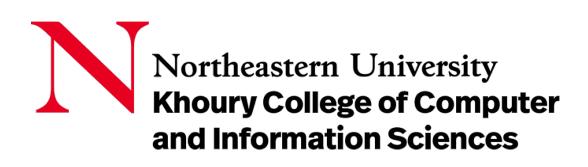
Practical Side-Channel Attacks against WPA-TKIP

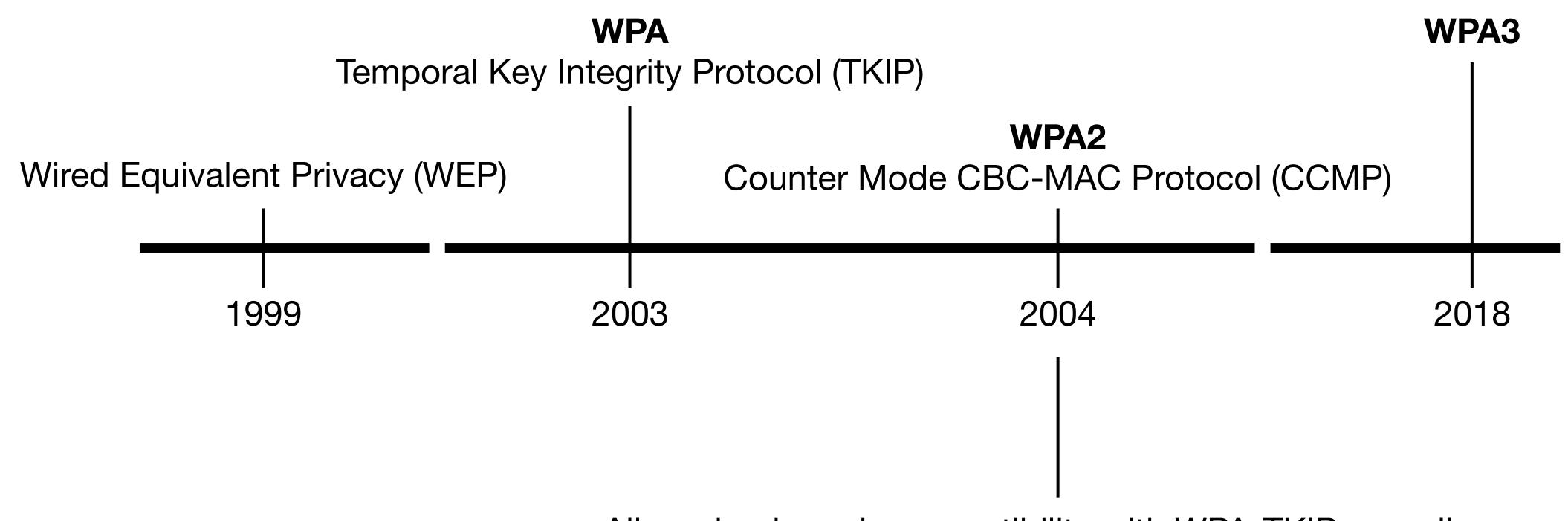
Domien Schepers In collaboration with Aanjhan Ranganathan and Mathy Vanhoef.

2019 Black Hat Europe





Brief Wi-Fi History



Allows backwards compatibility with WPA-TKIP as well.

Why are we interested in TKIP?

- Not marked as obsolete by the IEEE.
- Commercial products sold today still support WPA-TKIP.
- We still see a large support rate for WPA-TKIP in our Wi-Fi surveys.

2018 May

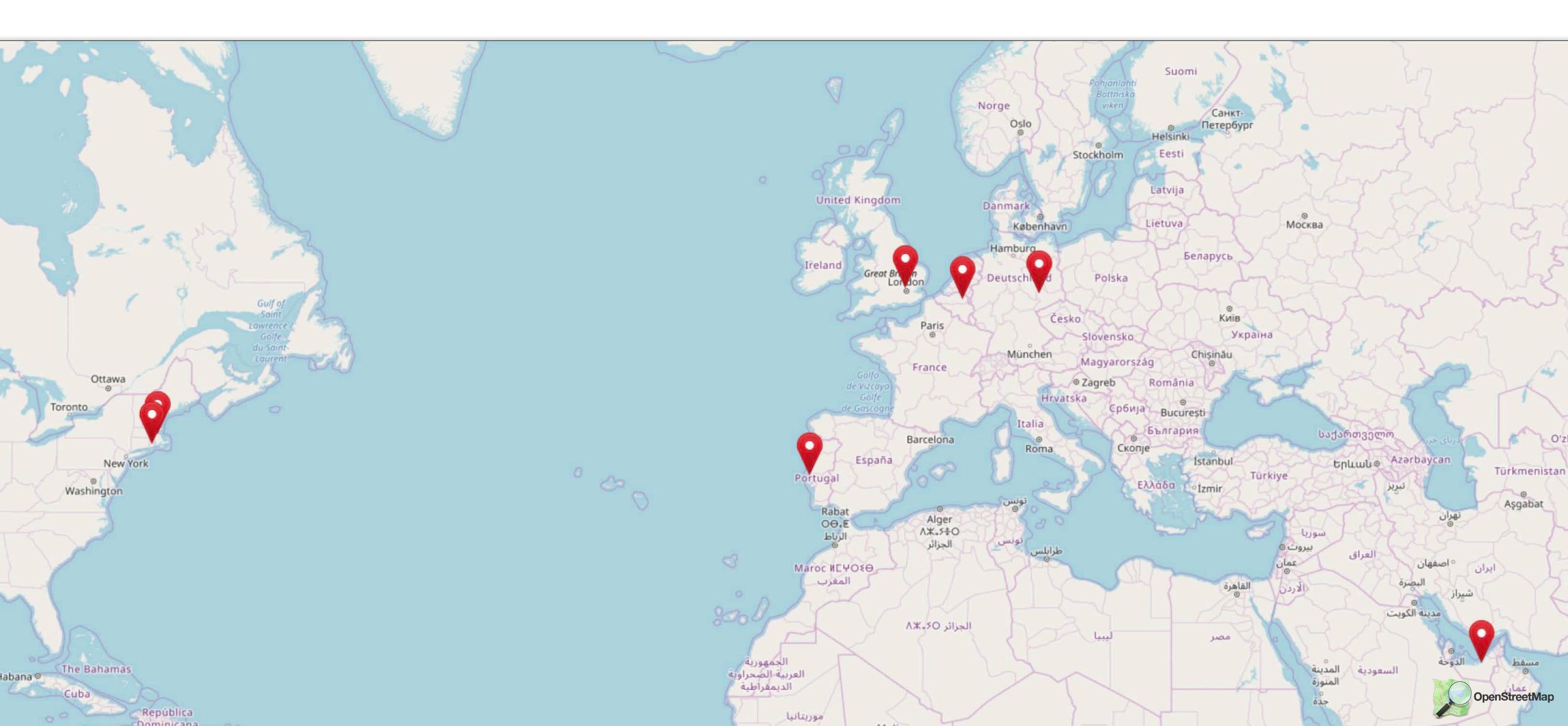
2.6.2.6.1. Discussion: cipher. 2.6.2.6.2. Straw Poll Results: A) 1 B) 10 C) 8 D) 11 2.6.2.7. Another Straw Poll (pick one): A) Delete WEP text only, make TKIP Obsolete. B) Delete WEP and remove TKIP as a pairwise cipher (continue supporting TKIP as a group cipher). C) Do nothing 2.6.2.7.1. Another Straw Poll Results: A) 6 B) 4 C) 12 2.6.2.8. Straw Poll #3: Mark TKIP as Obsolete? 2.6.2.8.1. Straw Poll #3 Results: Y: 6 N: 12 **2.6.2.9.** Therefore, resolve these 6 CIDs with Rejection, that we considered and decided to make no change. 2.6.2.10. Will resolve another CID (in PHY) about a WEP MIB variable, with the boilerplate answer about not maintaining a deprecated feature.

doc.: IEEE 802.11-18/0616r0

- 2.6.2.6.1.1. Don't think we can remove TKIP for pairwise



Wi-Fi Survey



Wi-Fi Survey

- Survey across several countries shows a large support for WPA-TKIP.
- On average 44.81% of networks supported TKIP in its Group Cipher.
 - Hasselt (Belgium) averages around ~58%.
 - Leipzig (Germany) averages around ~28%.



Outline

- Background on WPA-TKIP and existing attacks.
- Presentation of our side-channel attacks.
 - Tested against Linux and OpenBSD, MediaTek and Broadcom devices.
- Demonstration.
- Conclusion.

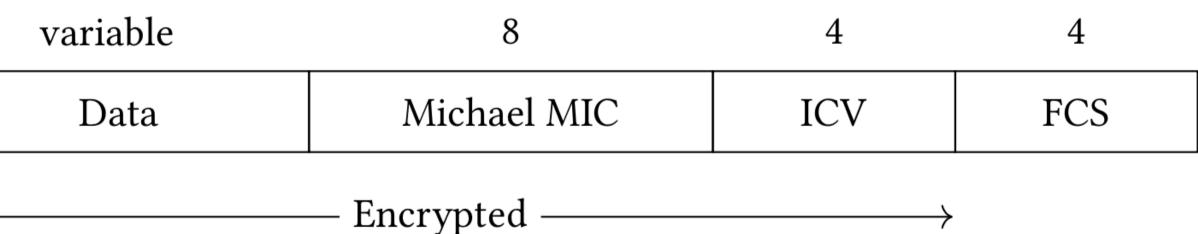


TKIP Header

		4	4	
MAC Header	Sequece No.	IV/KeyID	Ext. IV	

- Michael MIC provides better integrity protection than the ICV.

 - Added to the last fragment.
- TKIP Sequence Counter (TSC) for replay protection.



Covers plaintext data, destination and source address, and priority.



TKIP Countermeasures

- Michael MIC is invertible (Wool, 2004).
- Countermeasures state what to do if an invalid Michael MIC is received.
 - Client transmits a "Michael MIC Failure Report".
 - Access Point blocks TKIP traffic if two are received in one minute.
 - After one minute, clients may reassociate and negotiate fresh keys.

Quality of Service

- Defined in the IEEE 802.11e amendment (2005).
- Provides eight different channels for different QoS needs.
 - Michael MIC covers this channel number (i.e. priority).
 - TKIP Sequence Counter (TSC) for each channel.
- For example, used when making phone calls over Wi-Fi.

The ChopChop Attack

- The ChopChop attack originally targeted WEP (KoreK, 2004).
- Repeatedly chop off the last byte, and make corrections for the ICV.
 - Due to CRC linearity, ICV can be reconstructed if plaintext is known.
- Beck and Tews presented a ChopChop attack against WPA-TKIP (2009).

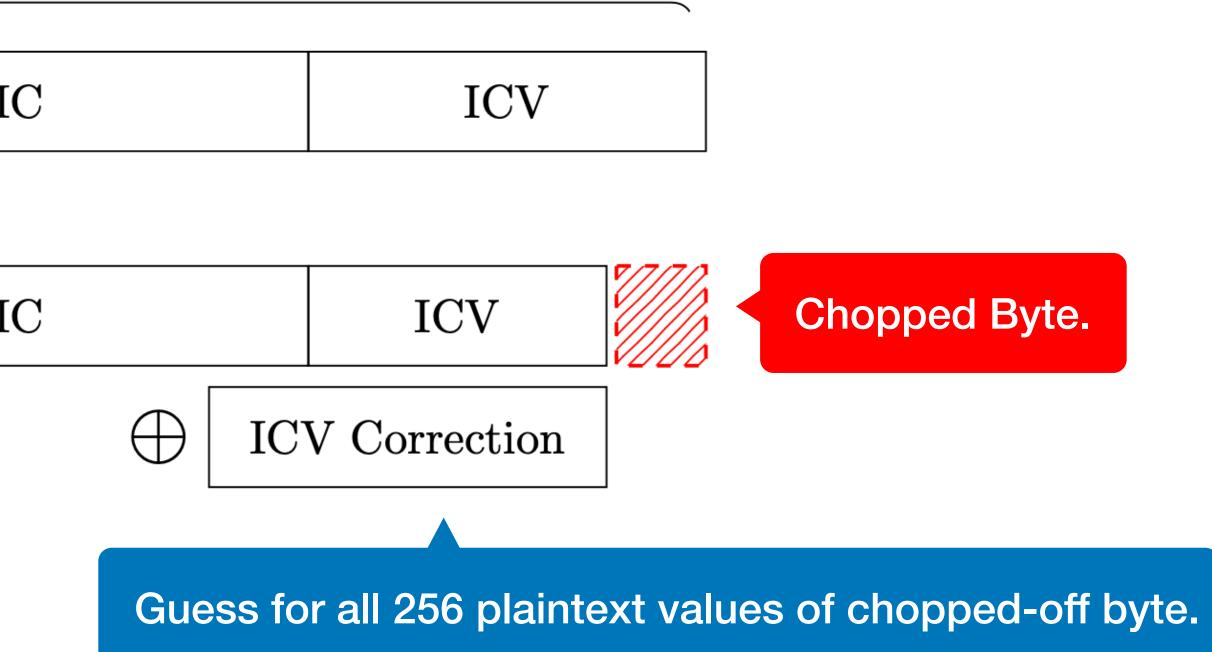
r		
Data	Michael MIC	ICV

	Encrypted		
Data	Michael MIC	ICV	
Data	Michael MIC	ICV	Chopped Byte.

Data	Michael MIC	ICV	
Data	Michael MIC	ICV	Chopped Byte.

	Encrypted
Data	Michael MIC

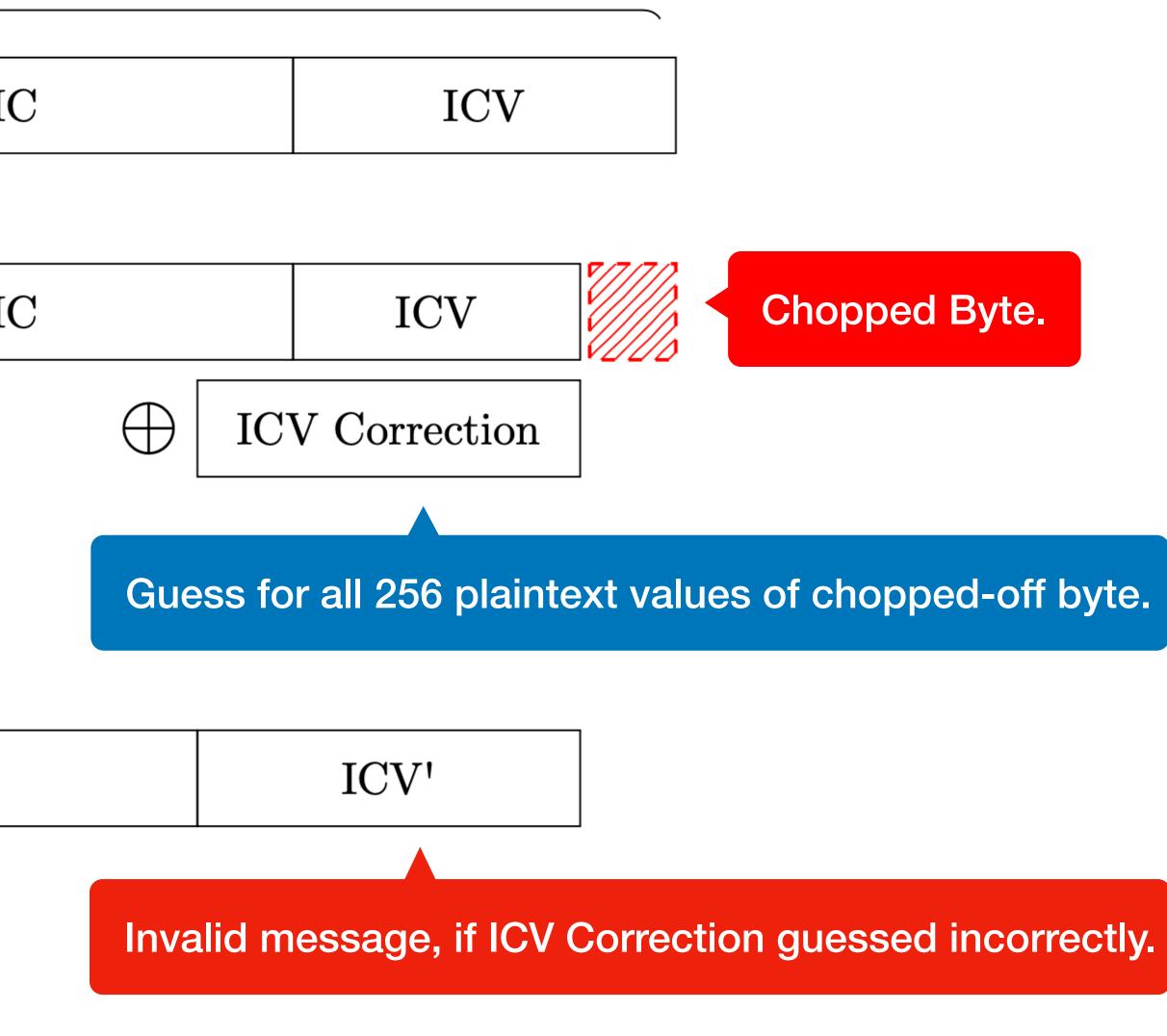
Data	chael MI
------	----------



	Encrypted
Data	Michael MIC
	•

Data	Michael MIC

Data'	Michael MIC'
-------	--------------

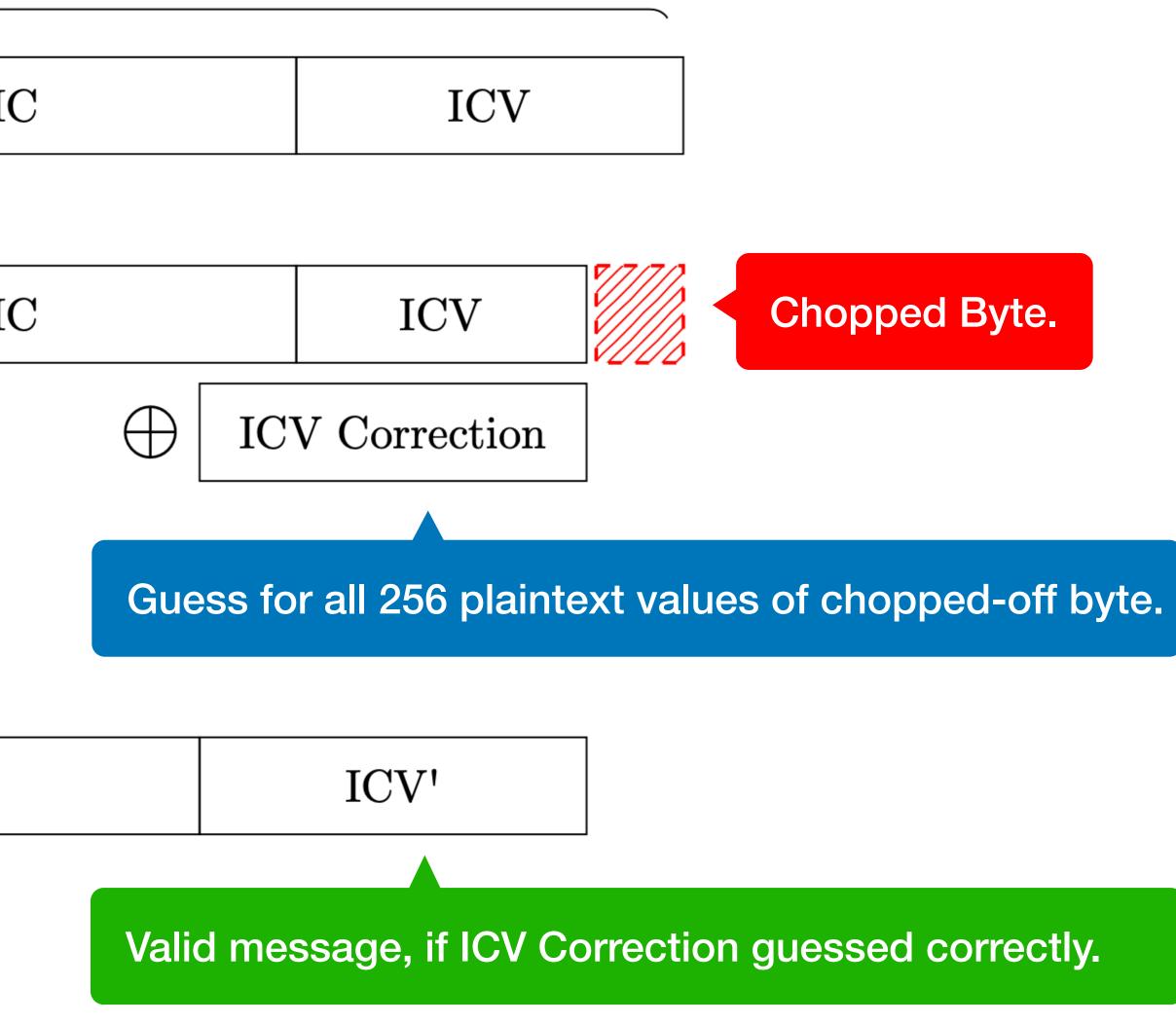




	Encrypted
Data	Michael MIC
	•

Data	Michael MIC

Data'	Michael MIC'
-------	--------------

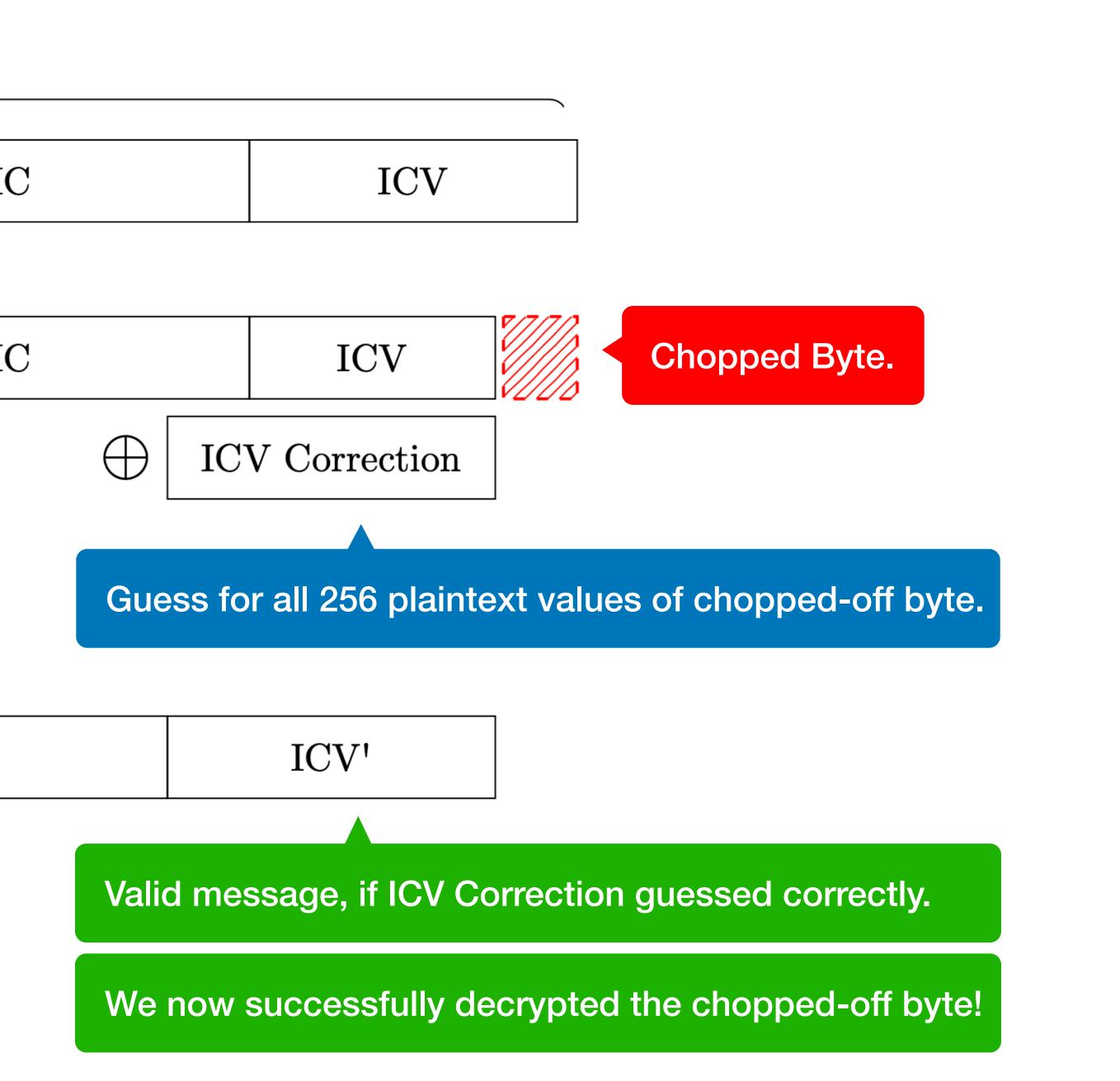




	Encrypted
Data	Michael MIC

Data	Michael MIC
1 1	

Data'	Michael MIC'
-------	--------------



How to detect a valid ICV Correction?

That is, having correctly guessed the plaintext value.

How to detect a valid CV Correction?

That is, having correctly guessed the plaintext value.

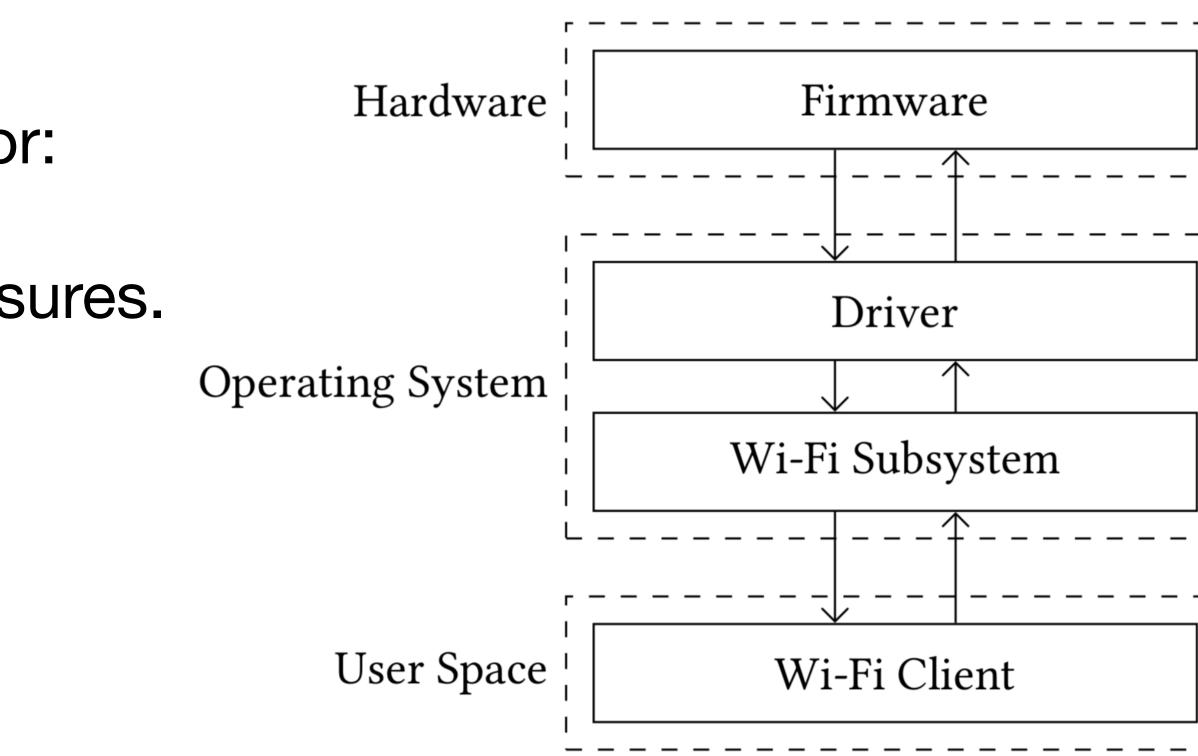
Side-Channels!

Side-Channels

- We want side-channels allowing us to detect if a TKIP frame is accepted.
 - Can we craft TKIP frames that influence the state of its receiver?
 - If so, can we find a way to detect the state change?
- We will analyze WPA-TKIP implementations to find such side-channels.
 - Allowing us to perform a ChopChop-like attack.

Methodology

- Inspect full Wi-Fi stack, looking for:
 - Implementation of countermeasures.
 - Side-channel vulnerabilities.
 - Logical bugs.



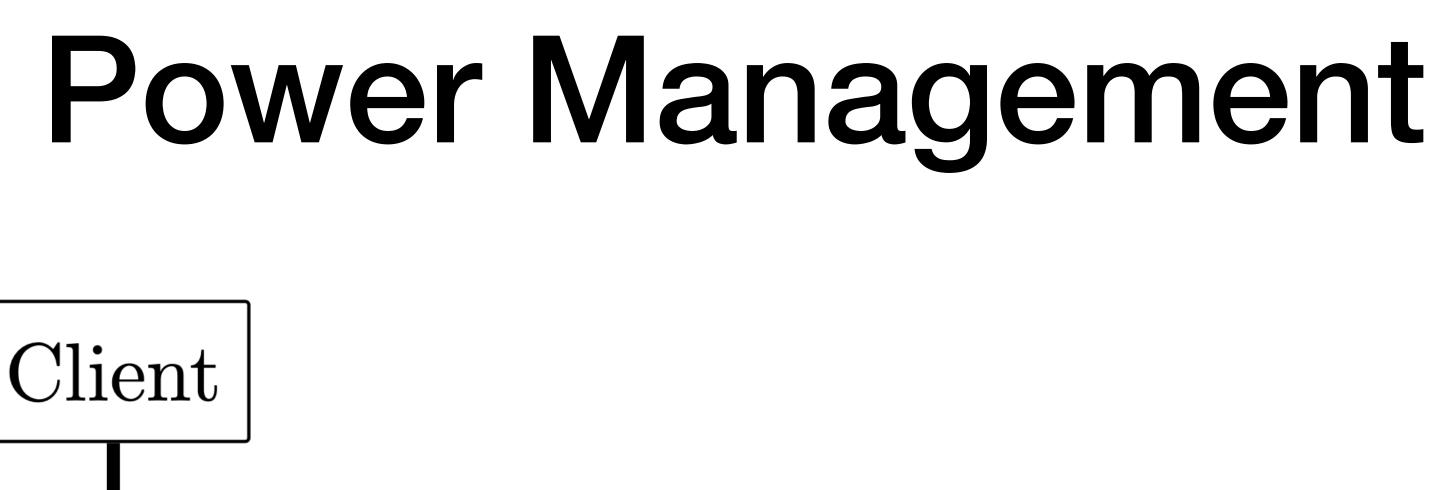
Results

- Making use of some of the following features:
 - Power Management.
 - Fragmentation.
 - Hardware Decryption.

Our approach resulted in novel (cross-layer) side-channel vulnerabilities.

Power Management

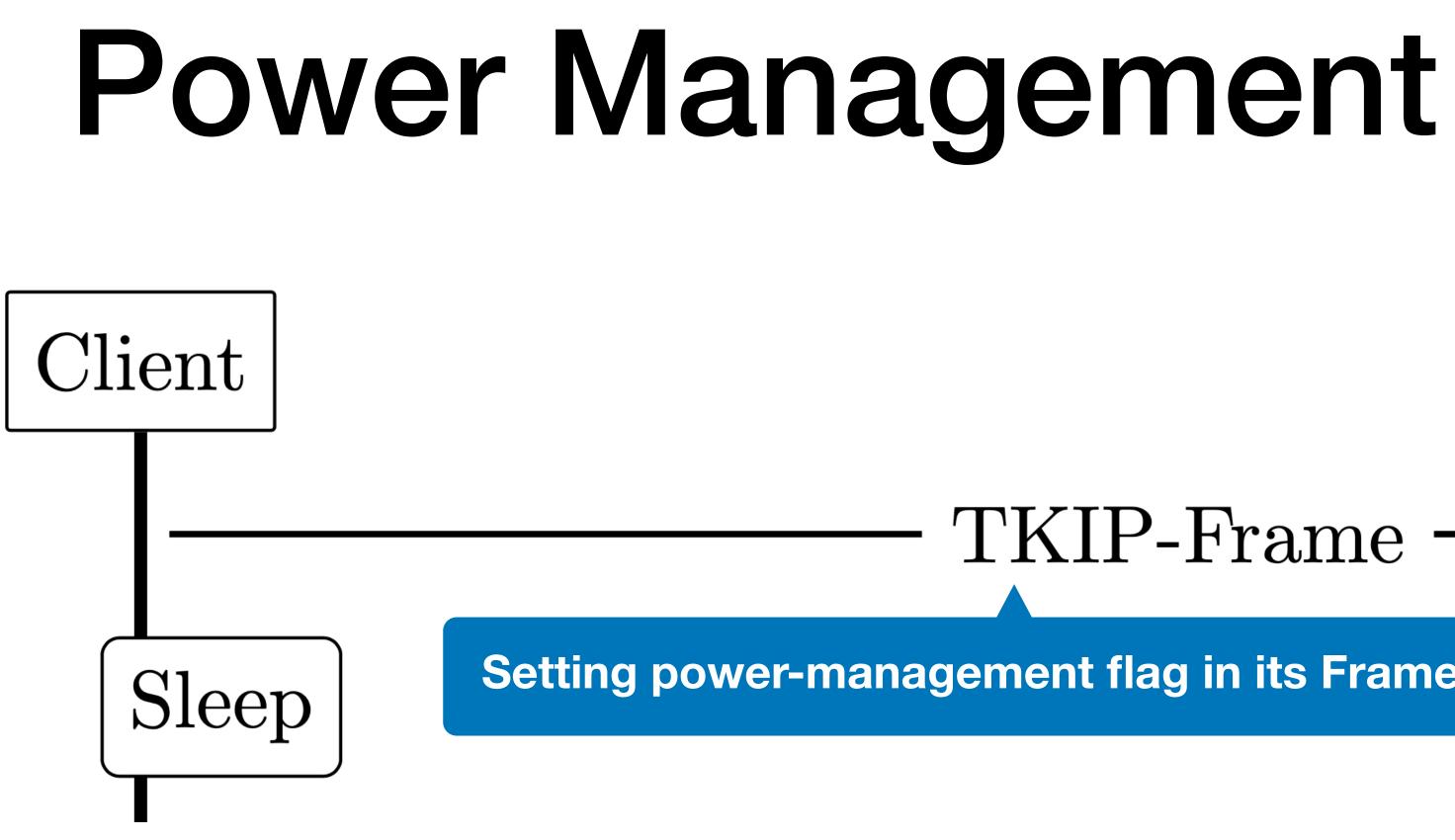
- IEEE 802.11e defines power save delivery and notification mechanisms.
 - E.g., Unscheduled Automatic Power Save Delivery (U-APSD).
 - E.g., Block Acknowledgements.
- Clients can indicate going to sleep, and wake up any time:
 - Requesting new data using a Power Save Poll (PS-Poll) frame.



TKIP-Frame

Setting power-management flag in its Frame Control (FC) field.



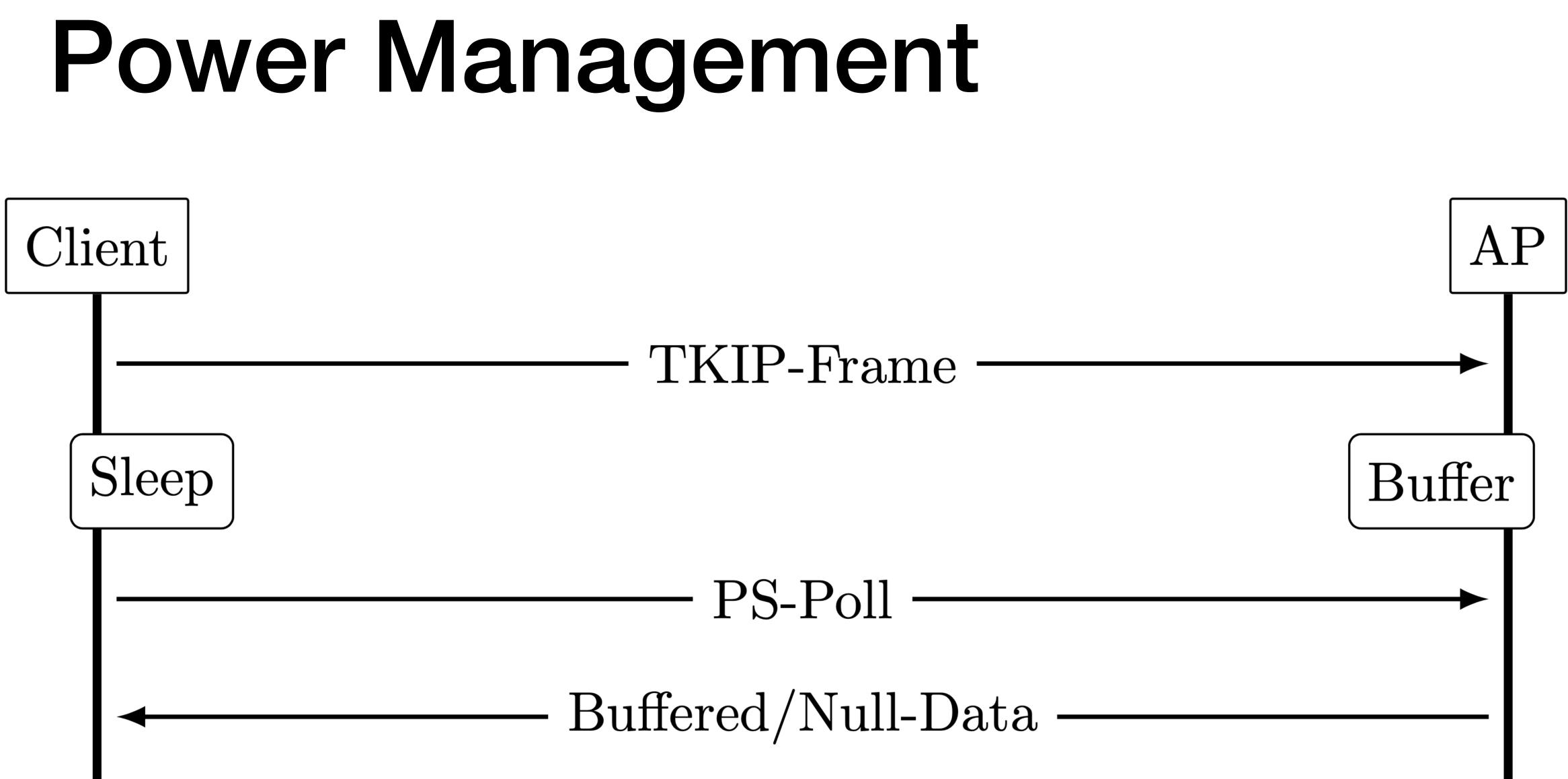


TKIP-Frame

Setting power-management flag in its Frame Control (FC) field.







Power Management

- Side-channels may abuse these power management features:
 - Spoofing target client to enter Power Save mode, then request data.
- Affects Linux kernel version 3.11 and lower.
 - Devices supporting WPA-TKIP are arguably running on older versions.
 - We will soon see another side-channel for newer kernel versions.

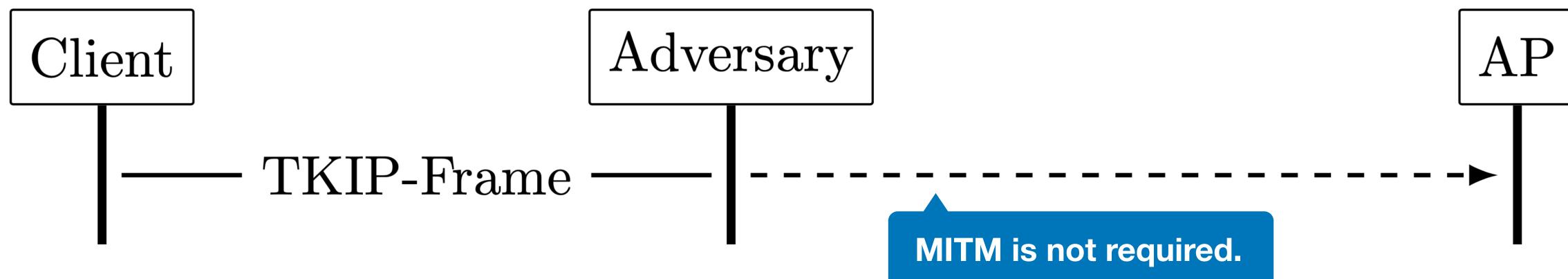
ieee80211_rx_result res = RX_DROP_MONITOR; struct sk_buff *skb;

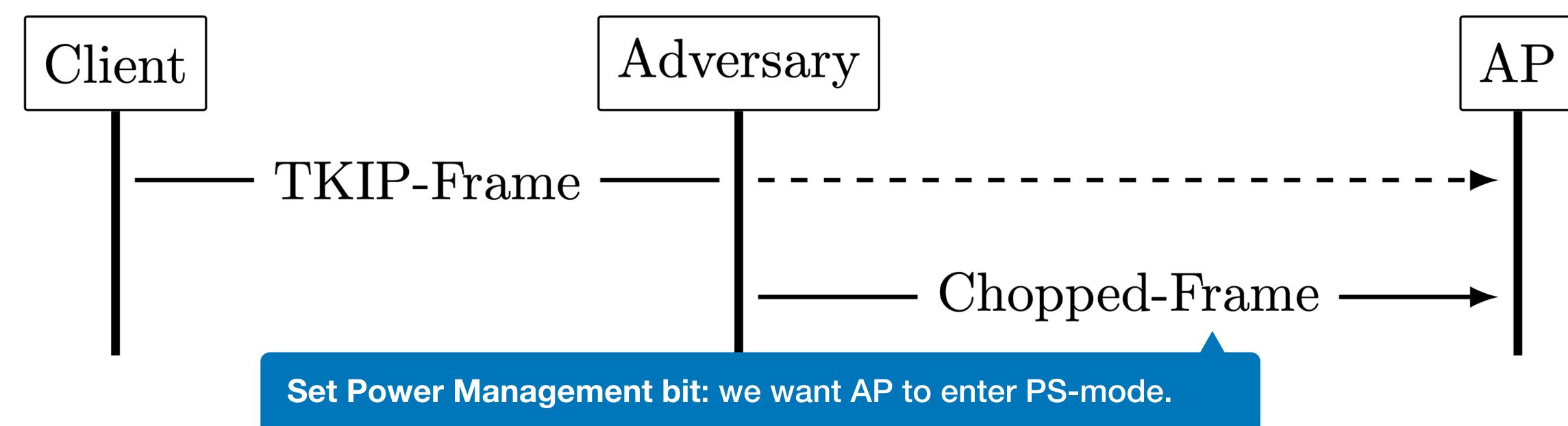
#define CALL_RXH(rxh) do { res = rxh(rx);if (res != RX CONTINUE)

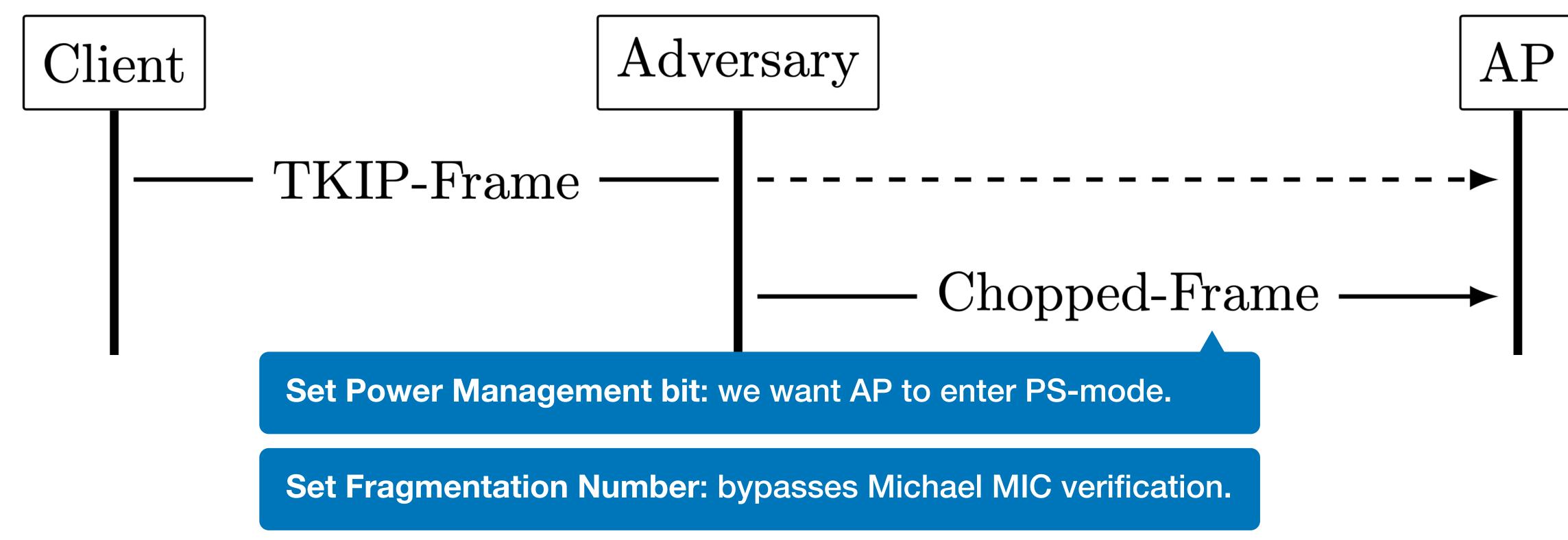
CALL_RXH(ieee80211_rx_h_decrypt) CALL_RXH(ieee80211_rx_h_check_more_data) CALL_RXH(ieee80211_rx_h_uapsd_and_pspoll) CALL_RXH(ieee80211_rx_h_sta_process) CALL_RXH(ieee80211_rx_h_defragment) CALL_RXH(ieee80211_rx_h_michael_mic_verify)

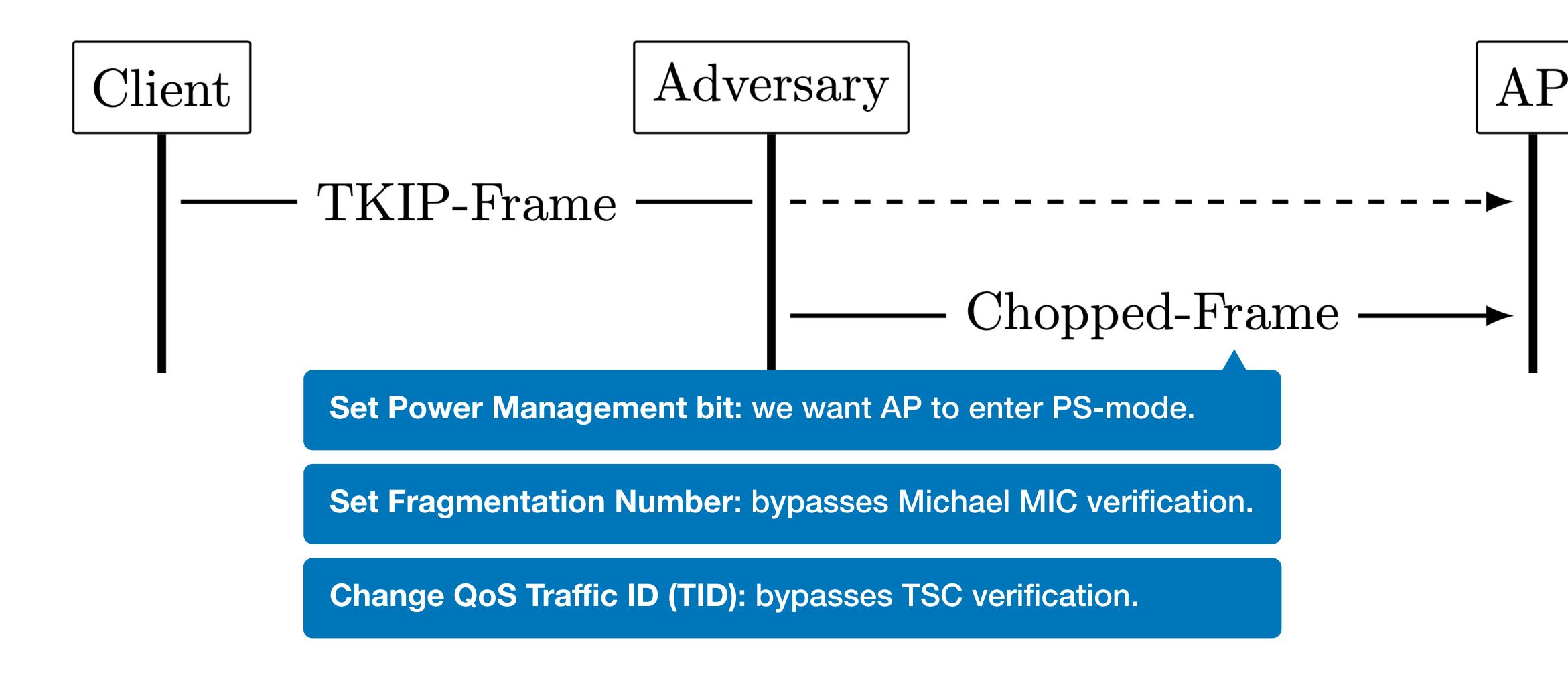
CALL_RXH(ieee80211_rx_h_check_more_data) CALL_RXH(ieee80211_rx_h_uapsd_and_pspoll) CALL_RXH(ieee80211_rx_h_sta_process) CALL_RXH(ieee80211_rx_h_defragment) CALL_RXH(ieee80211_rx_h_michael_mic_verify) /* must be after MMIC verify so header is counted in MPDU mic */ #ifdef CONFIG_MAC80211_MESH if (ieee80211_vif_is_mesh(&rx->sdata->vif)) CALL_RXH(ieee80211_rx_h_mesh_fwding); #endif

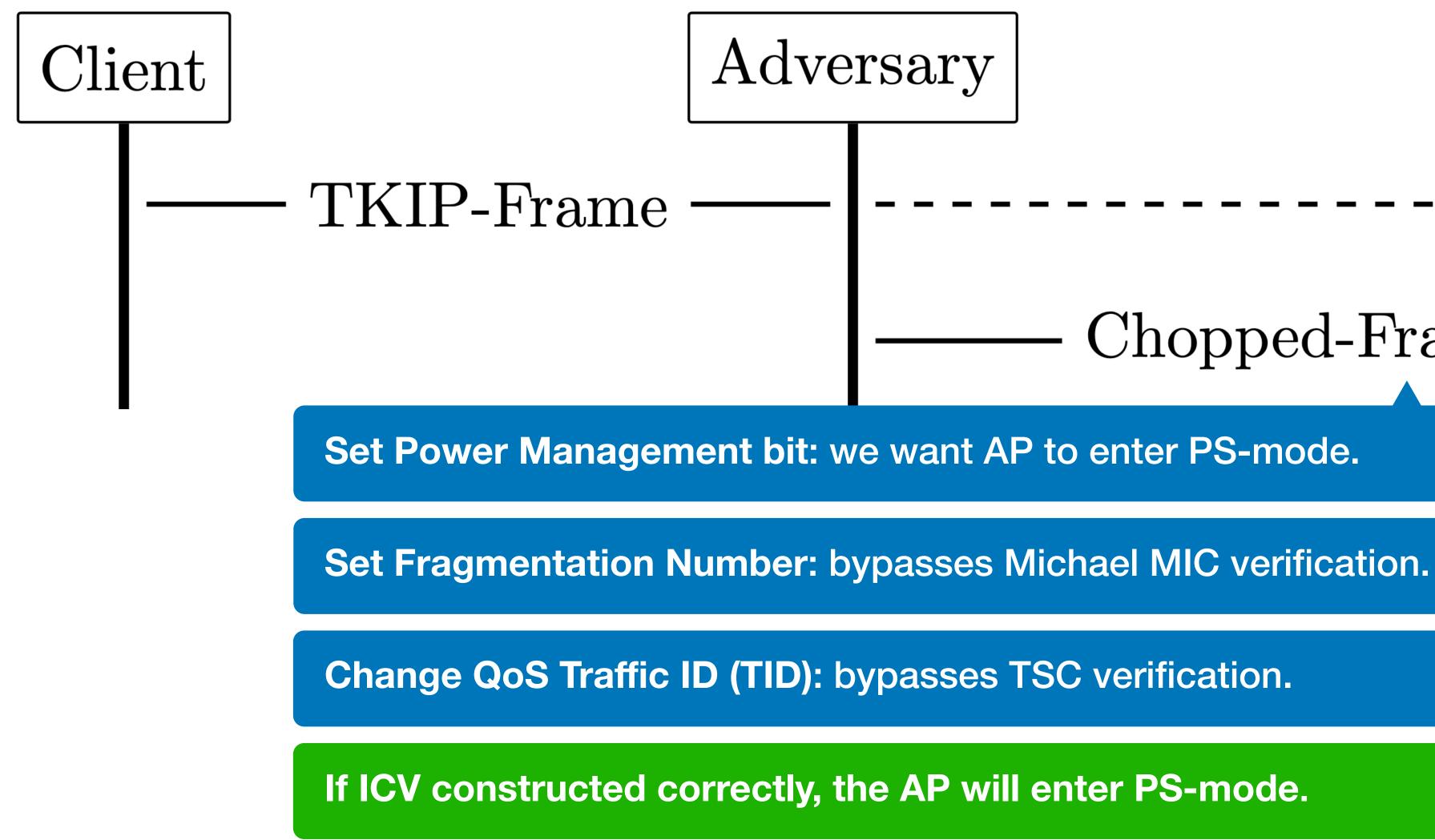
```
static void ieee80211_rx_handlers(struct ieee80211_rx_data *rx,
                                  struct sk_buff_head *frames)
```





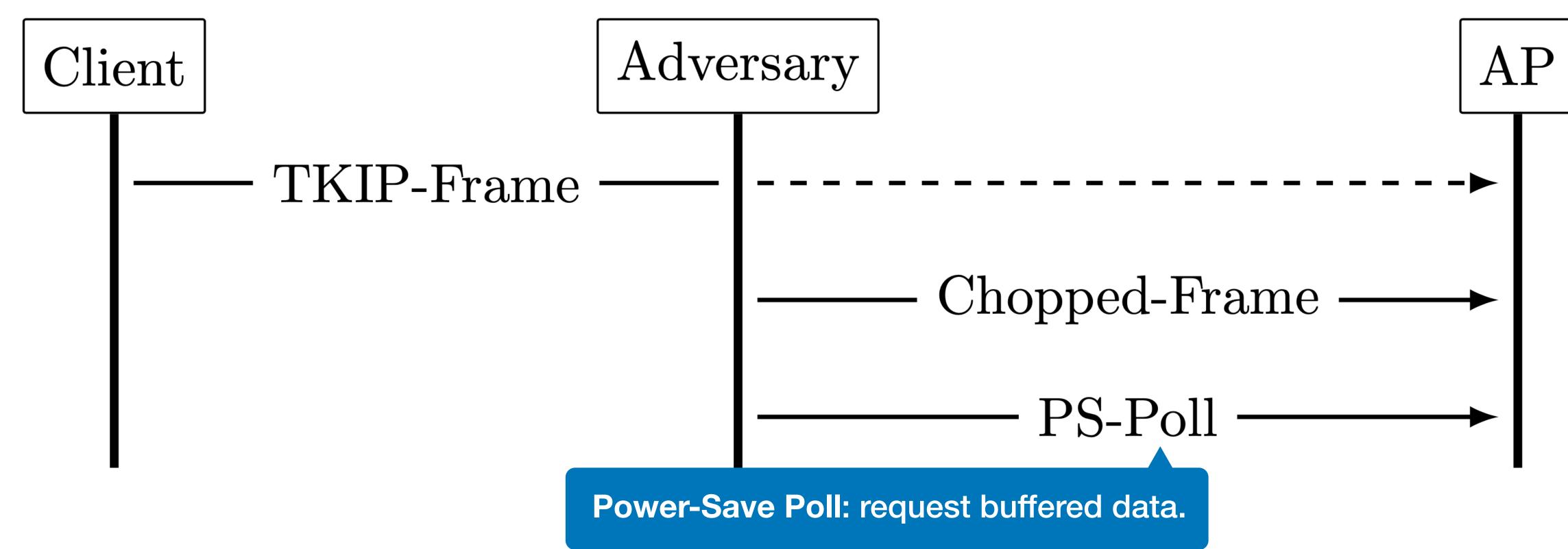


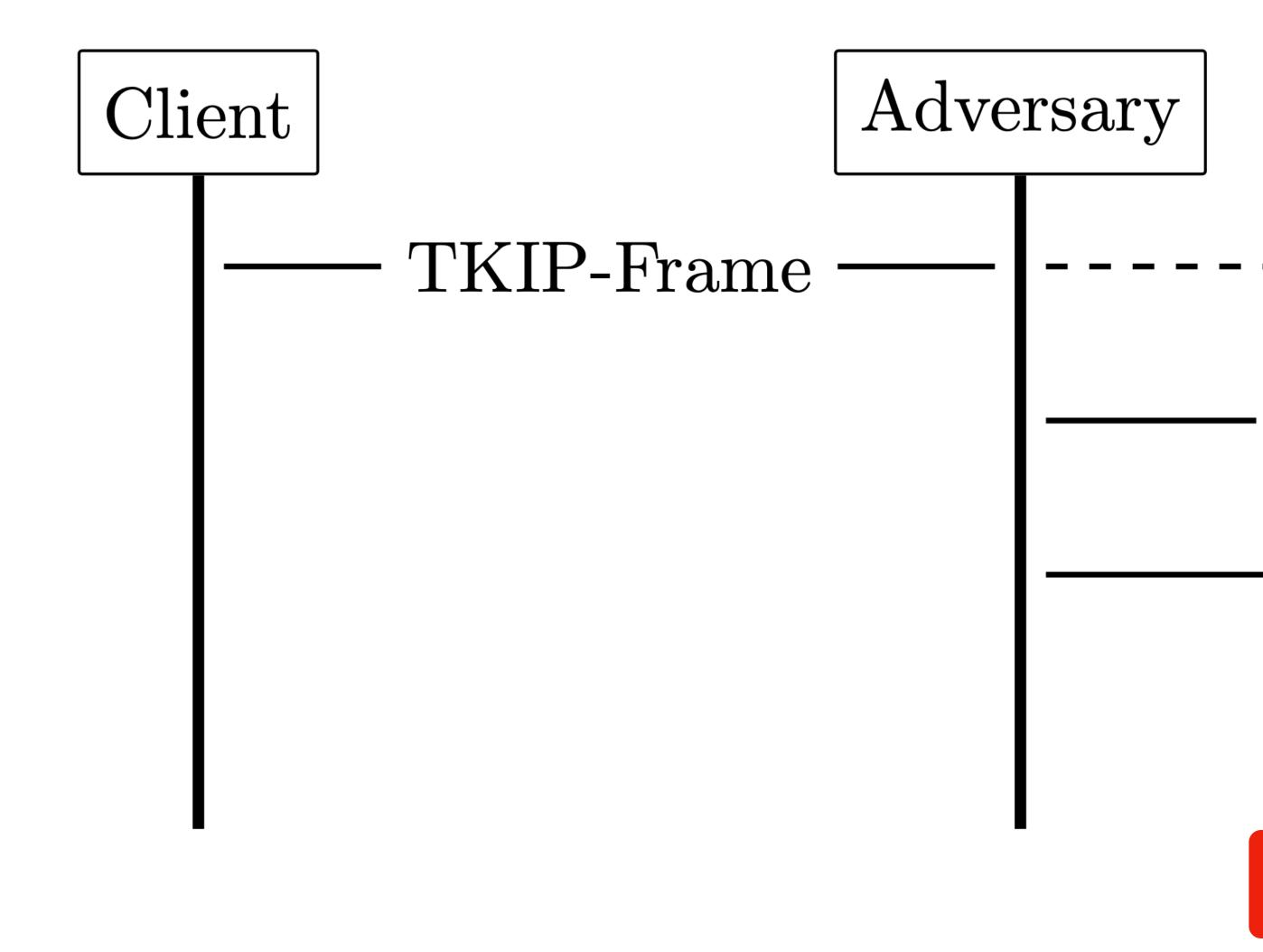




Chopped-Frame

AP





Chopped-Frame

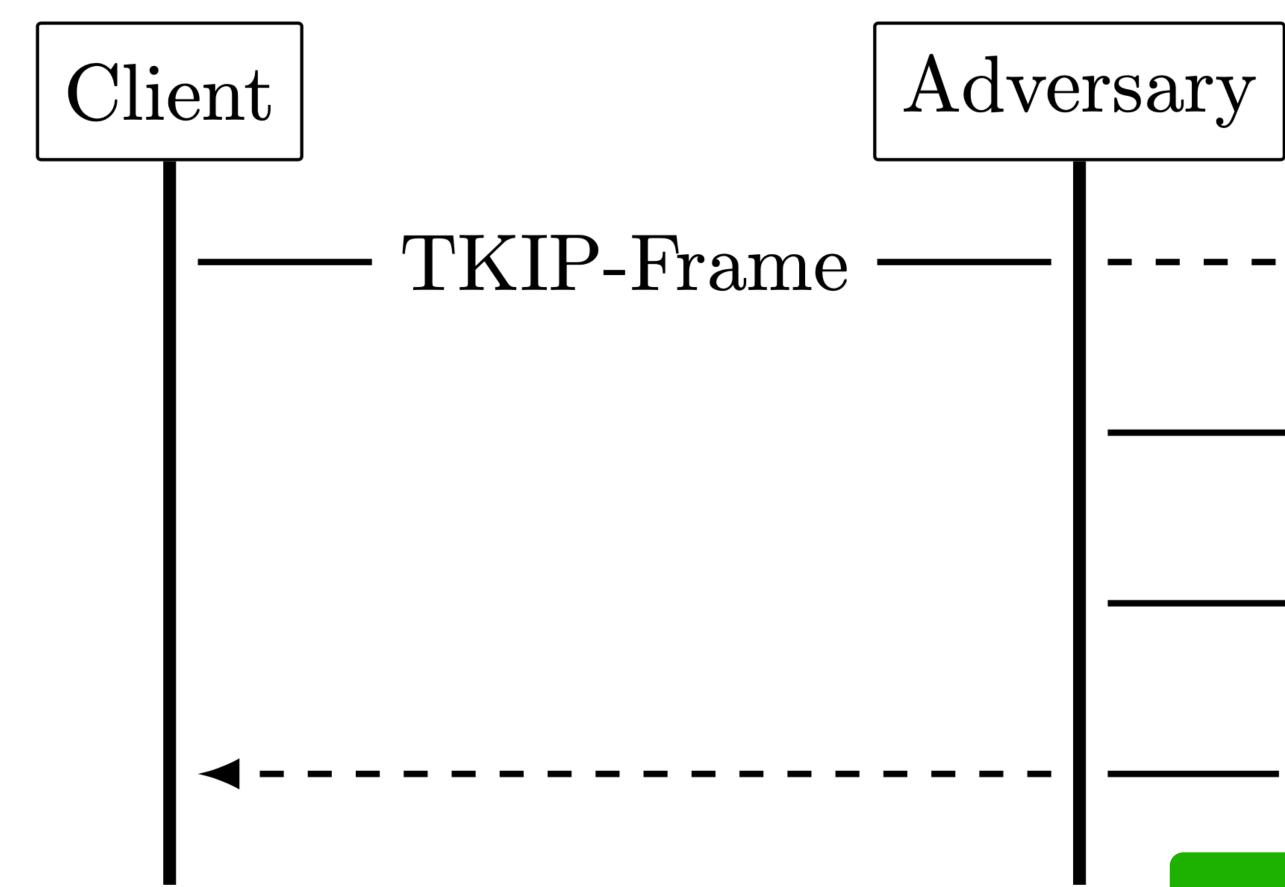
PS-Poll

No Reply

Failure: AP did not enter PS-mode.

AP

Power Management Side-Channel



Chopped-Frame

PS-Poll

Buffered/Null-Data

Success: AP entered PS-mode.



AP

	¥/>								Expression +
Apply a display filter < No. Time	Source	Destination		Protocol Le	nath I	nfo			
1 0.000000	02:00:00:00:01:00	Broadcast		802.11	New York Concerning Street	QoS Data, SN=45,	FN=0, Flags=.g	pT	
2 0.010055	02:00:00:00:01:00	Broadcast		802.11		QoS Data, SN=45,			
3 0.010811	02:00:00:00:01:00	02:00:00:00:00:00 (02	:00:00:00:00:00)	802.11	24	Power-Save poll,	Flags=	e 19	
4 0.010821	02:00:00:00:00:00	02:00:00:00:01:00		802.11	60 (QoS Null functio	n (No data), SM	N=4095, FN=15	, Flags=F.
1000 = Si ▼ Flags: 0x41	ta, Flags: .pT S Data (0x0028) eld: 0x8841 ersion: 0 ype: Data frame (2)	STA to DS via an AR (T							
	More Fragments: This i Retry: Frame is not be PWR MGT: STA will stay	s the last fragment ing retransmitted up	O DS: 1 From DS: 0	0) (0x1)					
0= 0= 0= 0= 0= 0= 0= .000 0000 0	More Fragments: This i Retry: Frame is not be PWR MGT: STA will stay More Data: No data buf Protected flag: Data i Order flag: Not strict 100 = Duration: 52 micr 202:00:00:00:00:00 (02	s the last fragment ing retransmitted up fered s protected ly ordered oseconds :00:00:00:00:00)	o DS: 1 From DS: 0	0) (0x1)					
0 = 0 = 0 = 0 = .1 = 0 = .000 0000 0011 0 Receiver address Transmitter addr Destination addr Source address:	More Fragments: This i Retry: Frame is not be PWR MGT: STA will stay More Data: No data buf Protected flag: Data i Order flag: Not strict 100 = Duration: 52 micr	<pre>s the last fragment ing retransmitted up fered s protected ly ordered oseconds :00:00:00:00:00) (02:00:00:00:01:00) f:ff:ff:ff) 0:00:00:01:00)</pre>	o DS: 1 From DS: 0	0) (0x1)					
	<pre>More Fragments: This i Retry: Frame is not be PWR MGT: STA will stay More Data: No data buf Protected flag: Data i Order flag: Not strict 100 = Duration: 52 micr 202:00:00:00:00:00:00(02) 02:00:00:00:01:00 ess: Broadcast (ff:ff:f 02:00:00:00:01:00 (02:0) 00:00:00:01:00 (02:00) 00:00:00:01:00 (02:00) 00:00:00:01:00 (02:00) 00:00:00:01:00 (02:00) 000 = Fragment number:</pre>	<pre>s the last fragment ing retransmitted up fered s protected ly ordered oseconds :00:00:00:00:00) (02:00:00:00:01:00) f:ff:ff:ff) 0:00:00:01:00) 0:00:01:00) 0</pre>	o DS: 1 From DS: 0	0) (0x1)					
0= 0= 0= 0= .1= 0= .000 0000 0011 0 Receiver address Transmitter addr Destination addr Source address: BSS Id: 02:00:00 STA address: 02: 0	<pre>More Fragments: This i Retry: Frame is not be PWR MGT: STA will stay More Data: No data buf Protected flag: Data i Order flag: Not strict Ode Duration: 52 micr 02:00:00:00:00:00:00:00 02:00:00:00:01:00 Cess: Broadcast (ff:ff:f 02:00:00:00:01:00 (02:00:00:00:00) 00:00:00:01:00 (02:00:00:00:00:00) 00:00:00:01:00 (02:00:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:01:00 (02:00:00) 00:00:00:00:00:00:00:00 00:00:00:00:00:</pre>	<pre>s the last fragment ing retransmitted up fered s protected ly ordered oseconds :00:00:00:00:00) (02:00:00:00:01:00) f:ff:ff:ff) 0:00:00:01:00) 0:00:01:00) 0</pre>	o DS: 1 From DS: 4	0) (0x1)					

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	2 0.010055	02:00:00:00:01:00	Broadcast		802.11		QoS Data, SN=4			ar 	
	3 0.010811 4 0.010821	02:00:00:00:00:00:00 02:00:00:00:00:00	02:00:00:00:00:00 (0) 02:00:00:00:01:00	/2:00:00:00:00:00)	802.11		Power-Save pol QoS Null funct			5, FN=15, Flag	s=F.
Radio 802.11 IEEE 8 Typ V Fra	<pre>tap Header v0, 1 radio informa 802.11 QoS Data be/Subtype: QoS ame Control Fie 10 = Typ 1000 = Sub Flags: 0x51</pre>	tion , Flags: .p.PT Data (0x0028) ld: 0x8851 sion: 0 e: Data frame (2) type: 8									
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Rec Tra Des Sou BSS	<pre></pre>	More Fragments: This i Retry: Frame is not be PWR MGT: STA will go t More Data: No data buf Protected flag: Data i Order flag: Not strict 00 = Duration: 52 micro 02:00:00:00:00:00 (02	<pre>s the last fragment ing retransmitted o sleep fered s protected ly ordered oseconds :00:00:00:00:00) (02:00:00:00:01:00) f:ff:ff:ff) 0:00:00)</pre>			er m	odificati	ons to	replaye	ed frame	

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4 0.010821		:00:00:00:00:00:00 (02:00:00				on (No data), SN=4095,	FN=15, Flags=F.
 Frame 3: 24 bytes o Radiotap Header v0, 	n wire (192 bits), 24 byte Length 8	s captured (192 bits)					
802.11 radio inform							
	ave poll, Flags: ver-Save poll (0x001a)						
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🔵 🎽 IEEE 802.11 wireless	LAN (wlan), 16 bytes				Packets: 4 · [Displayed: 4 (100.0%)	Profile: Default

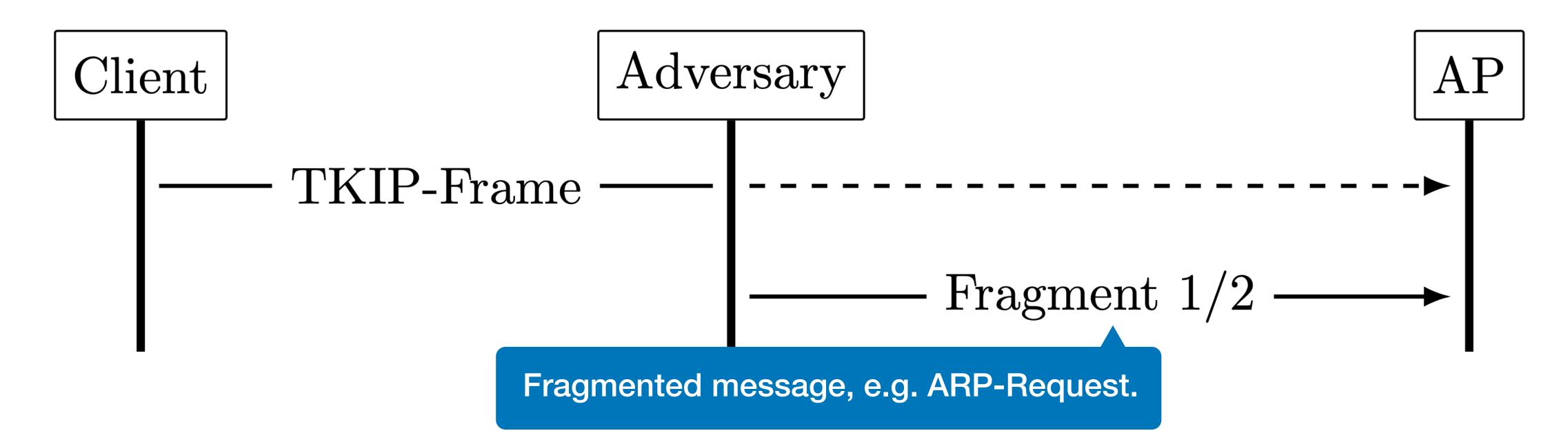
o o o	
📘 Apply a display filter <郑/> Expression	+
No. Time Source Destination Protocol Length Info	
1 0.000000 02:00:00:00:01:00 Broadcast 802.11 116 QoS Data, SN=45, FN=0, Flags=.pT	
2 0.010055 02:00:00:00:01:00 Broadcast 802.11 90 QoS Data, SN=45, FN=1, Flags=.p.PT 3 0.010811 02:00:00:00:01:00 … 02:00:00:00:00:00:00 (02:00:00:00:00)… 802.11 24 Power-Save poll, Flags=	
4 0.010821 02:00:00:00:00:00:00:00:00:00:00:00:00:0	- -
► Frame 4: 60 bytes on wire (480 bits), 60 bytes captured (480 bits) ► Padiatap Header v@ _longth 34	
Radiotap Header v0, Length 34 802.11 radio information	
▼ IEEE 802.11 QoS Null function (No data), Flags:F.	
Type/Subtype: QoS Null function (No data) (0x002c)	
▼ Frame Control Field: 0xc802	
00 = Version: 0	
10 = Type: Data frame (2)	
1100 = Subtype: 12	
▼ Flags: 0x02	
0 = Retry: Frame is not being retransmitted	
0 = PWR MGT: STA will stay up	
= More Data: No data buffered	
.0 = Protected flag: Data is not protected	
0 = Order flag: Not strictly ordered	
.000 0000 0011 0100 = Duration: 52 microseconds	
Receiver address: 02:00:00:00:01:00 (02:00:00:00:01:00) Transmitter address: 02:00:00:00:00:00:00:00:00:00:00:00:00:0	
Destination address: 02:00:00:00:00:00:00:00:00:00:00:00:00)	
Source address: 02:00:00:00:00:00:00 (02:00:00:00:00) BSS Id: 02:00:00:00:00:00 (02:00:00:00:00:00)	
STA address: 02:00:00:00:00:00:00:00:00:00:00:00:00:0	
1111 = Fragment number: 15	
1111 1111 1111 = Sequence number: 4095	
▶ Qos Control: 0x0007	

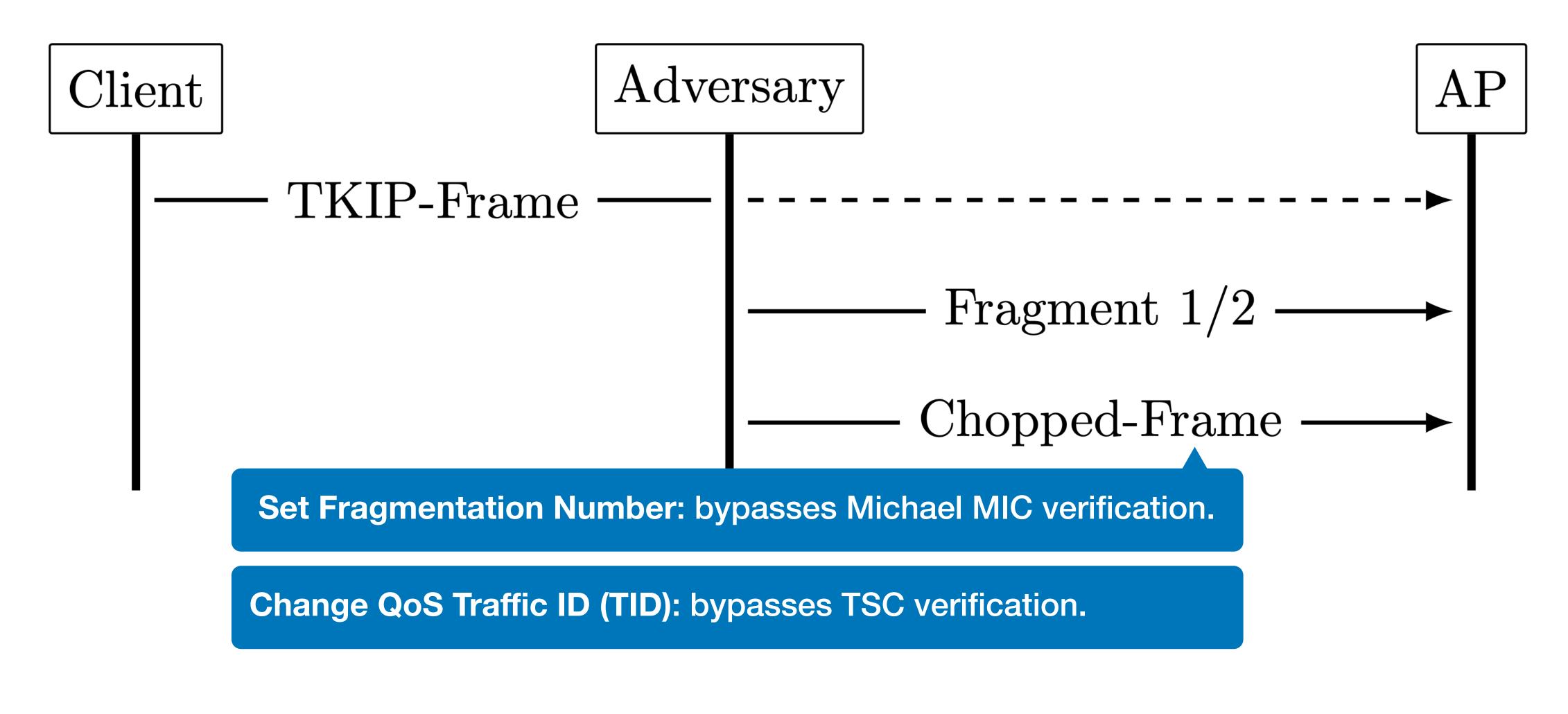
Power Management

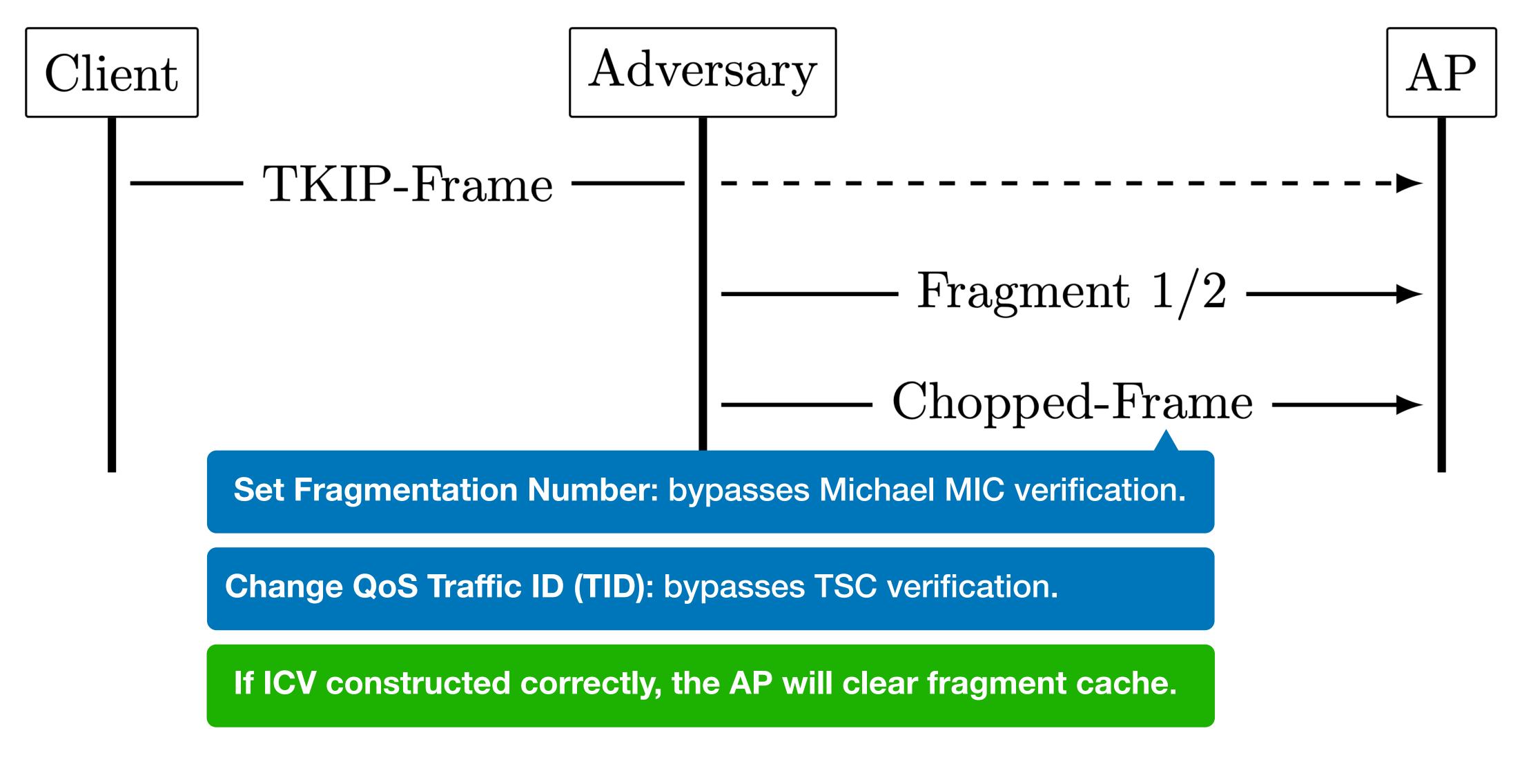
- Bypasses all existing TKIP Countermeasures.
 - E.g. Michael MIC Failure Reports, increased rekeying interval.
- Recovers the Michael MIC key (message authentication) in 1 to 4 minutes.
 - Much faster than previous attacks, needing 7 to 8 minutes.
 - Due to not relying on TKIP's Michael MIC Failure Reports.

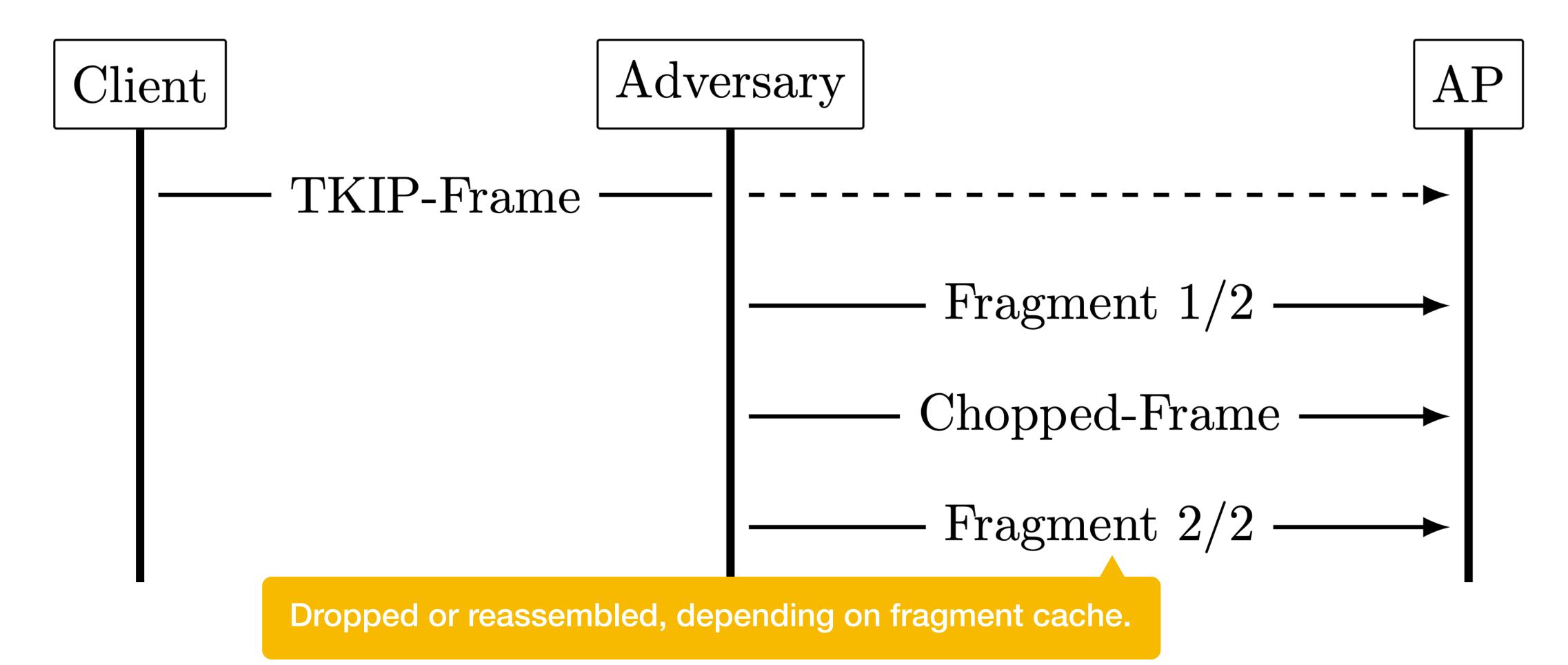
Fragmentation

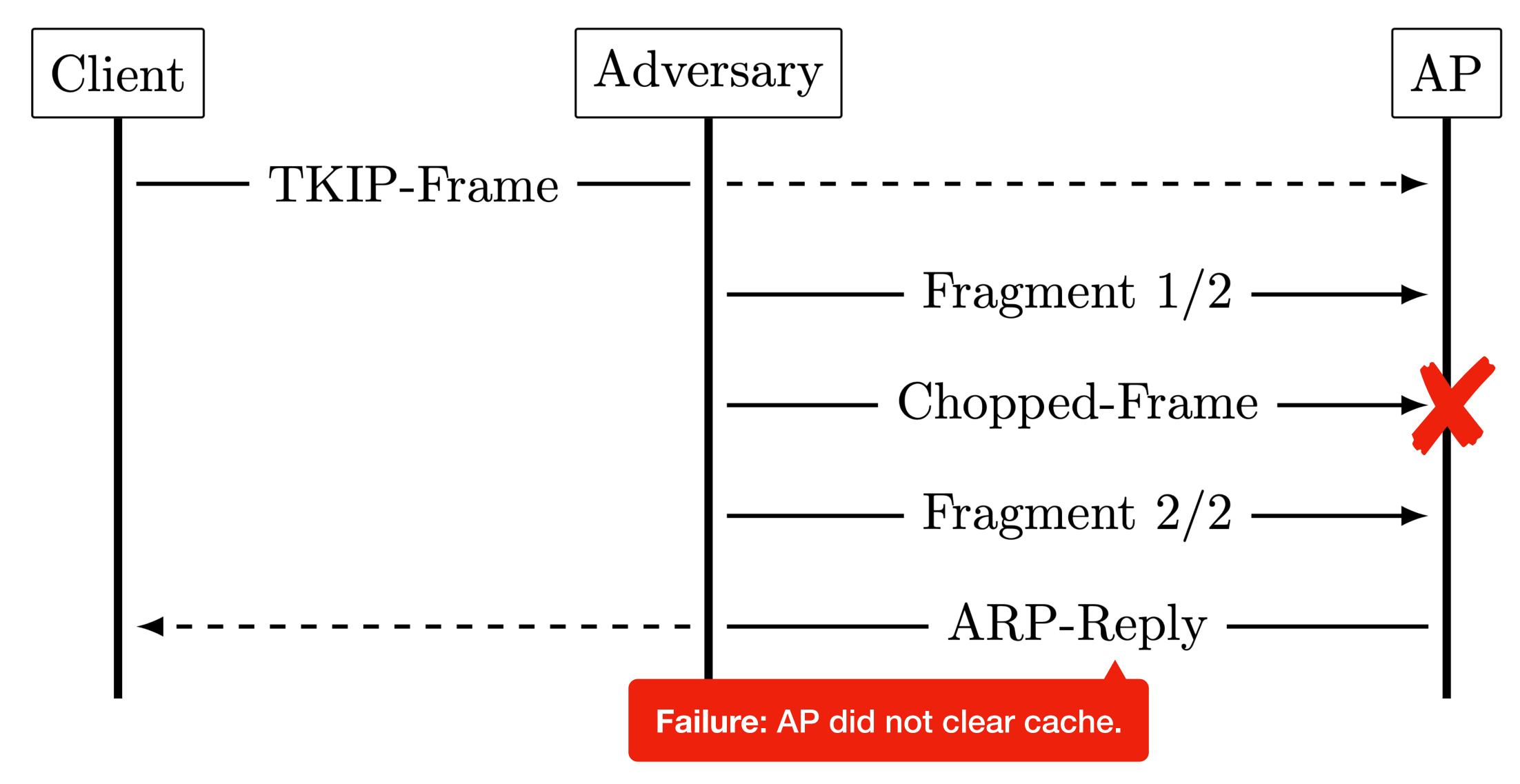
- Side-Channel abuses fragmentation features in MediaTek devices.
- Fragmented WPA-TKIP frame with valid ICV clears fragmentation cache.
 - If the ICV is wrong, the frame is silently dropped (cache unaffected).
 - If the ICV is correct, the cache is cleared (cache affected).
- Proof-of-Concept with fragmented ARP-Request.

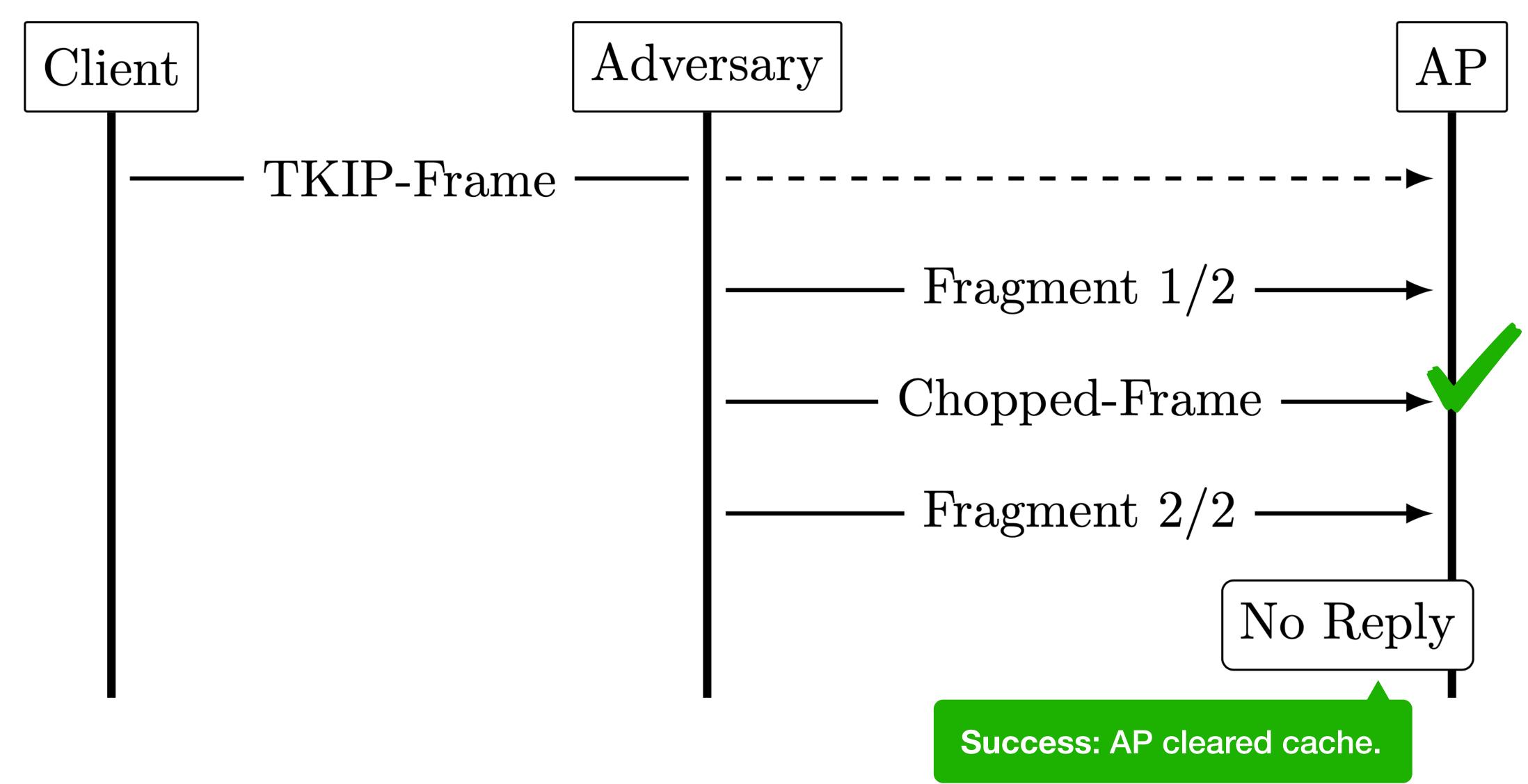










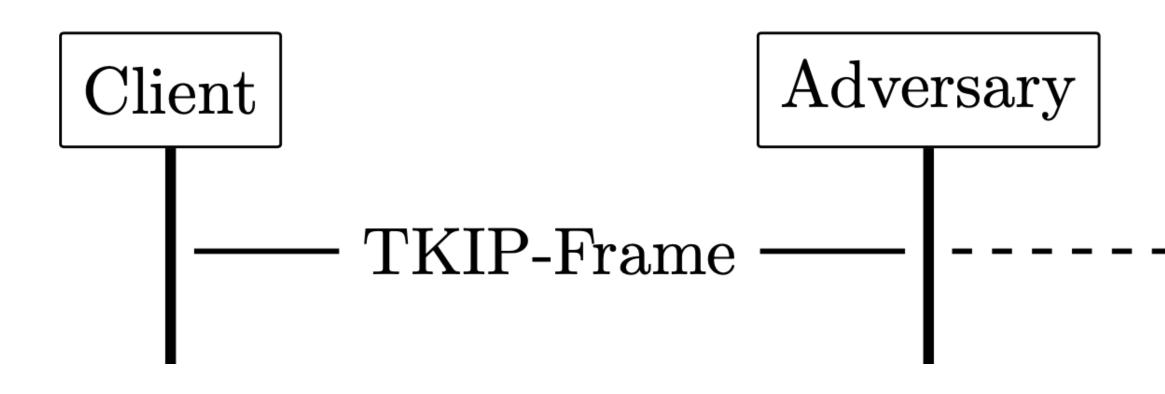


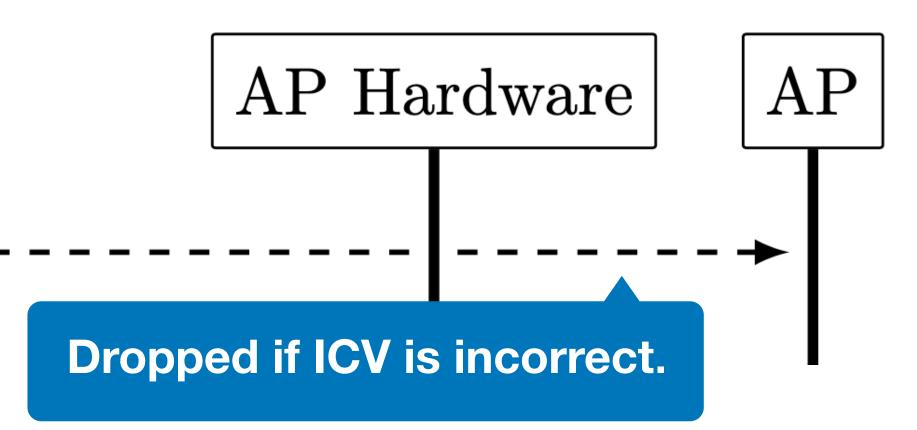
Hardware Decryption

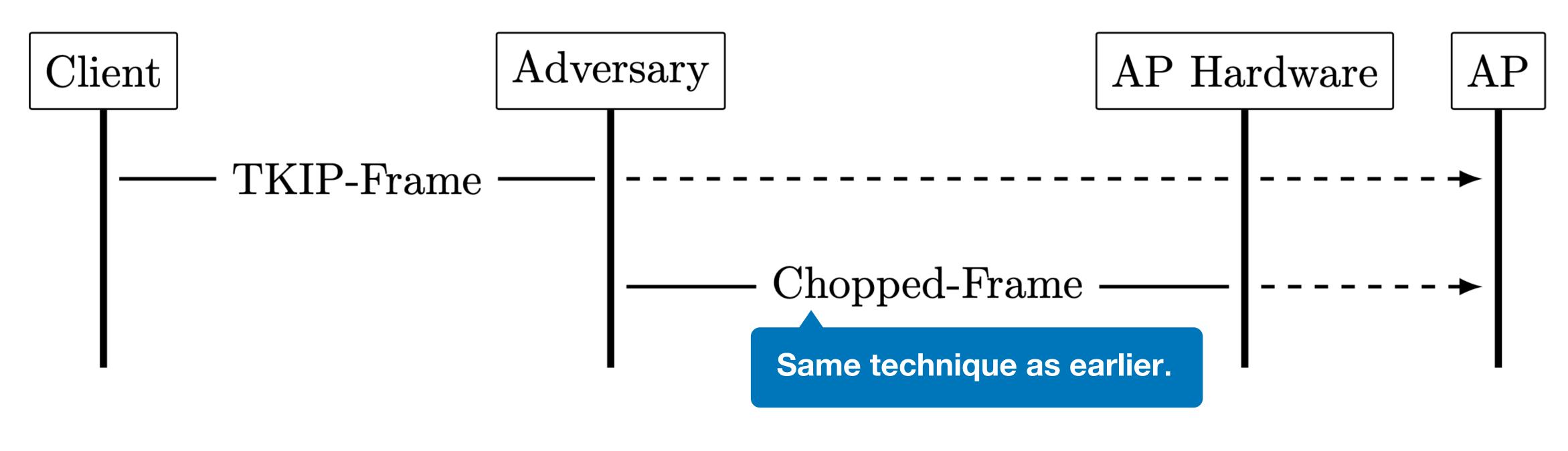
- Investigate the hardware decryption features of Wi-Fi chips.
- Some drivers drop incoming frames if hardware decryption failed.
 - For example, due to an incorrect ICV.
 - Frames only forwarded if ICV is correct, potentially leaking information.

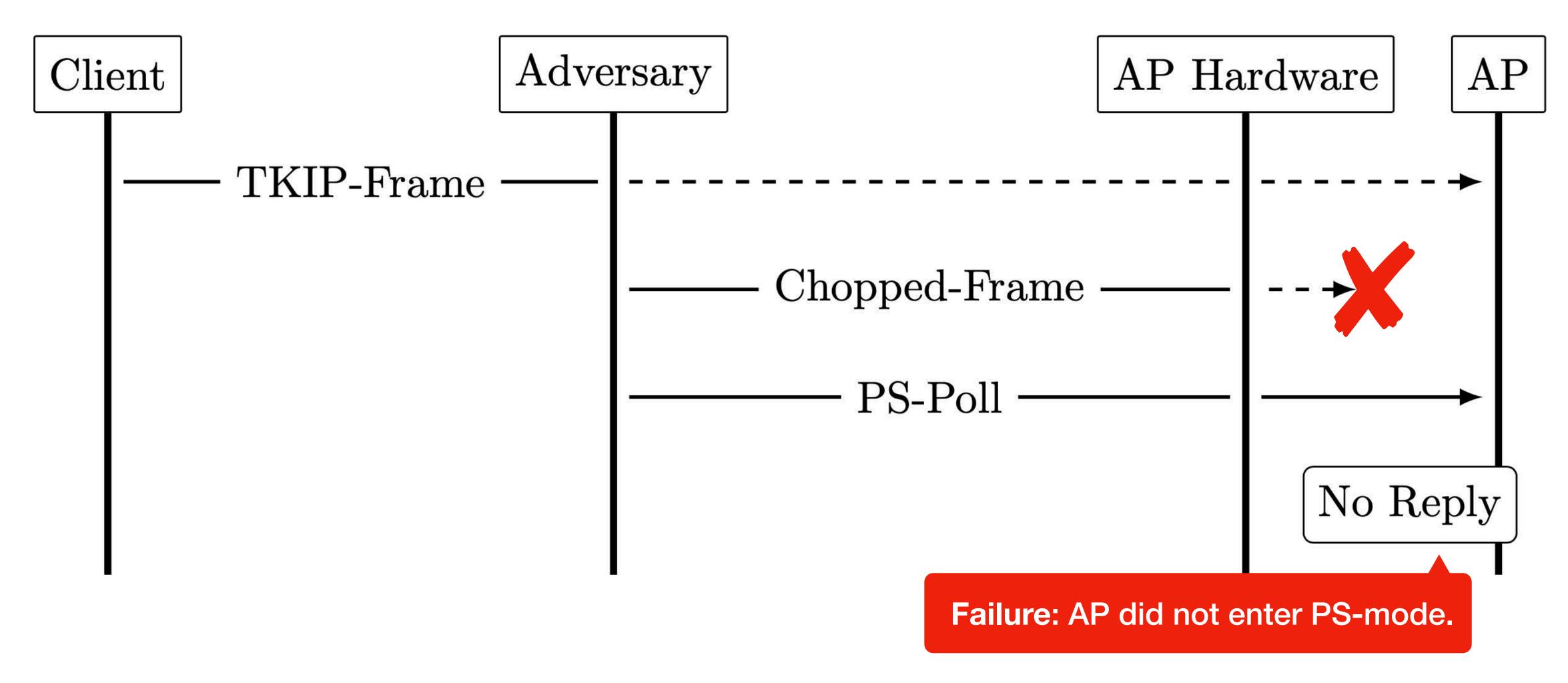
Return of Power Management

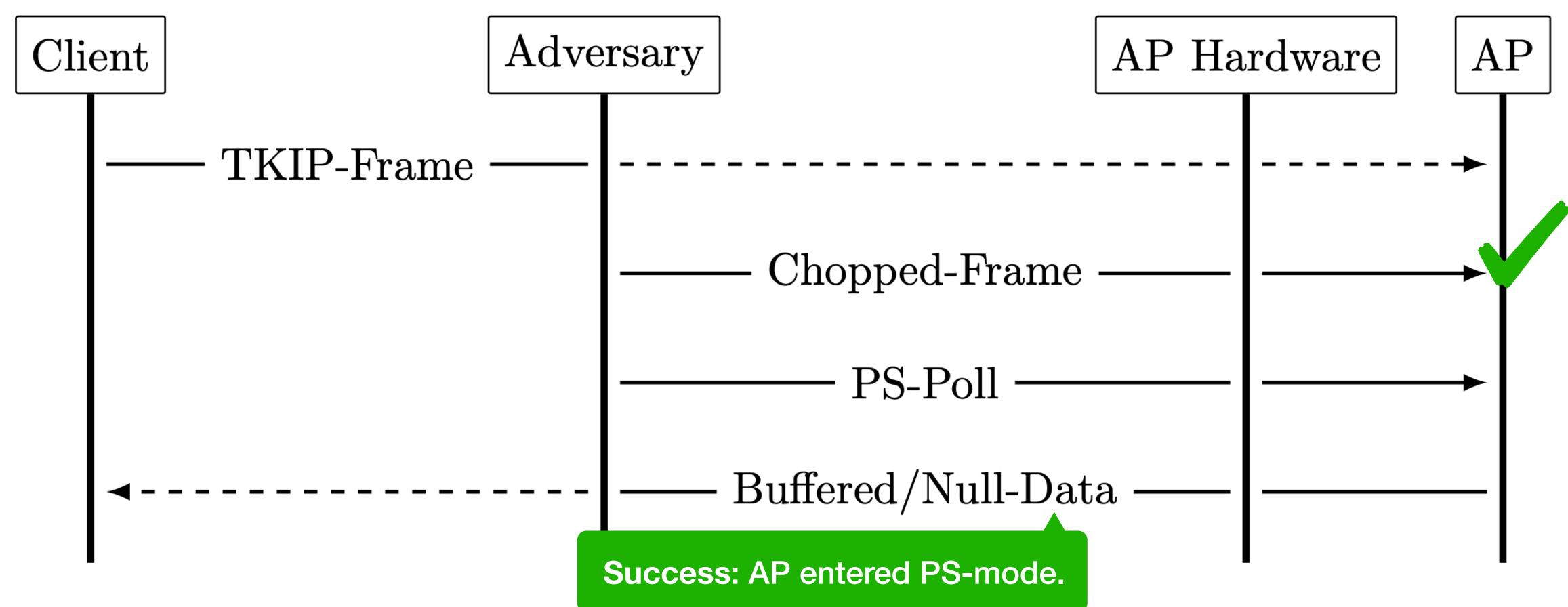
- Hardware decryption performed before OS handles power management.
 - Flags inspected only when hardware decryption succeeded.
- Works against Linux 3.12 and higher if hardware decryption is enabled.
 - Unlike previous side-channel, which worked on Linux 3.11 and lower.

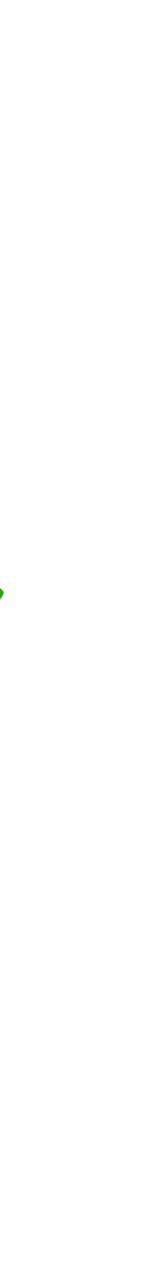












Hardware Decryption

- Investigation reveals several new (similar) side-channels and attacks.
 - Linux Power Management, Retransmission Detection.
 - OpenBSD Block Acknowledgement, Fragment and Replay Attack.
- See our white paper for more details.

Let's Summarize:

- We can decrypt frames.

- Numerous Side-Channels.

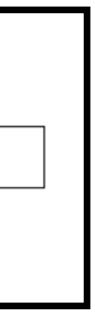
Cool! Now, what is the impact?

Recall '	TKIP			de	r			
	Slide 9		4	4	variable	8	4	4
	MAC Header	Sequece No.	IV/KeyID	Ext. IV	Data	Michael MIC	ICV	FCS
					<	— Encrypted ———	>	

- We now have a means for decrypting one byte at a time.

Decrypting the last 12-bytes of a frame reveals the Michael MIC and ICV.

• Recovers Michael MIC key since the algorithm is invertible (Wool, 2004).



We want more keystream!

- Allows us to inject messages.
- Michael MIC recovered after first message.
 - Then, an adversary needs to decrypt the last 4-bytes (ICV) only.
- Quickly learn keystream by making educated guesses for data contents.
 - E.g., most bytes of an ARP-Reply are known thus brute-forced quickly.

Impact

- We have now obtained keystream and the Michael MIC key.
- Allows us to decrypt, and inject, network traffic:
 - The Michael MIC key can be used for decryption using a Michael Reset Attack (Vanhoef et al, 2013).
 - Forged messages are injected into QoS channels with a lower TSC (Tews et al, 2009; Vanhoef et al, 2013).

How is this an improvement?

Improvements

- We bypass all existing TKIP Countermeasures.
- Consequently we do not rely on any of them either.
 - E.g., related work relies on Michael MIC Failure Reports.
- As a result, our decryption techniques are significantly faster.
 - Recovery of Michael message authentication key in 1 to 4 minutes.

Countermeasures to our Attacks

- Our side-channel attacks bypass all existing countermeasures.
 - E.g., short rekeying interval, delayed or disabled failure reports.
- We can further reduce rekeying interval, or try to fix an outdated protocol:
 - E.g., software changes to process reassembled frames only.

Countermeasures to our Attacks

- Our side-channel attacks bypass all existing countermeasures.
 - E.g., short rekeying interval, delayed or disabled failure reports.
- We can further reduce rekeying interval, or try to fix an outdated protocol:
 - E.g., software changes to process reassembled frames only.
- Or... let us simply disable WPA-TKIP and move on to better protocols.

What about WPA2-CCMP?

CTR mode with CBC-MAC Protocol

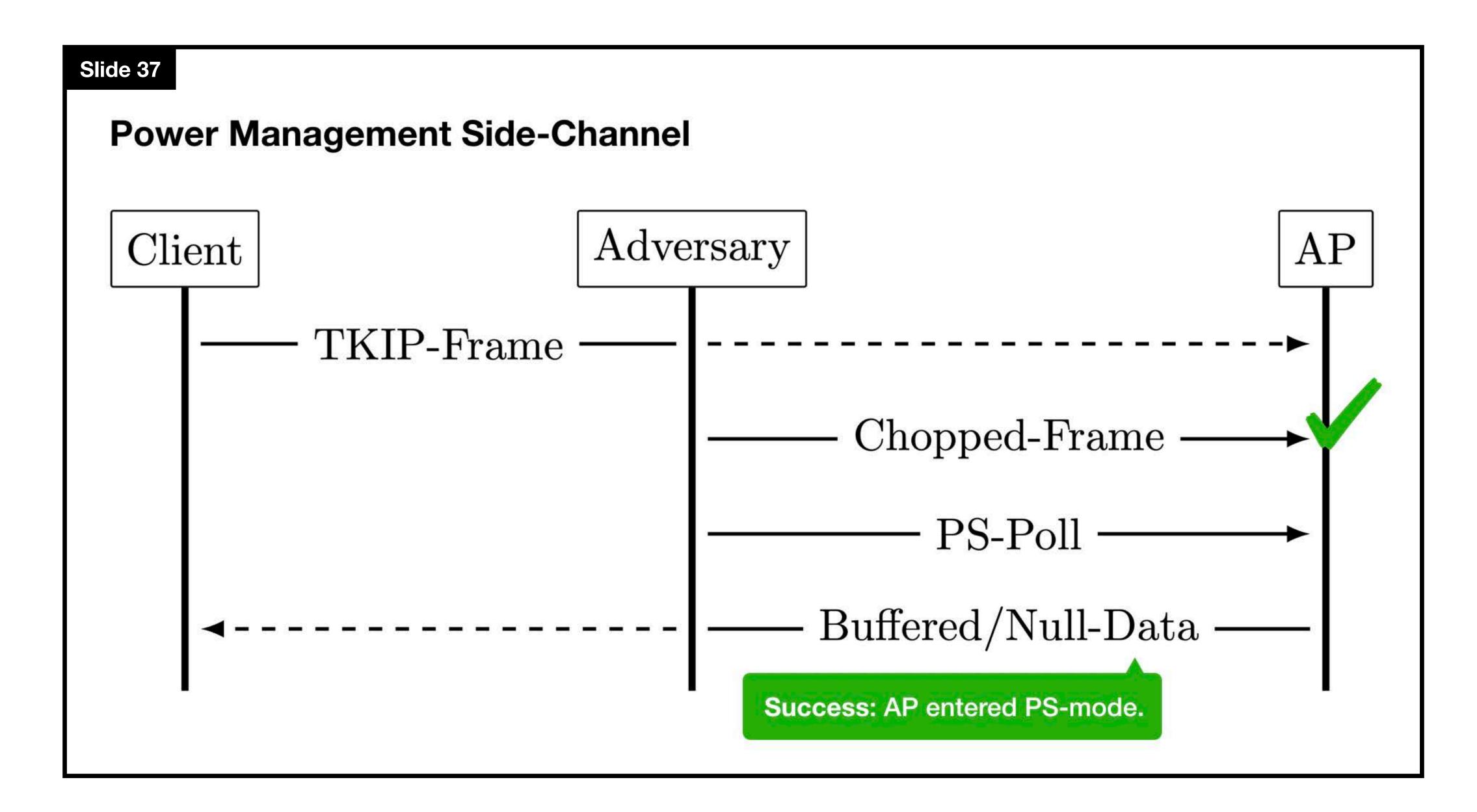
- WPA2 with CCMP uses an AES block cipher in counter mode.
- CCMP MIC (CBC-MAC) is calculated over every frame.
 - Unlike TKIP MIC, where Michael MIC is transmitted in the last fragment.
- CCMP MIC is not known to be invertible.
 - Unlike TKIP's Michael MIC (Wool, 2004).
- CCMP has no ICV.

CTR mode with CBC-MAC Protocol

- Immediately verifies for data integrity and data origin authentication.
 - Unlike TKIP, which for example carries Michael MIC in last fragment.
- Our side-channels work against WPA-TKIP because:
 - Operations are performed on decrypted but unauthenticated data.
 - Decrypted frames are processed before being reassembled.

Demonstration.

https://github.com/domienschepers/asiaccs-practical-side-channel

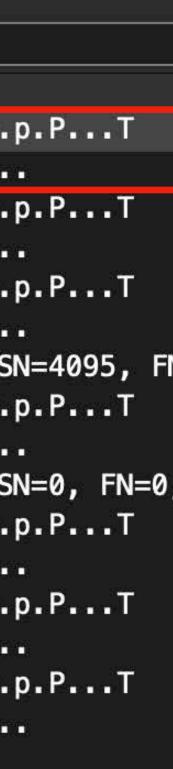


😣 🗐 🗊 🛛 root@ubuntu: /bheu

```
root@ubuntu:/bheu# python demo.py
[+] Sniffing for Target 02:00:00:00:01:00 and BSSID 02:00:00:00:00:00 on Interface wlan2.
[+] Capturing a TKIP Frame...
[+] 802.11 Data 8 02:00:00:00:01:00 > 02:00:00:00:00:00 / Dot11QoS / Dot11TKIP
[+] Received a Null Data Frame for guess 163 (Plaintext 0xa3).
[+] Received a Null Data Frame for guess 221 (Plaintext 0x6e).
[+] Received a Null Data Frame for guess 202 (Plaintext 0xa2).
[+] Received a Null Data Frame for guess 49 (Plaintext 0x96).
[+] Received a Null Data Frame for guess 71 (Plaintext 0x73).
[+] Received a Null Data Frame for guess 74 (Plaintext 0xa0).
[+] Received a Null Data Frame for guess 186 (Plaintext 0x87).
[+] Received a Null Data Frame for guess 85 (Plaintext 0xc2).
[+] Received a Null Data Frame for guess 5 (Plaintext Oxaa).
[+] Received a Null Data Frame for guess 173 (Plaintext 0xd9).
[+] Received a Null Data Frame for guess 28 (Plaintext 0xc4).
[+] Received a Null Data Frame for guess 15 (Plaintext 0x7c).
[+] Recovered the plaintext of the ARP Message for guess #111.
###[ LLC ]###
 dsap
           = 0xaa
           = 0xaa
 ssap
 ctrl
           = 3
###[ SNAP ]###
              = 0 \times 0
    OUI
              = ARP
    code
###[ ARP ]###
                 = 0 \times 1
       hwtype
                 = IPv4
        ptype
       hwlen
                 = None
        plen
                 = None
                 = who-has
        ор
        hwsrc
                 = 02:00:00:00:01:00
                 = 192.168.0.10
        psrc
        hwdst
                 = ff:ff:ff:ff:ff
        pdst
                  = 192.168.0.111
root@ubuntu:/bheu#
```

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No.		Time	Source		Destination					Protoc	ol Len		•		
	1	0.00000	02:00:00:00:01	:00	Broadcast					802.1	1	79	QoS Data, SN=4	2, FN=1,	, Flags=.p
	2	0.001077	02:00:00:00:01	:00	02:00:00:0	00:00:00	(02:00:	:00:00	:00:00)	. 802.1	.1		Power-Save pol		
	3	0.053206	02:00:00:00:01	:00	Broadcast					802.1	1	79	QoS Data, SN=4	2, FN=1,	, Flags=.p
	4	0.054175							00:00)	. 802.1	.1		Power-Save pol		
	5	0.107561	Failure: clier	nt no	t register	red in	PS-mo	ode.		802.1	1		QoS Data, SN=4		
	6	0.108973			U				00:00)	. 802.1	.1		Power-Save pol		
	7	0.109012	02:00:00:00:00	:00	02:00:00:0	00:01:00				802.1	1	60	QoS Null funct	ion (No	data), SN
	8	0.114788	02:00:00:00:01	:00	Broadcast					802.1	.1	79	QoS Data, SN=4	2, FN=1	, Flags=.p
	9	0.115859	02:00:00:00:01	:00	02:00:00:0	00:00:00	(02:00:	:00:00	:00:00)	. 802.1	1	24	Power-Save pol	l, Flag	s=
	10	0.115871	02:00:00:00:00	:00	02:00:00:0	00:01:00	l.			802.1	.1	60	QoS Null funct	ion (No	data), SN
	11	0.122050	02:00:00:00:01	:00	Broadcast					802.1	1	78	QoS Data, SN=4	2, FN=1,	, Flags=.p
	12	0.123099	02:00:00:00:01	:00	02:00:00:0	00:00:00	(02:00:	:00:00	:00:00)	. 802.1	1	24	Power-Save pol	l, Flag	s=
	13	0.174995	02:00:00:00:01	:00	Broadcast					802.1	1	78	QoS Data, SN=4	2, FN=1,	, Flags=.p
	14	0.176063	02:00:00:00:01	:00	02:00:00:0	00:00:00	(02:00:	:00:00	:00:00)	. 802.1	1	24	Power-Save pol	l, Flag	s=
	15	0.227746	02:00:00:00:01	:00	Broadcast					802.1	1	78	QoS Data, SN=4	2, FN=1,	, Flags=.p
	16	0.228729	02:00:00:00:01	:00	02:00:00:0	00:00:00	(02:00)	:00:00	:00:00)	. 802.1	1	24	Power-Save pol	l, Flag	s=
			\sim $($.6						

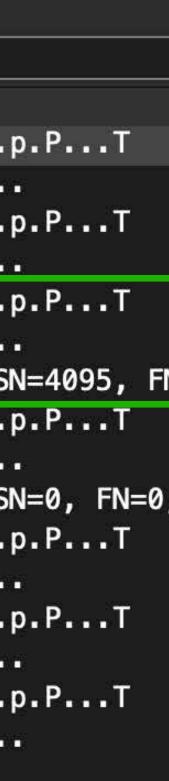
Frame 1: 79 bytes on wire (632 bits), 79 bytes captured (632 bits) Radiotap Header v0, Length 8 802.11 radio information IEEE 802.11 QoS Data, Flags: .p.P...T ▶ Data (37 bytes)



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📘 Apply	a display filter	. <郑/>				
No.	Time	Source	Destination	Pro	otocol Length	Info
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	2 0.001077	02:00:00:00:01:00	02:00:00:00:00:00 (0	02:00:00:00:00:00) 80	2.11 24	Power-Save poll, Flags=
	3 0.053206	02:00:00:00:01:00	Broadcast	803	2.11 79	QoS Data, SN=42, FN=1, Flags=.p
	4 0.054175	02:00:00:00:01:00	02:00:00:00:00:00 (0	02:00:00:00:00:00) 80	2.11 24	Power-Save poll, Flags=
	5 0.107561	02:00:00:00:01:00	Broadcast	803	2.11 79	QoS Data, SN=42, FN=1, Flags=.p
	6 0.108973	02:00:00:00:01:00	02:00:00:00:00:00 (0	02:00:00:00:00:00) 80	2.11 24	Power-Save poll, Flags=
	7 0.109012	02:00:00:00:00:00	02:00:00:00:01:00	803	2.11 60	QoS Null function (No data), SN
	8 0.114788	02:00:00:00:01:00	Broadcast	803	2.11 79	QoS Data, SN=42, FN=1, Flags=.p
	9 0.115859				2.11 24	Power-Save poll, Flags=
	10 0.115871	Success: client re	egistered in PS-m	node. 802	2.11 60	QoS Null function (No data), SN
	11 0.122050			803	2.11 78	QoS Data, SN=42, FN=1, Flags=.p
	12 0.123099	02:00:00:00:01:00	02:00:00:00:00:00 (0	02:00:00:00:00:00) 80	2.11 24	Power-Save poll, Flags=
	13 0.174995	02:00:00:00:01:00	Broadcast	803	2.11 78	QoS Data, SN=42, FN=1, Flags=.p
	14 0.176063	02:00:00:00:01:00	02:00:00:00:00:00 (0	02:00:00:00:00:00) 80	2.11 24	Power-Save poll, Flags=
	15 0.227746	02:00:00:00:01:00	Broadcast	803	2.11 78	QoS Data, SN=42, FN=1, Flags=.p
	16 0.228729	02:00:00:00:01:00	02:00:00:00:00:00 (0	02:00:00:00:00:00) 80	2.11 24	Power-Save poll, Flags=
		Apply a display filter No. Time 1 0.000000 2 0.001077 3 0.053206 4 0.054175 5 0.107561 6 0.108973 7 0.109012 8 0.114788 9 0.115859 10 0.115871 11 0.122050 12 0.123099 13 0.174995 14 0.176063 15 0.227746	Apply a display filter < ℜ/>No. Time Source 1 0.000000 02:00:00:00:01:00 2 0.001077 02:00:00:00:01:00 3 0.053206 02:00:00:00:01:00 4 0.054175 02:00:00:00:01:00 5 0.107561 02:00:00:00:01:00 6 0.108973 02:00:00:00:01:00 7 0.109012 02:00:00:00:01:00 8 0.114788 02:00:00:00:01:00 9 0.115859 02:00:00:00:01:00 10 0.122050 02:00:00:00:01:00 13 0.174995 02:00:00:00:01:00 14 0.176063 02:00:00:00:01:00 15 0.227746 02:00:00:00:01:00	Apply a display filter <\%/>No. Time Source Destination 1 0.000000 02:00:00:00:01:00 Broadcast 2 0.001077 02:00:00:00:01:00 Broadcast 3 0.053206 02:00:00:00:01:00 Broadcast 4 0.054175 02:00:00:00:01:00 Broadcast 5 0.107561 02:00:00:00:01:00 Broadcast 6 0.108973 02:00:00:00:01:00 Broadcast 7 0.109012 02:00:00:00:01:00 Broadcast 9 0.114788 02:00:00:00:01:00 Broadcast 9 0.115871 02:00:00:00:01:00 Broadcast 10 0.122050 02:00:00:00:01:00 Broadcast 11 0.122050 02:00:00:00:01:00 02:00:00:00:00:00:00 13 0.174995 02:00:00:00:01:00 Broadcast 14 0.176063 02:00:00:00:01:00 02:00:00:00:00:00 15 0.227746 02:00:00:00:01:00 Broadcast	▲ ▲	Image: Constraint of the second s

Frame 1: 79 bytes on wire (632 bits), 79 bytes captured (632 bits) Radiotap Header v0, Length 8 802.11 radio information IEEE 802.11 QoS Data, Flags: .p.P...T ▶ Data (37 bytes)

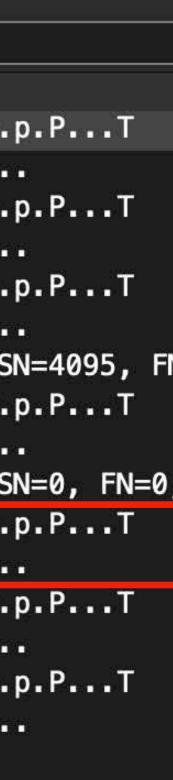




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No.	Time	Source	Destination			Protocol	to the set of the set	Info	
1	0.00000	02:00:00:00:01:00	Broadcast			802.11	79	QoS Data, SN=42,	FN=1, Flags=.p
2	0.001077	02:00:00:00:01:00	02:00:00:00:00:00	(02:00:00:00:	00:00)	802.11	24	Power-Save poll,	Flags=
3	0.053206	02:00:00:00:01:00	Broadcast			802.11	79	QoS Data, SN=42,	FN=1, Flags=.p
4	0.054175	02:00:00:00:01:00	02:00:00:00:00:00	(02:00:00:00:	00:00)	802.11	24	Power-Save poll,	Flags=
5	0.107561	02:00:00:00:01:00	Broadcast			802.11	79	QoS Data, SN=42,	FN=1, Flags=.p
6	0.108973	02:00:00:00:01:00	02:00:00:00:00:00	(02:00:00:00:	00:00)	802.11	24	Power-Save poll,	Flags=
7	0.109012	02:00:00:00:00:00	02:00:00:00:01:00			802.11	60	QoS Null function	on (No data), SN
8	0.114788	02:00:00:00:01:00	Broadcast			802.11	79	QoS Data, SN=42,	FN=1, Flags=.p
9	0.115859	02:00:00:00:01:00	02:00:00:00:00:00	(02:00:00:00:	00:00)	802.11	24	Power-Save poll,	Flags=
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12	0.123099	02:00:00:00:01:00	02:00:00:00:00:00	(02:00:00:00:	00:00)	802.11	24	Power-Save poll,	Flags=
13	0.174995	02:00:00:00:01:00	Broadcast			802.11	78	QoS Data, SN=42,	FN=1, Flags=.p
14	0.176063				00:00)	802.11	24	Power-Save poll,	Flags=
15	0.227746	Failure: client not	registered in F	PS-mode.		802.11	78	QoS Data, SN=42,	FN=1, Flags=.p
16	0.228729				00:00)	802.11	24	Power-Save poll,	Flags=

 Frame 1: 79 bytes on wire (632 bits), 79 bytes captured (632 bits)
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 Data (37 bytes) 🖉 attack.pcap

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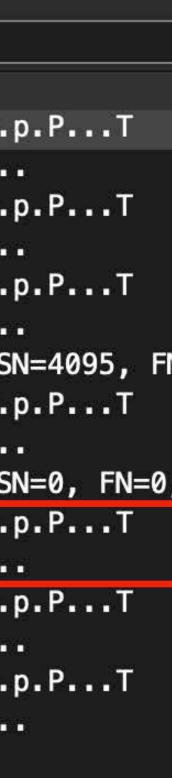
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5	0.107561	02:00:00:00:01:00	Broadcast		802.11	79 QoS 🗖 ata	a, SN=42, F	FN=1, Flags=.p
6	0.108973	02:00:00:00:01:00	02:00:00:00:00:00	(02:00:00:00:00:00)	802.11	24 Por er-Sa	av poll, F	=lags=
7	0.109012	02:00:00:00:00:00	02:00:00:00:01:00		802.11	60 S Nu ¹	function	(No data), SN
8	0.114788	02:00:00:00:01:00	Broadcast		802.11	79 QoS 🗖 ata	a, SN=42, F	N=1, Flags=.p
9	0.115859	02:00:00:00:01:00	02:00:00:00:00:00	(02:00:00:00:00:00)	802.11	24 Poter-Sa	ave poll, F	ags=
10	0.115871	02:00:00:00:00:00	02:00:00:00:01:00		802.11	60 🥂 S Nul	function	(No data), SN
11	0.122050	02:00:00:00:01:00	Broadcast		802.11	78 QoS Data	a, SN=42, F	™=1, Flags=.p
12	0.123099	02:00:00:00:01:00	02:00:00:00:00:00	(02:00:00:00:00:00)	802.11	24 Power-Sa	ave poll, F	-lags=
13	0.174995	02:00:00:00:01:00	Broadcast		802.11	78 QoS Data	a, SN=42, F	N=1, Flags=.p
14	0.176063						ve poll, F	=lags=
15	0.227746	We decrypted the	e previous byte	e, so we can cho	p off ano	ther one.	, SN=42, F	N=1, Flags=.p
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 Frame 1: 79 bytes on wire (632 bits), 79 bytes captured (632 bits)
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Conclusion

- Large support for WPA-TKIP.
- We found several novel side-channels in WPA-TKIP implementations.
 - Implementing WPA-TKIP without side-channel vulnerabilities is hard.
- We should do more cross-layer analysis to find vulnerabilities.
- Hopefully we accelerate the process of deprecating support for WPA-TKIP.

Conclusion

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- We found several novel side-channels in WPA-TKIP implementations.
 - Implementing WPA-TKIP without side-channel vulnerabilities is hard.
- We should do more cross-layer analysis to find vulnerabilities.
- Hopefully we accelerate the process of deprecating support for WPA-TKIP. \bullet

RECOMMENDATION: Abandon WPA-TKIP, and move towards WPA3.



Features we're no longer developing

We're no longer actively developing these features and may remove them from a future update. Some features have been replaced with other features or functionality, while others are now available from different sources.

If you have feedback about the proposed replacement of any of these features, you can use the Feedback Hub app.

_	Feature	Details
	Taskbar settings roaming	Roaming of taskbar settings is no longer being de
	Wi-Fi	In this release a warning message will appear whe
	WEP and	as secure as those using WPA2 or WPA3. In a futu
	TKIP	be disallowed. Wi-Fi routers should be updated to
	Windows	Windows To Go is no longer being developed.
	To Go	
		The feature does not support feature updates and type of USB that is no longer supported by many

leveloped and we plan to disable this capability in a future release

en connecting to Wi-Fi networks secured with WEP or TKIP, which are not ure release, any connection to a Wi-Fi network using these old ciphers will to use AES ciphers, available with WPA2 or WPA3.

Ind therefore does not enable you to stay current. It also requires a specific y OEMs.



THANK YOU Happy to answer your questions.

Practical Side-Channel Attacks against WPA-TKIP

2019 Black Hat Europe



