 B8 7C 69 61 4E			
Sys		C:0x2865	A3	INC	DPTR		X:0x000070: 4F 0D 2E 2C E4 FD 66 C4 8E BF 0B			
a	0x00	C:0x2866	DFFC	DJNZ	R7,C:2864		X:0x000080: EA 68 D1 70 8B 07 39 26 FA F8 39	19 23	B9 59	98
b	0x00	C:0x2868	DEFA	DJNZ	R6,C:2864		X:0x000090: D0 E4 EC E4 52 4E AE BF AC 01 D2			
<mark>sp</mark> dptr	0×07	C:0x286A	758161	MOV	SP(0x81),#0x61		X:0x0000A0: F5 E2 F0 A6 2C C5 45 17 F3 76 90	B4 OE	OF 4B	3 0B
PC \$	0x0000 0x2856	C:0x286D	0228AB	LJMP	C:28AB		X:0x0000B0: 61 EC F8 E0 2B 3F 2F 2D 75 B7 E5	8A 92	AF 63	2B
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		C:0x2874	93	MOVC	A, @A+DPTR		X:0x0000E0: 20 84 C0 CE 01 0D BD 85 B8 DB F1	DO 2A	1F D7	5D
		C:0x2875	A3	INC	DPTR		X:0x0000F0: 90 17 74 E4 C1 0F 0E 87 F0 F7 F1	61 6D	CD E7	0D
		C:0x2876	F8	MOV	R0,A		X:0x000100: 24 10 F4 FD F0 FB 58 7C 2E 0E 8D	OF C6	94 D5	76
		C:0x2877	E4	CLR	A		X:0x000110: 8F 8A CC 87 B7 41 65 78 1F 0F 0F	A6 D5	D5 78	60
		C:0x2878	93	MOVC	A, @A+DPTR		X:0x000120: CD 89 1F BC AE B8 72 FA AB 24 26	E5 D9	D0 59	9 C8
		C:0x2879	A3	INC	DPTR		X:0x000130: 2E 09 DB 4D 68 B0 61 E0 0D D4 75	7D F5	C4 C5	• E7
		C:0x287A	4003	JC	C:287F		X:0x000140: CF OC 9F BF D0 6A 6A CD 3F 0F 8F	07 F2	9E E9	71
		C:0x287C	F6	MOV	@RO,A		X:0x000150: 0A 09 C9 8B 74 74 E2 E6 0B 9A 3F	9F A8	79 D0) F2
		C:0x287D	8001	SJMP	C:2880		X:0x000160: AE OF 06 DF F2 C0 F2 B0 9F 07 57	67 96	96 B6	5 F6
		C:0x287F	F2	MOVX	@RO,A		X:0x000170: 1E 12 27 46 38 EA E0 F2 F4 34 E8	7A 87	OE OD) OF
		C:0x2880	08	INC	RO		X:0x000180: 9B C3 C2 6B 07 AB AB 60 5E 75 44	CF 4B	7F A5	6F
		C:0x2881	DFF4	DJNZ	R7,C:2877		X:0x000190: B3 B2 D0 F0 01 A9 88 0D DB CB 78	D9 B8	B6 1C	; OD
		C:0x2883	8029	SJMP	C:28AE		X:0x0001A0: D5 E1 E0 EB 4E CD 8C DA A0 80 90	B5 99	9D DF	5 9E
		C:0x2885	E4	CLR	A		X:0x0001B0: 35 B9 F0 36 53 2B AF 0D D2 F2 E0	F0 F9	AF 25	47
		C:0x2886	93	MOVC	A, @A+DPTR		X:0x0001C0: 15 18 30 98 C5 0B 38 9B EC F0 76	F6 05	05 88	EO
		C:0x2887	A3	INC	DPTR		X:0x0001D0: 20 1C 5E DA 77 3D 26 2E F2 F0 08	44 OE	4C A9) 3E 🗸
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		C:0x288B	240C	ADD	A,#0x0C		Address: X:0xF9			
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🗉 Project 🛛 🚟 Reg	isters						X:0x000119: OF OF A6 D5 D5 78 60 CD 89 1F BC	AE B8	72 FA	AB 🗸



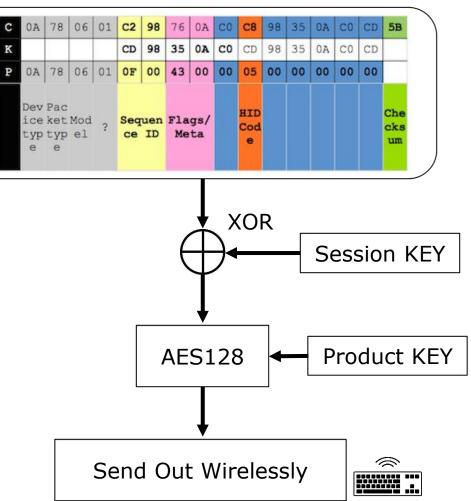
nRFgo Start Kit & nRF24LE1

Programmi	ng nRFModule on Board 2 •
Programming	g Advanced settings
Se <u>c</u> ond file:	Browse
	☑ Verify after programming
	🗹 Enable hardware debugging
	Block readback of mainmemory
	Protected program memory
	Enable 512 bytes 🗣 No pages protected
	Program <u>V</u> erify <u>R</u> ead nRF Re <u>s</u> et

Dynamic Analysis – Verifying Encryption functions



Plain Packet



- Firmware Modification Modifying codes for Repeated execution
 - Initial input of all zero bytes
 - Output \rightarrow Next Encryption input
 - Inserting artificial trigger signal before each encryption
 - 100 ms delay

- Firmware Modification Reprogramming of Keyboard firmware
 - Write protection is set for the target keyboard
 - So, We changed it to a reprogrammed chip for Side channel analysis

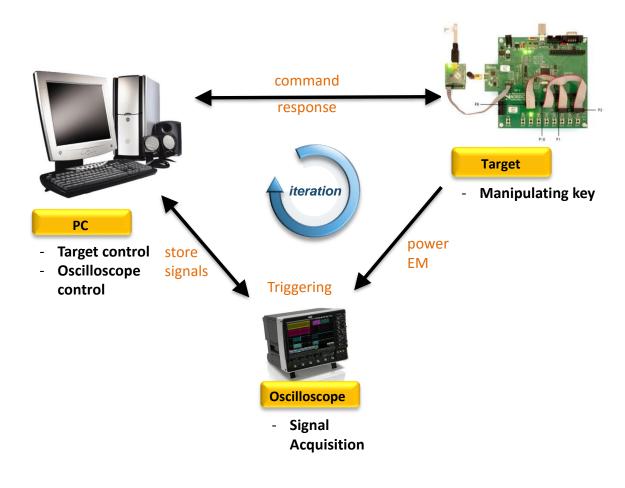


AES Wireless Keyboard

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- Side Channel Attacks, Template Attacks -

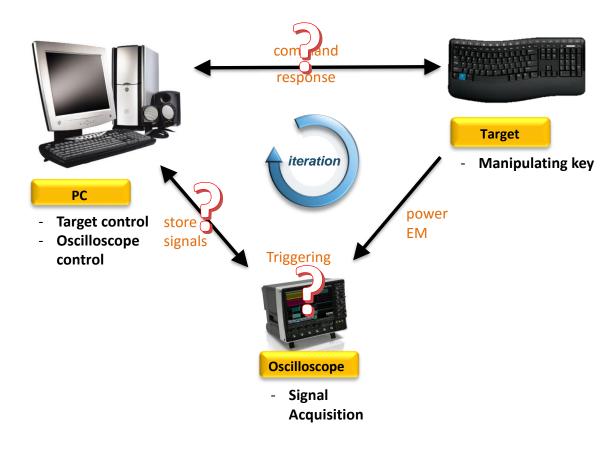
Side Channel Attacks in Theory



White-box Attack Model

- All information known and controllable
 - Known Specification
 - Known Interfaces
 - Known Design (by oneself)
 - Known Implementation details
 - Ideal Setup

Side Channel Attacks in Real World



Black-box Attack Model

- Most information unknown and controllable
 - Known Specification Only
 - Unknown Interfaces
 - Unknown SW/HW Design
 - Unknown Implementation details
 - Many **unknown** factors for setup

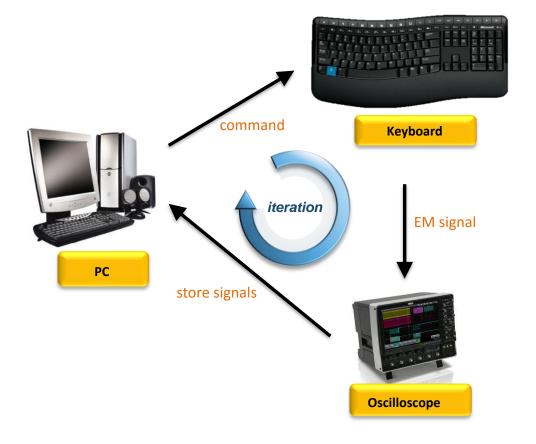
Side Channel Attack Environment



<u>Setup</u>

- ① Target Keyboard with AES key
- 2 Langer LF-1 near-field probe set
- (3) Langer PA303 pre-amplifier
- (4) Digital Oscilloscope
- (5) Software to control a target and an oscilloscope

Side Channel Attack

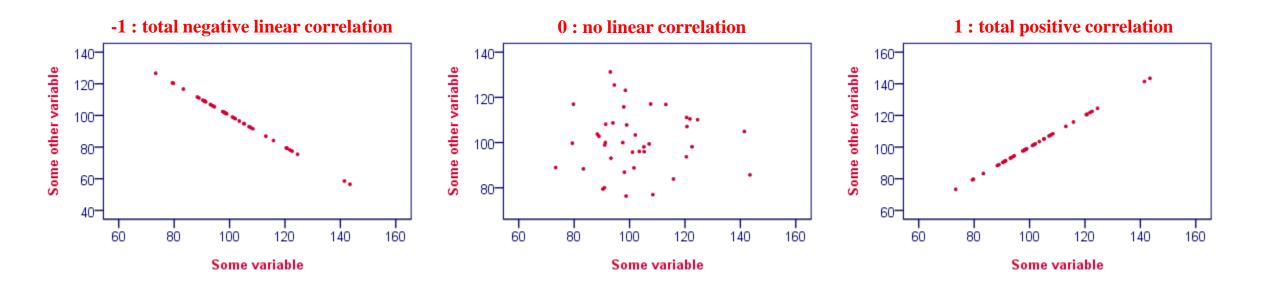


Attack Process

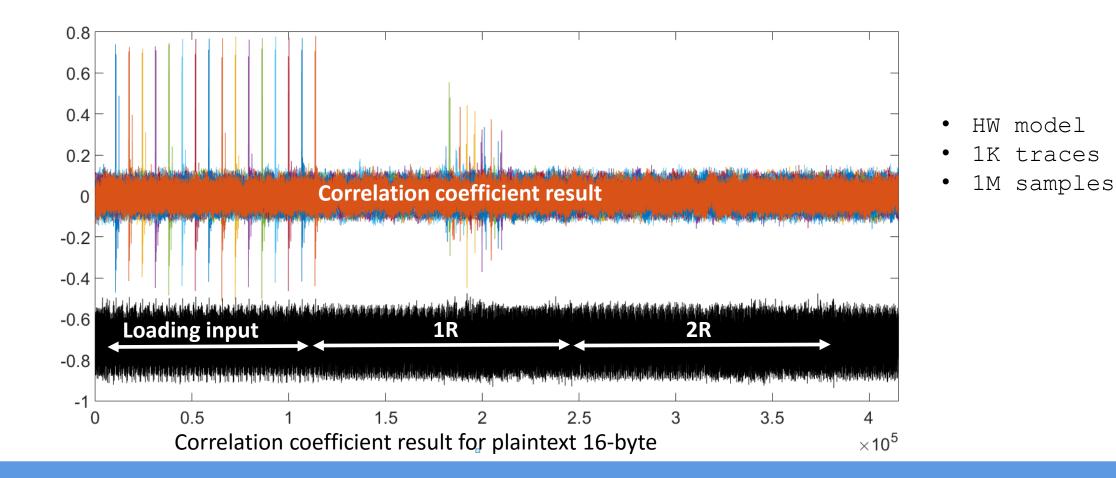
- Acquisition
- Alignment
- Correlation with input/output
- Profiling AES algorithm
- CPA
- Building templates
- Template attack

- CPA (Correlation Power Analysis)
 - Pearson Correlation Coefficient
 - *T* : measurement, *P* : predictions
 - Number between -1 and 1

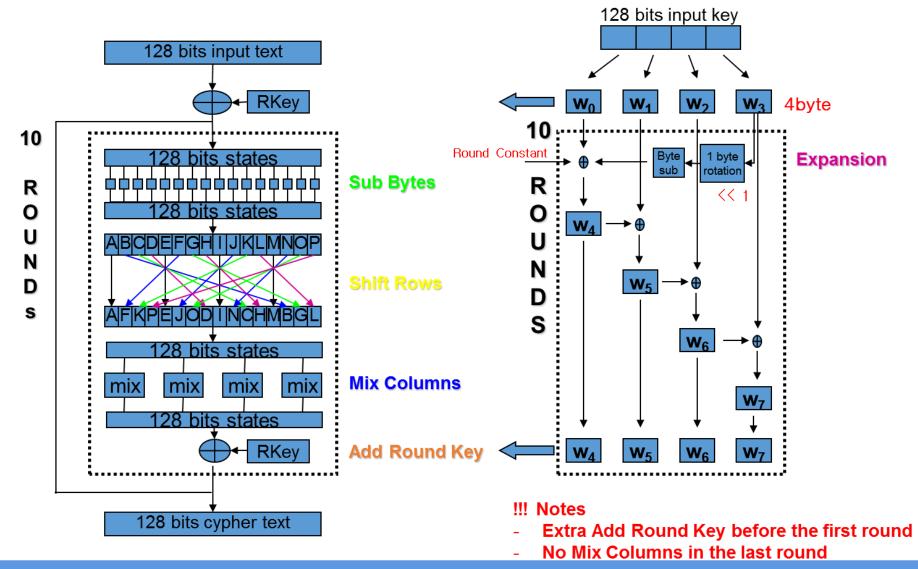
$$\rho(T,P) = \frac{E(TP) - E(T)E(P)}{\sqrt{Var(T)Var(P)}}$$



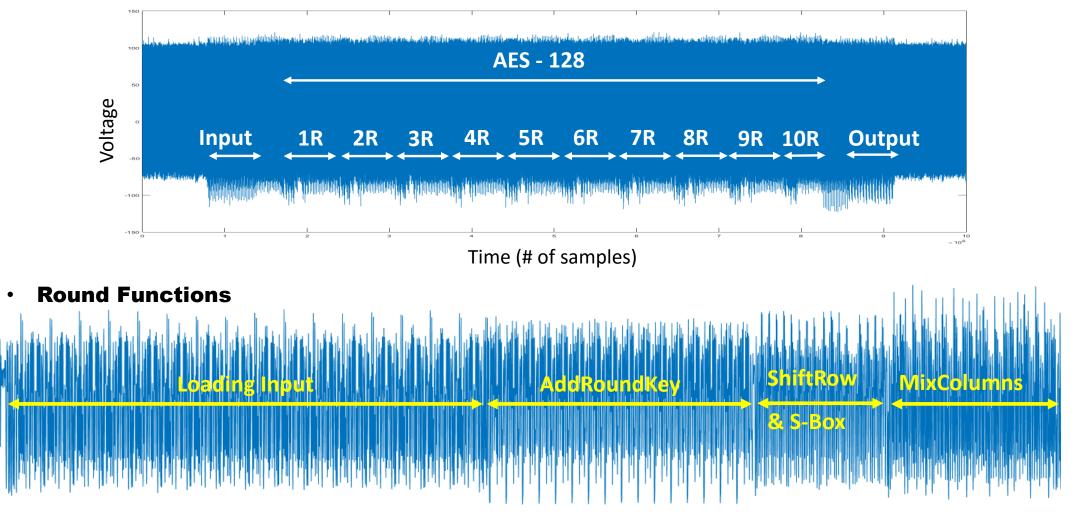
- CPA (Correlation Power Analysis) with input
 - Confirm plaintext loading operation



AES (Advanced Encryption Standard)

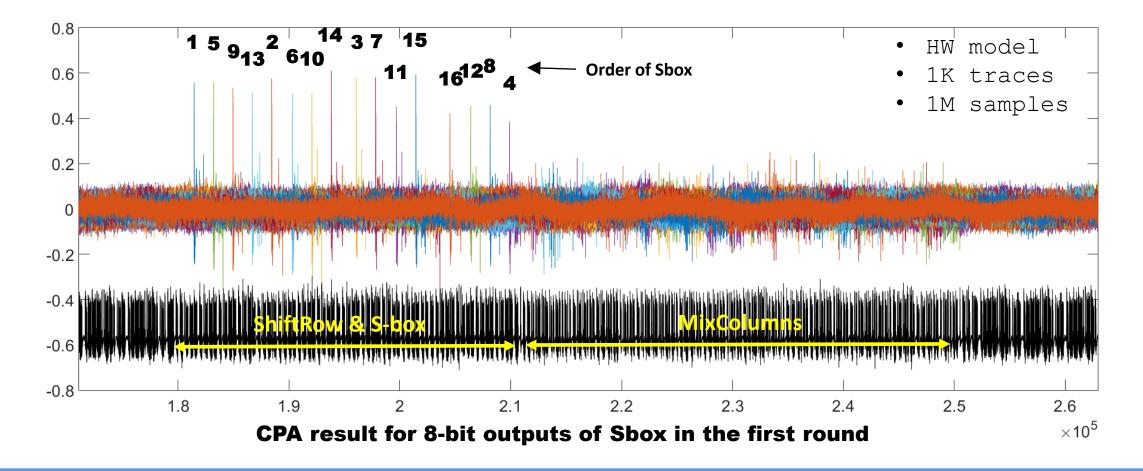


Profiling EM Trace

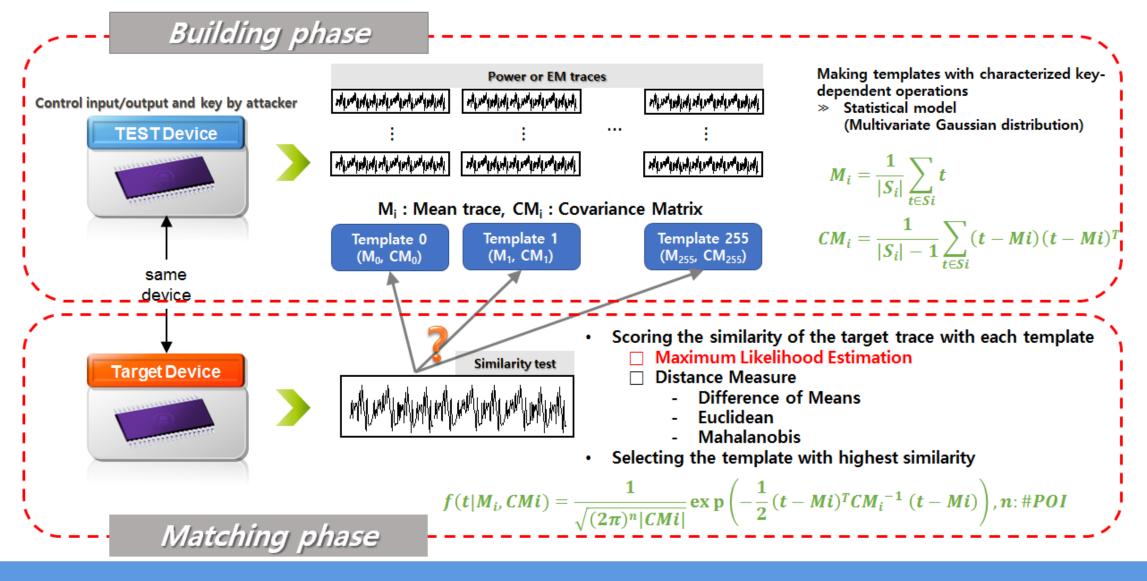


X Hardware Co-processor of MixColumns operation

- CPA (Correlation Power Analysis)
 - Recover the secret key
 - Target variable : S-Box output



Basic Template Attacks



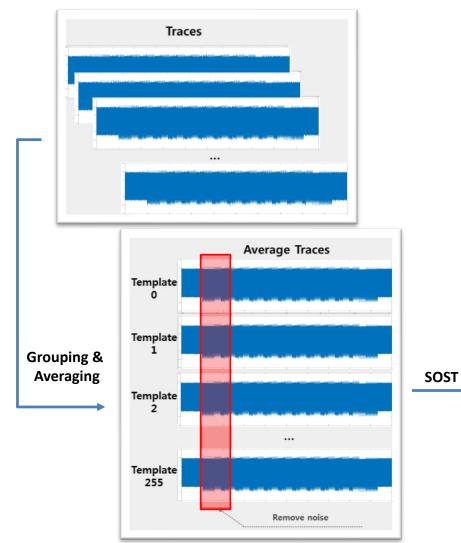
Selecting Interesting Points

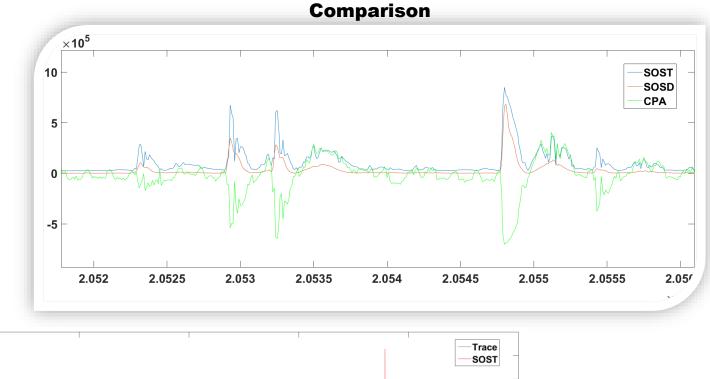
- Computational or memory restrictions
- Reducing the number of samples and the size of templates
- Selection of some special samples that contain the most information about the characterized key-dependent operation
 - ✓ DOM (Sum of pairwise Differences)
 - ✓ SOSD (Sum of Squared pairwise Differences)
 - ✓ SOST (Sum of Squared pairwise T-differences)
 - ✓ NICV (Normalized Inter-Class Variance)
 - ✓ PCA (Principal Component Analysis)
 - ✓ LDA (Linear Discriminant Analysis)

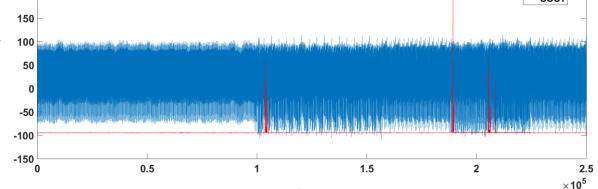
250

200

Selecting Interesting Points – SOST





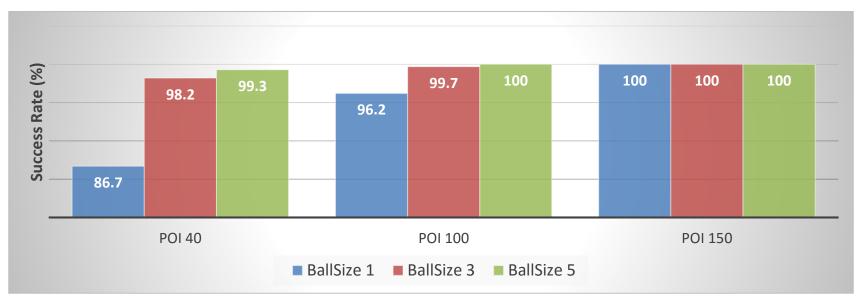


Building Templates

- 256 templates for a bytes AES key
- Training set : 1K traces for each template (total: 256K)
- Test set : 1K traces for each value (total: 256K)
- Point of Interest : 150 points

- Method : SOST, PCA, Pooled Covariance Matrix
 - Case 1: same device (profiling = matching device)
 - ✓ Case 2: different device (profiling ≠ matching device)

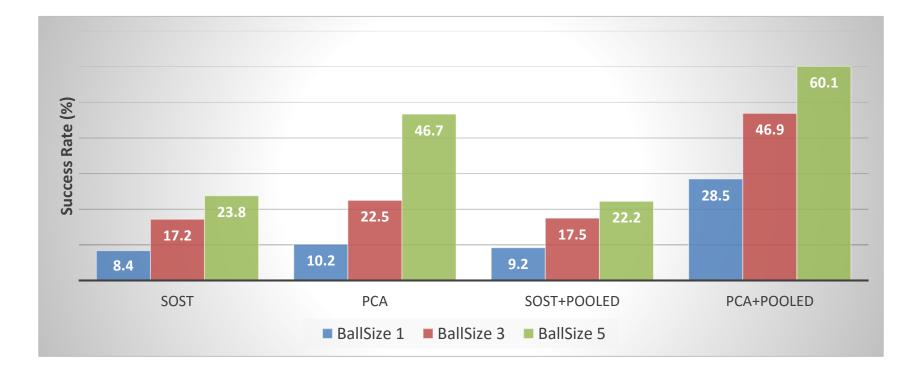
- Template Attacks : Case 1 Same device
 - Selecting POI : SOST
 - Template Matching : Maximum Likelihood



- Success Rate is the probability that the correct key is ranked among the first *n* candidates.
- **BallSize** *n* means the top *n* candidates ranked according to score, highest to lowest.

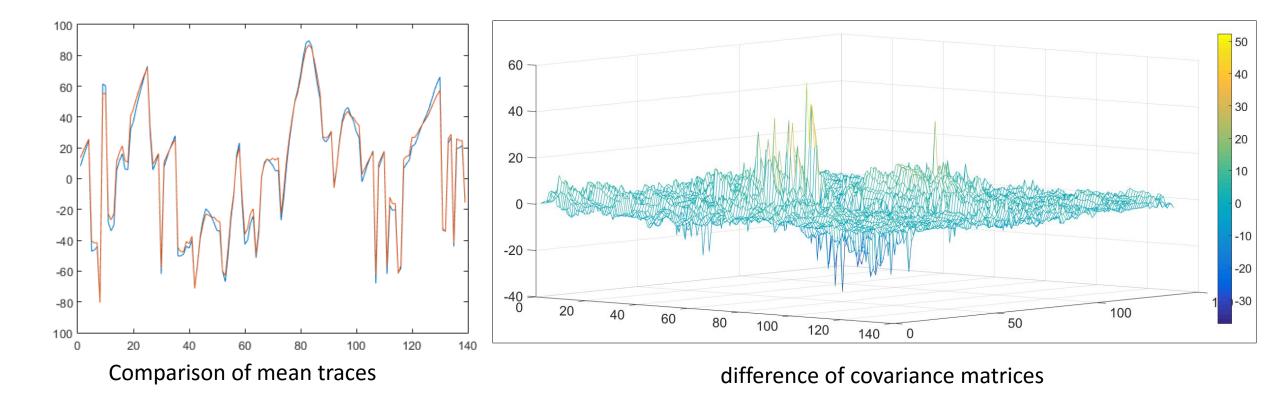
Template Attacks : Case 2 – Difference device

- Selecting POI : SOST, PCA, Pooled Covariance Matrix
 - ✓ Common POI of Device 1 (training) and POI of Device 2 (target)
 - ✓ 139 Points in 150 points
- Template Matching : Maximum Likelihood



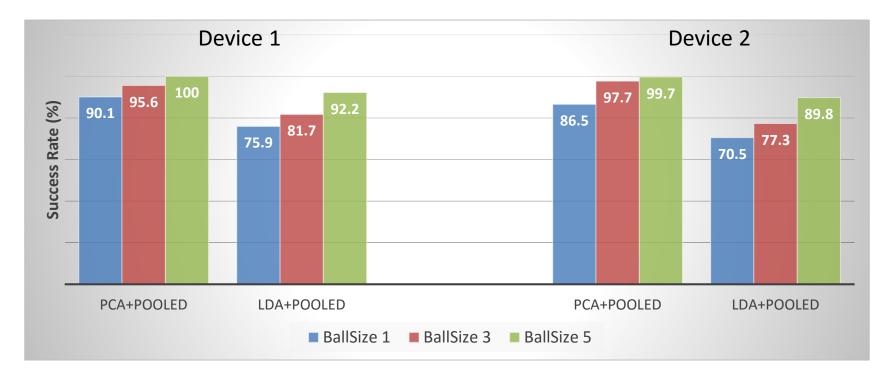
Template Attacks : Case 2 – Difference device

Characteristics comparison between two devices



Enhanced Attacks : Case 2 – Difference device

- Rebuilding robust templates using both training device and target devices^{*}
- Matching with robust templates to device 1, 2 respectively



*

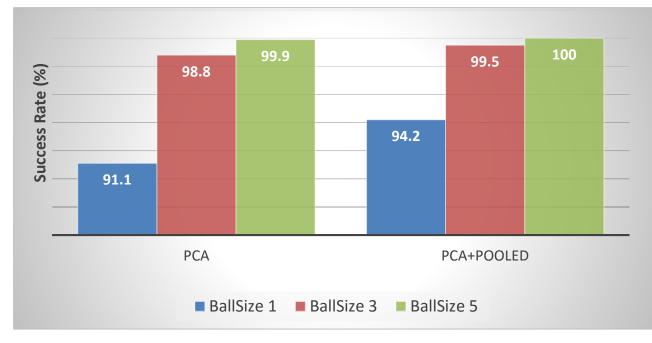
O. Choudary and M. G. Kuhn, "Template Attacks on Different Devices", in COSADE 2014

M. Renauld, F.-X. Standaert, N. Veyrat-Charvillon, D. Kamel, and D. Flandre, "A Formal Study of Power Variability Issues and Side-Channel Attacks for Nanoscale Devices", in EUROCRYPT 2011

- Proposed Attacks : Case 2 Difference device
 - Adjusting the difference between means of two devices by correction vector

 $Correction \ vector = \frac{mean \ of \ target \ traces}{mean \ of \ template \ traces}$

- Rebuilding templates using corrected traces on training device
- Matching with corrected templates to target device



- Result
 - Success rate for case 1, 2

					(70)
		the nur	The used number of		
		1	3	5	target trace
Same device		100	100	100	1
Difference device	Single	28.5	46.9	60.1	1
	Robust	88.3	96.65	99.85	1,000
	Corrected	94.2	99.5	100	256

※ Our attack requires the number of traces much less than the robust template attack

(%)

- Side-channel attack on secure wireless keyboard
- Setup for SCA environment using reverse engineering in Real World
- Performing EM analysis attacks to recover the AES key
- Template attack on the other keyboard

- Future research
 - We need more improved template attack
 - We need Far-field EM attack (long distance)
 - We need to apply it to other devices of similar types

• Side-channel attack is a serious threats in real-world.

• Wireless keyboards with AES is not secure.

Manufacturers need to defend against reverse engineering and side-channel attacks.



Thanks!

Any Questions?

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